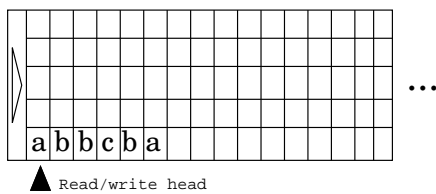


12-0: **Extending Turing Machines**

- When we added a stack to NFA to get a PDA, we increased computational power
- Can we do the same thing for Turing Machines?
  - That is, can we add some new “feature” to TMs that will increase their computational power?

12-1: **Multi-Track Tape**

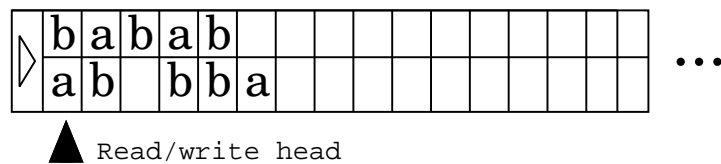
- Instead of each tape location holding a single symbol, we add several “tracks” to the tape
  - Based on contents of all tracks, either move head left, move head right, or write new values to any of the tracks



12-2: **Multi-Track Tape**

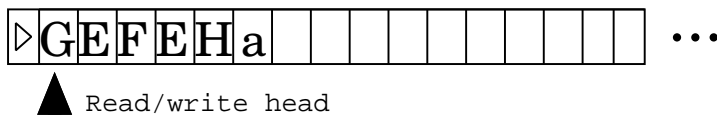
- Can simulate a mutli-track machine with a standard TM
  - Increase the size of the tape alphabet
  - $k$  tracks, each with an alphabet of  $n$  symbols
  - New alphabet of size  $n^k$

12-3: **Multi-Track Tape**



$$= \begin{bmatrix} \\ \end{bmatrix} \quad a = \begin{bmatrix} \\ a \end{bmatrix} \quad b = \begin{bmatrix} \\ b \end{bmatrix} \quad C = \begin{bmatrix} a \\ \end{bmatrix} \quad D = \begin{bmatrix} a \\ a \end{bmatrix}$$

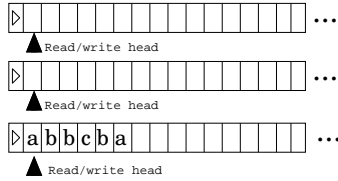
$$E = \begin{bmatrix} a \\ b \end{bmatrix} \quad F = \begin{bmatrix} b \\ \end{bmatrix} \quad G = \begin{bmatrix} b \\ a \end{bmatrix} \quad H = \begin{bmatrix} b \\ b \end{bmatrix}$$



12-4: **Multiple Tapes**

- Several tapes, with independent read/write heads
- Reach symbol on each tape, and based on contents of all tapes:

- Write or move each tape independently
- Transition to new state



12-5: Multiple Tapes

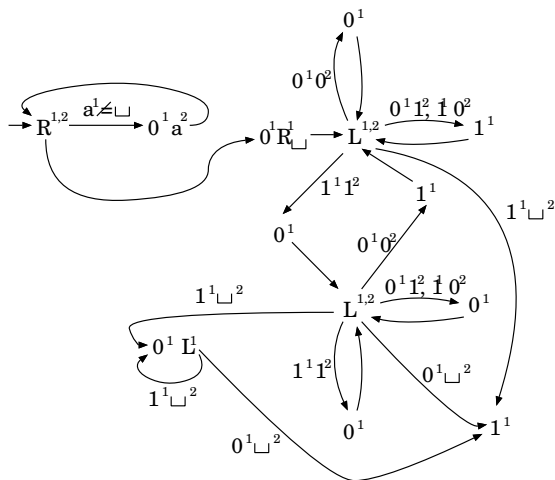
- Create a 2-Tape Machine that adds two numbers
  - Convert  $\triangleright \sqcup w; v$  to  $\triangleright \sqcup w + v$  (leading zeros OK)
- Assume that tape 1 holds input (and output), and tape 2 starts out with blanks

12-6: Multiple Tapes

- Create a 2-Tape Machine that adds two numbers
  - Convert  $\triangleright \sqcup w; v$  to  $\triangleright \sqcup w + v$  (leading zeros OK)
- Copy first # to second tape (zeroing out first # on first tape)
- Do “standard addition”, keeping track of carries.

12-7: Multiple Tapes

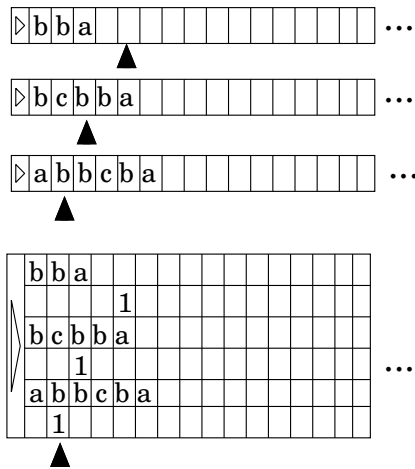
- Create a 2-Tape Machine that adds two numbers



12-8: Multiple Tapes

- Are  $k$ -tape machines more powerful than 1-tape machines?

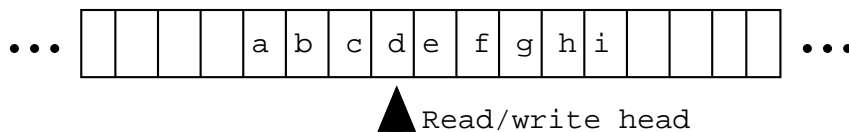
12-9: Multiple Tapes



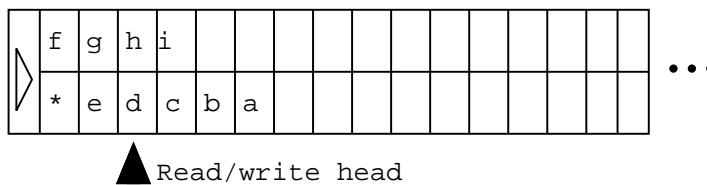
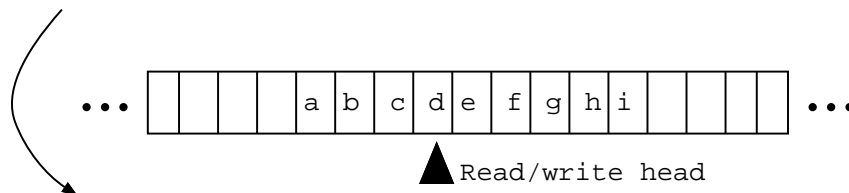
12-10: Multiple Tapes

- Each transition from the original, multi-tape machine will require several transitions from the simulated machine – and each state in the multiple-tape machine will be represented by a set of states in the simulation machine
  - First, need to scan tape head to find all “virtual heads”, and remember what symbol is stored at each head location
    - Use state to store this information
  - Next, scan tape to implement the action on each tape (moving head, rewriting symbols, etc)
  - Finally, transition to a new set of states

12-11: 2-Way Infinite Tape



12-12: 2-Way Infinite Tape



12-13: 2-Way Infinite Tape

- Make 2 copies of states in original machine: One set for top tape, one set for bottom tape
- Top Tape States
  - Use the top track
  - Execute as normal
  - When “Move Left” command, and beginning of tape symbol is on the bottom tape, move Right instead, switch to Bottom Tape States

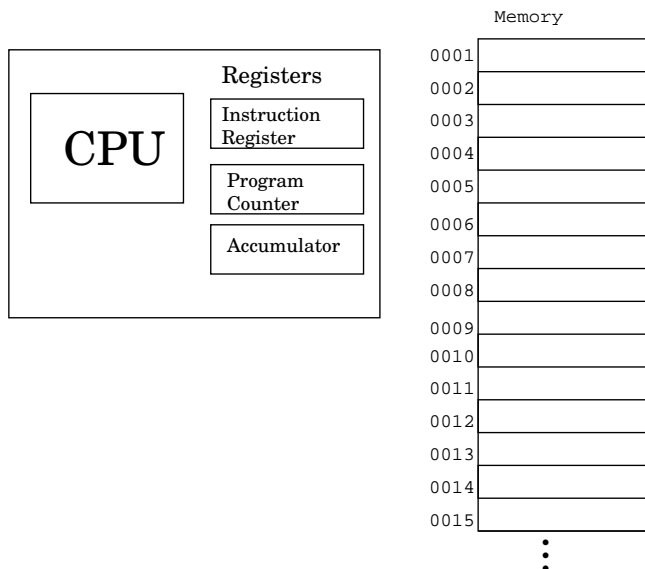
12-14: **2-Way Infinite Tape**

- Make 2 copies of states in original machine: One set for top tape, one set for bottom tape
- Bottom Tape States
  - Use the bottom track
  - Move left on a “Move Right” command, move right on a “Move Left” command
  - When the beginning of tape symbol is encountered, switch to Top Tape States

12-15: **Simple Computer**

- CPU
- 3 Registers (Instruction Register (IR), Program Counter (PC), Accumulator (ACC))
- Memory
- Operation:
  - Set IR → MEM[PC]
  - Increment PC
  - Execute instruction in IR
  - Repeat

12-16: **Simple Computer**



12-17: Simple Computer

	Instruction	Meaning
00	HALT	Stop Computation
01	LOAD x	ACC ← MEM[x]
02	LOADI x	ACC ← x
03	STORE x	MEM[x] ← ACC
04	ADD x	ACC ← ACC + MEM[x]
05	ADDI x	ACC ← ACC + x
06	SUB x	ACC ← ACC - MEM[x]
07	SUBI x	ACC ← ACC - x
08	JUMP x	IP ← x
09	JZERO x	IP ← x if ACC = 0
10	JGT x	IP ← x if ACC > 0

Write a program that multiplies two numbers (in locations 1000 & 1001), and stores the result in 1002 12-18:

Simple Computer

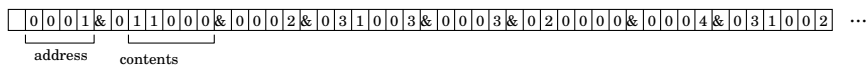
Memory	Machine Code	Assembly
0001	011000	LOAD 1000
0002	031003	STORE 1003
0003	020000	LOADI 0
0004	031002	STORE 1002
0005	021003	LOAD 1003
0006	090012	JZERO 0012
0007	070001	SUBI 1
0008	031003	STORE 1003
0009	011002	LOAD 1002
0010	041001	ADD 1001
0011	080004	STORE 1002
0012	000000	HALT

12-19: Computers & TMs

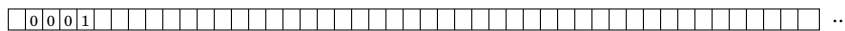
- We can simulate this computer with a multi-tape Turing machine:
  - One tape for each register (IR, IP, ACC)
  - One tape for the Memory
    - Memory tape will be entries of the form <address> <contents>

12-20: Computers & TMs

Memory



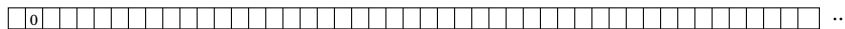
Instruction Pointer



Instruction Register



Accumulator

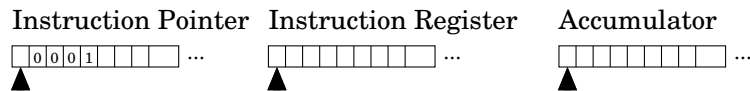
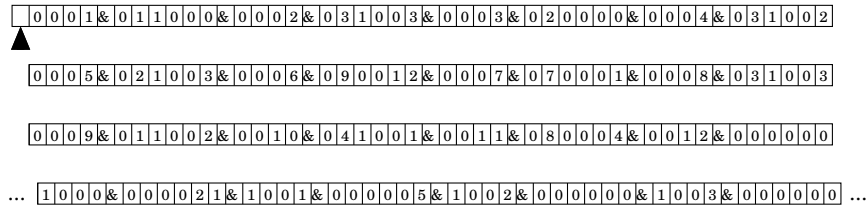


12-21: Computers & TMs

- Operation:
  - Scan through memory until reach an address that matches the IP
  - Copy contents of memory at that address to the IR
  - Increment IP
  - Based on the instruction code:
    - Copy value into IP
    - Copy a value into Memory
    - Copy a value into the ACC
    - Do addition/subtraction

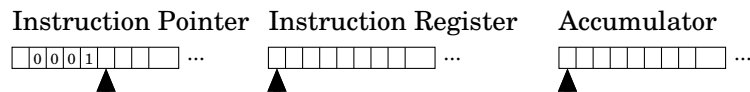
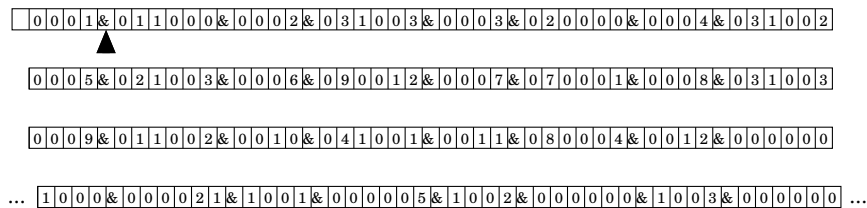
12-22: Computers & TMs

Memory



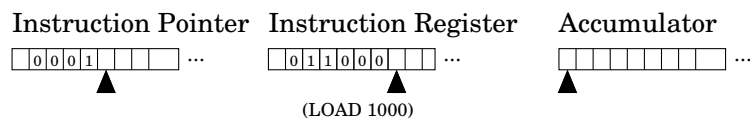
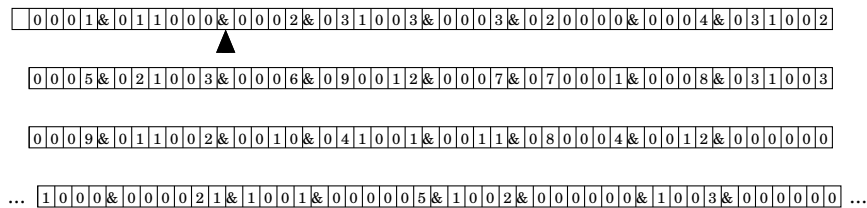
12-23: Computers & TMs

Memory



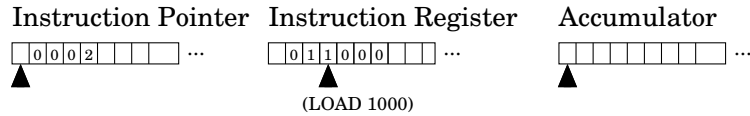
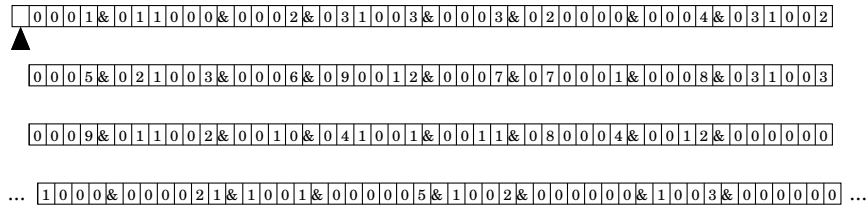
12-24: Computers & TMs

Memory



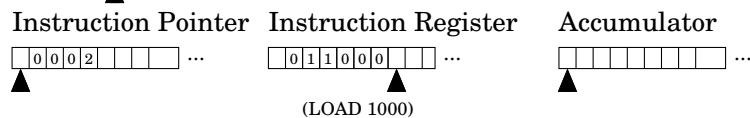
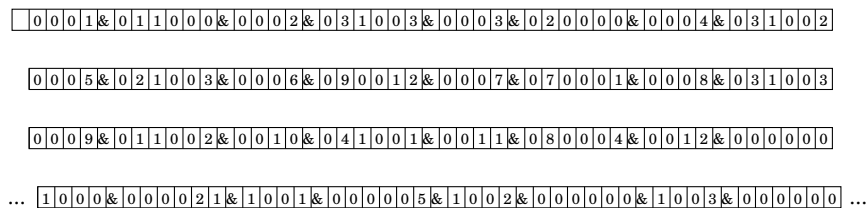
12-25: Computers & TMs

Memory



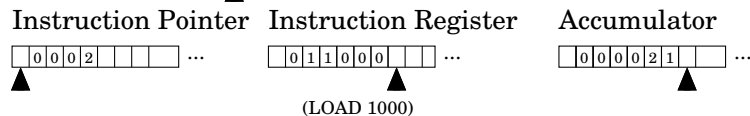
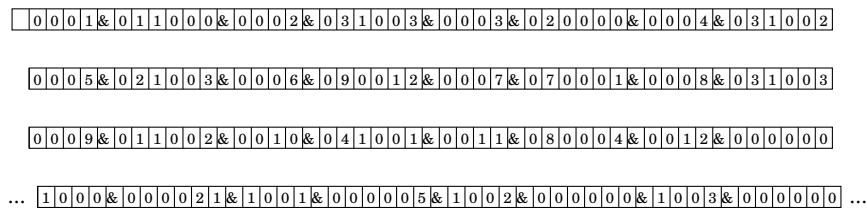
12-26: Computers & TMs

Memory



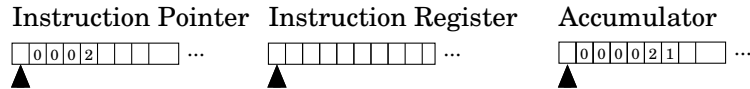
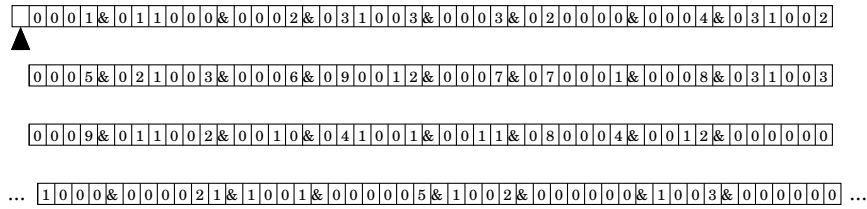
12-27: Computers & TMs

Memory



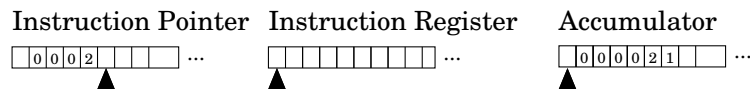
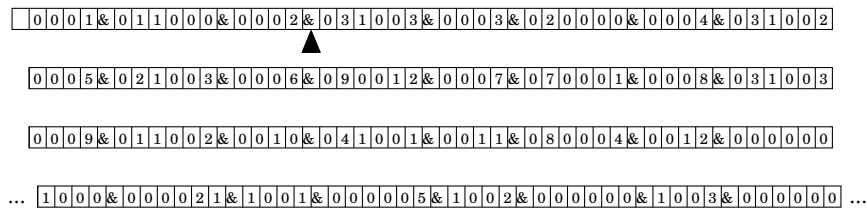
12-28: Computers & TMs

Memory



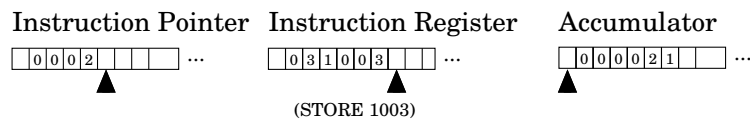
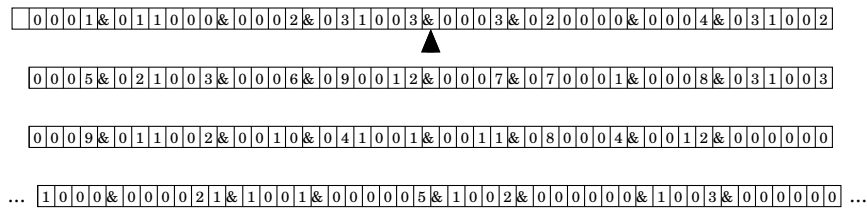
12-29: Computers & TMs

Memory



12-30: Computers & TMs

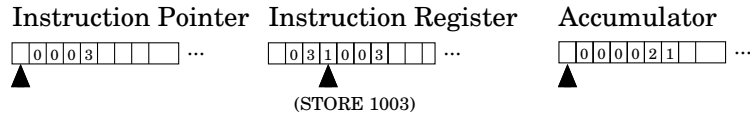
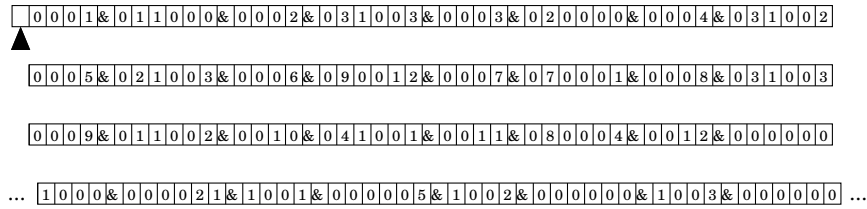
Memory



12-31: Computers & TMs

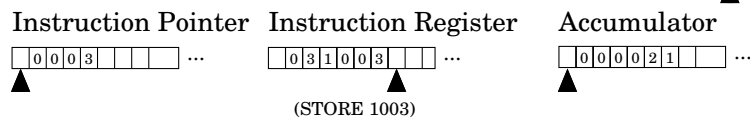
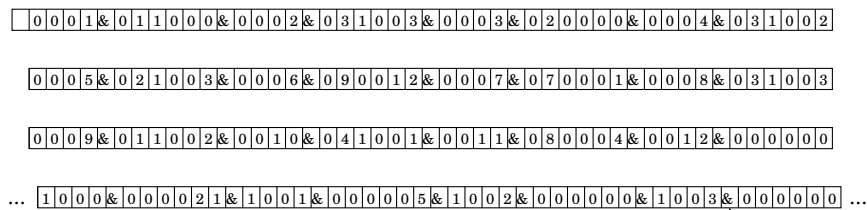


Memory



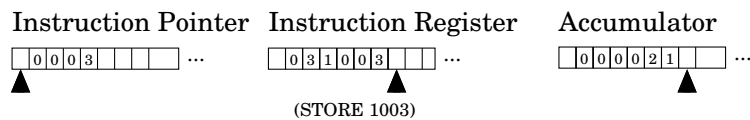
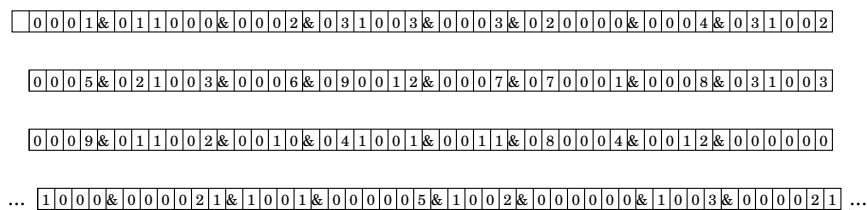
12-32: Computers & TMs

Memory

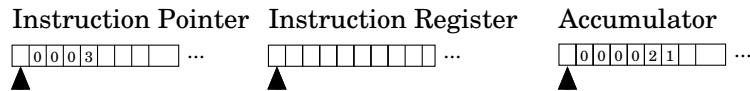
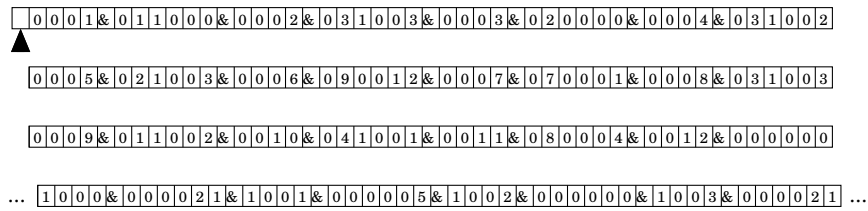


12-33: Computers & TMs

Memory



12-34: Computers & TMs

**Memory****12-35: Computers & TMs**

- “Simple Computer” can be modeled by a Turing Machine
- Any current machine can be modeled in the same way by a Turing Machine
- If there is an algorithm for it, a Turing Machine can do it
  - Note that at this point, we don’t care *how long* it might take, just that it can be done

**12-36: Turing Complete**

- A computation formalism is “Turing Complete” if it can simulate a Turing Machine
- Turing Complete  $\Rightarrow$  can compute anything
  - Of course it might not be convenient ...

**12-37: Non-Determinism**

- Final extension to Turing Machines: Non-Determinism
  - Just like non-determinism in NFAs, PDAs
  - String is accepted by a non-deterministic Turing Machine if there is at least one computational path that accepts

**12-38: Non-Determinism**

A Non-Deterministic Machine  $M$  *Decides* a language  $L$  if:

- All computational paths halt
- For each  $w \in L$ , at least one computational path for  $w$  accepts
- For all  $w \notin L$ , no computational path accepts

**12-39: Non-Determinism**

A Non-Deterministic Machine  $M$  *Semi-Decides* a language  $L$  if:

- For each  $w \in L$ , at least one computational path for  $w$  halts and accepts
- For all  $w \notin L$ , no computational path halts and accepts

**12-40: Non-Determinism**

A Non-Deterministic Machine  $M$  *Computes a Function* if:

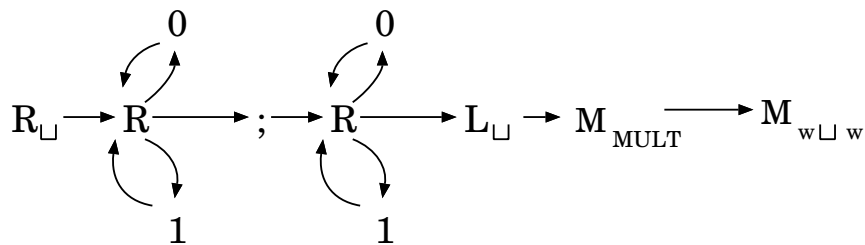
- All computational paths halt
- Every computational path produces the same result

12-41: **Non-Determinism**

- Non-Deterministic TM for  $L = \{w \in \{0,1\} : w \text{ is composite}\}$
- (semi-decides is OK)

12-42: **Non-Determinism**

- Non-Deterministic TM for  $L = \{w \in \{0,1\} : w \text{ is composite}\}$



How could we make this machine decide (instead of semi-decide)  $L$ ? 12-43: **Non-Determinism**

How we can make this machine decide (instead of semi-decide)  $L$

- First, transform  $w$  into  $w \sqcup w; w$
- Non-deterministically modify the second  $2 w$ 's
- Multiply the second  $2 w$ 's
- Check to see if the resulting string is  $w \sqcup w$

12-44: **Non-Determinism**

- Are Non-Deterministic Turing Machines more powerful than Deterministic Turing machines?
  - Is there some  $L$  which can be semi-decided by a non-deterministic Turing Machine, which cannot be semi-decided by a Deterministic Turing Machine?
- Non-determinism in Finite Automata didn't buy us anything
- Non-determinism in Push-Down Automata did

12-45: **Non-Determinism**

- How to Simulate a Non-Deterministic Turing Machine with a Deterministic Turing Machine

12-46: **Non-Determinism**

- How to Simulate a Non-Deterministic Turing Machine with a Deterministic Turing Machine
  - Try one computational path – if it says yes, halt and say yes. Otherwise, try a different computational path. Repeat until success

12-47: **Non-Determinism**

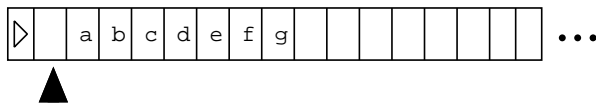
- How to Simulate a Non-Deterministic Turing Machine with a Deterministic Turing Machine
  - Try one computational path – if it says yes, halt and say yes. Otherwise, try a different computational path. Repeat until success
    - But what if the first computational path runs forever . . .

12-48: **Non-Determinism**

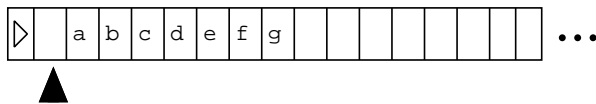
- How to Simulate a Non-Deterministic Turing Machine with a Deterministic Turing Machine
  - Try all computational paths of length 1
  - Try all computational paths of length 2
  - Try all computational paths of length 3
  - . . .
- If there is a halting configuration, you will find it eventually. Otherwise, run forever.

12-49: **Non-Determinism**

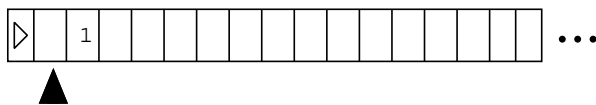
**Original Tape**



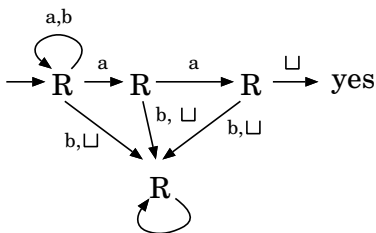
**Work Tape**



**Control Tape**

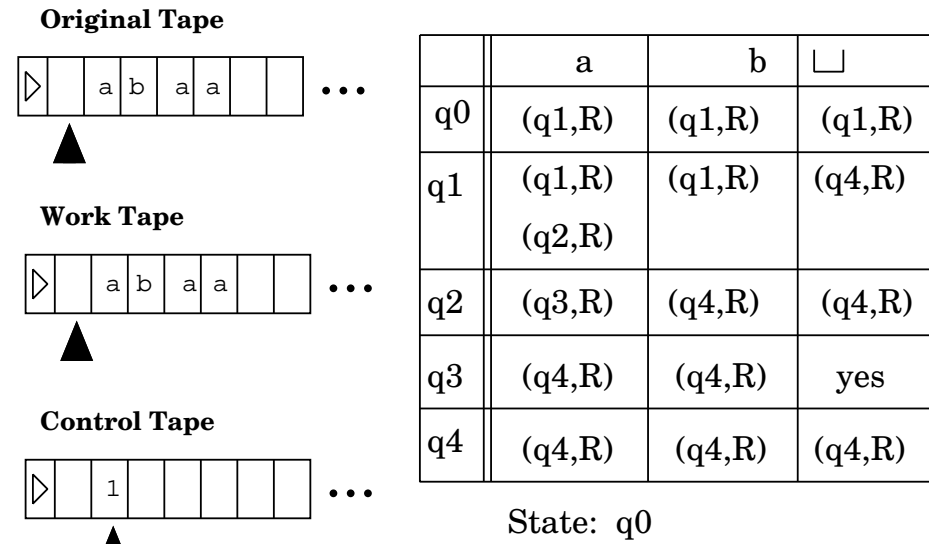


12-50: **Non-Determinism**

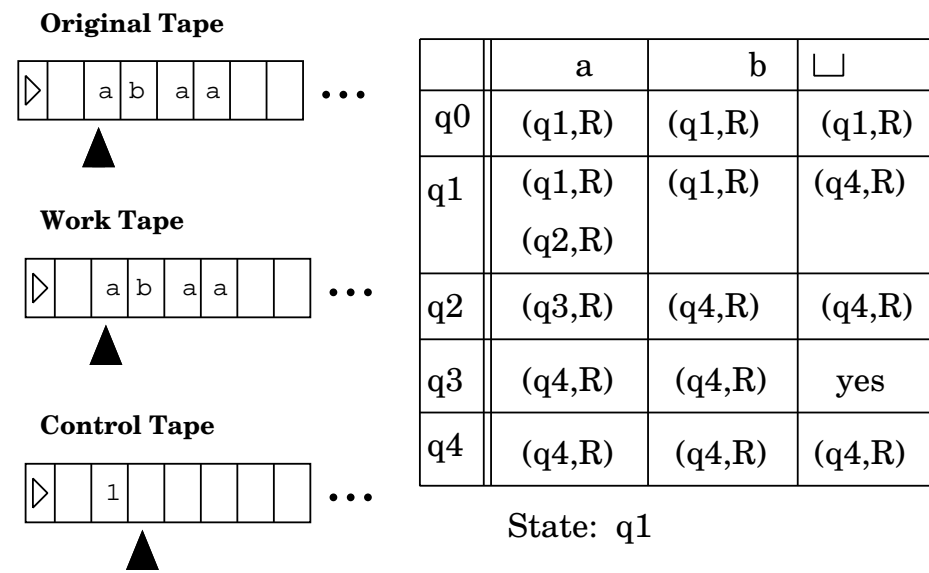


	a	b	$\square$
q0	(q1,R)	(q1,R)	(q1,R)
q1	(q1,R)	(q1,R)	(q4,R)
q2	(q3,R)	(q4,R)	(q4,R)
q3	(q4,R)	(q4,R)	yes
q4	(q4,R)	(q4,R)	(q4,R)

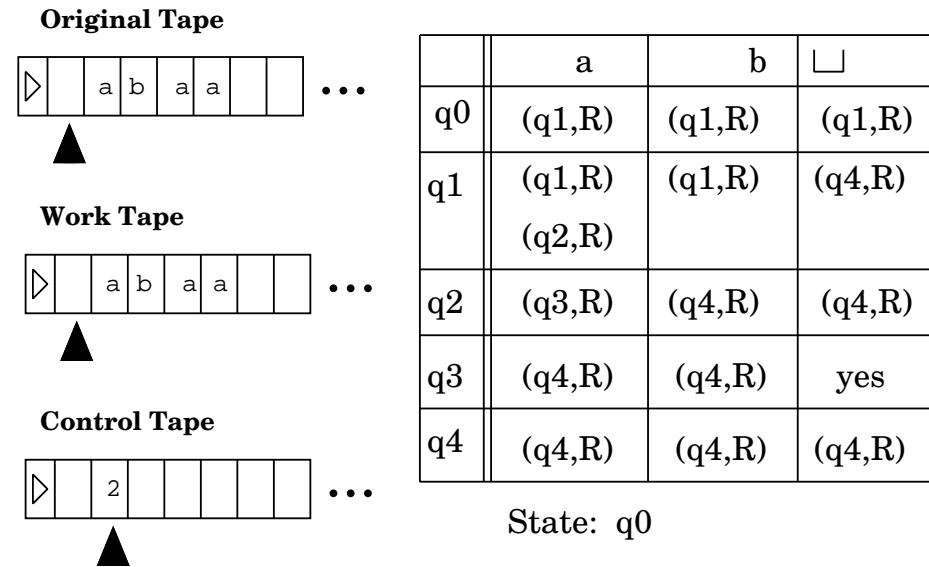
12-51: **Non-Determinism**



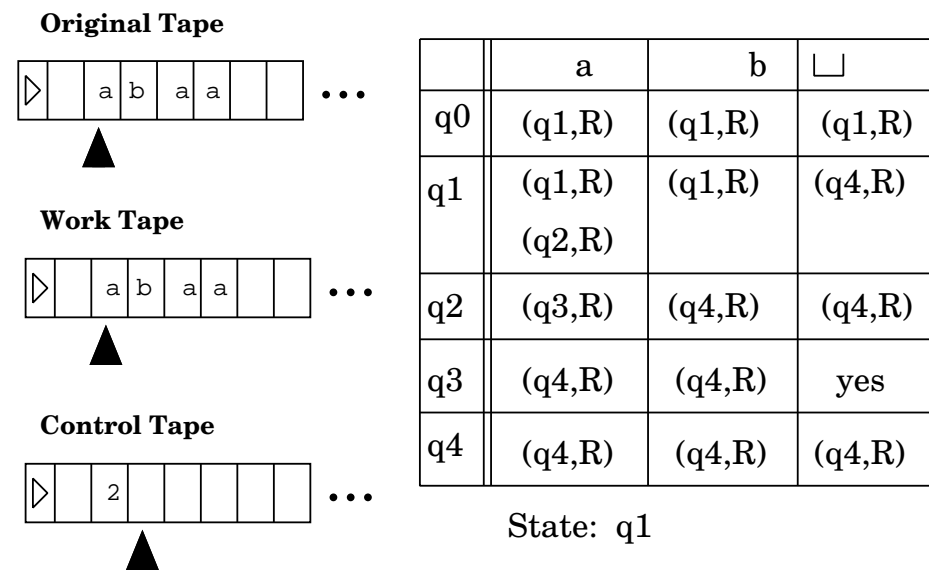
12-52: Non-Determinism



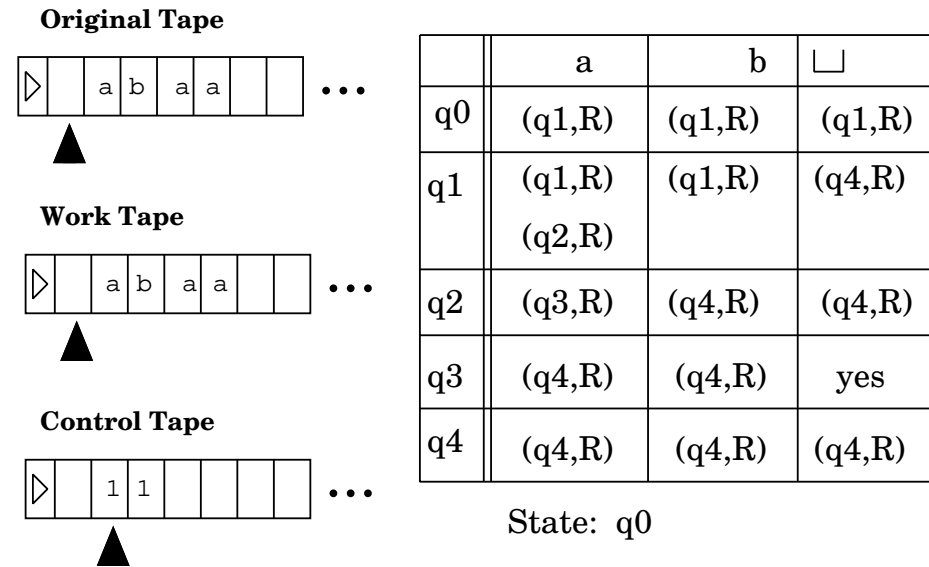
12-53: Non-Determinism



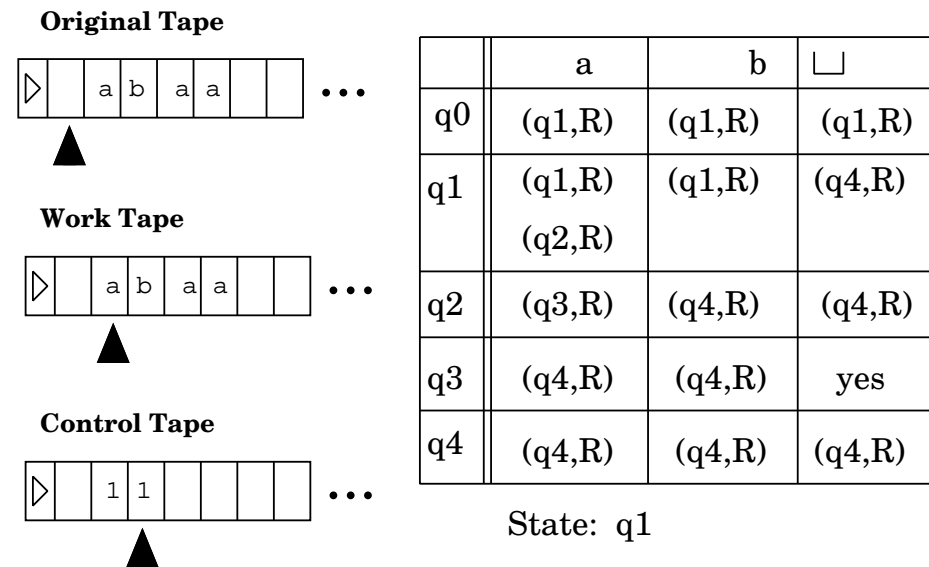
12-54: Non-Determinism



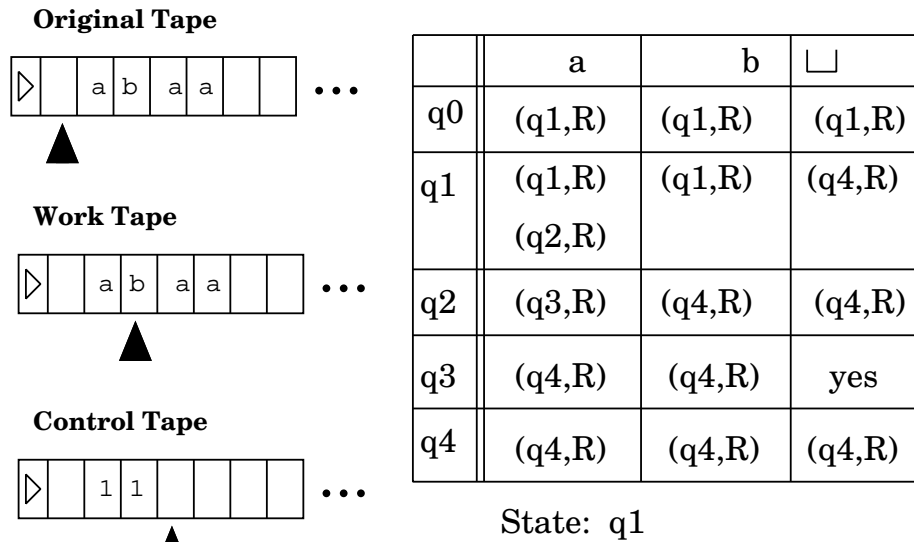
12-55: Non-Determinism



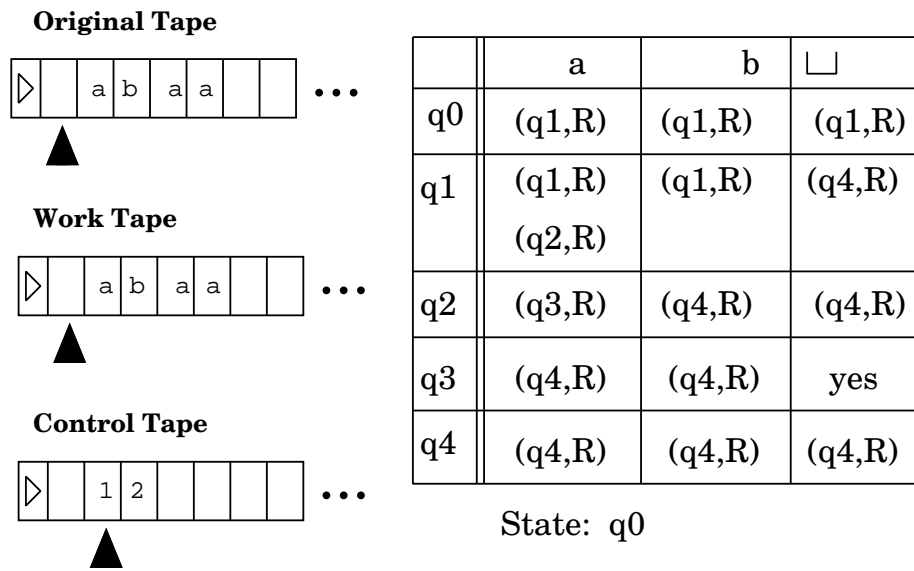
12-56: Non-Determinism



12-57: Non-Determinism

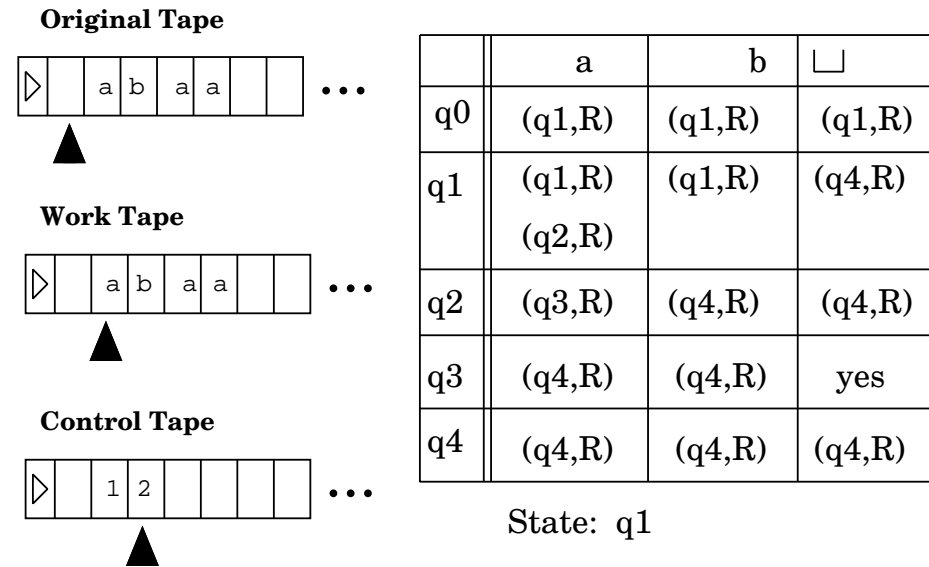


12-58: Non-Determinism

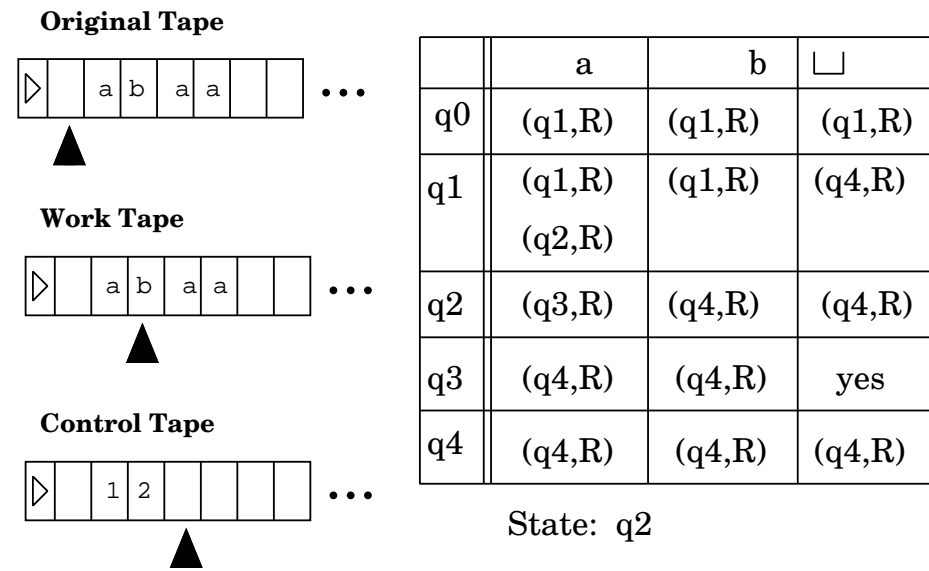


12-59: Non-Determinism

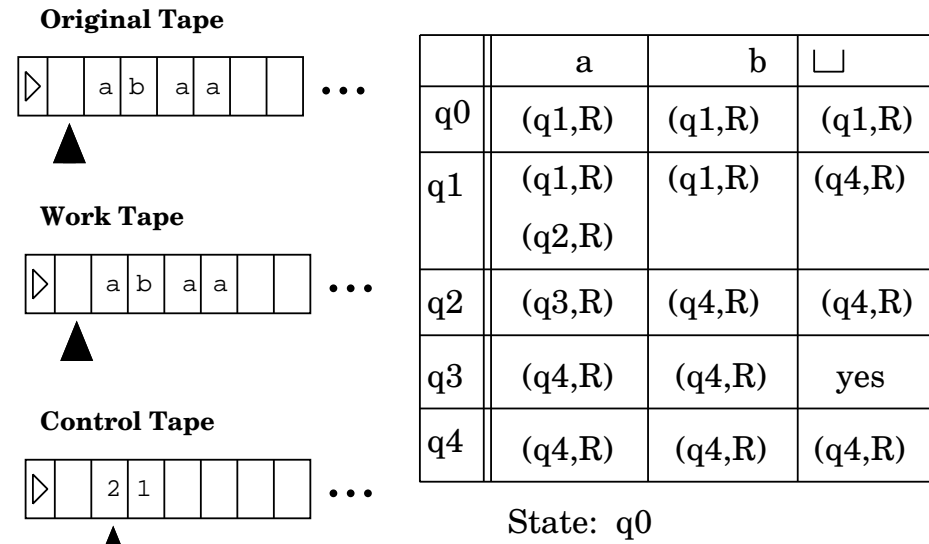




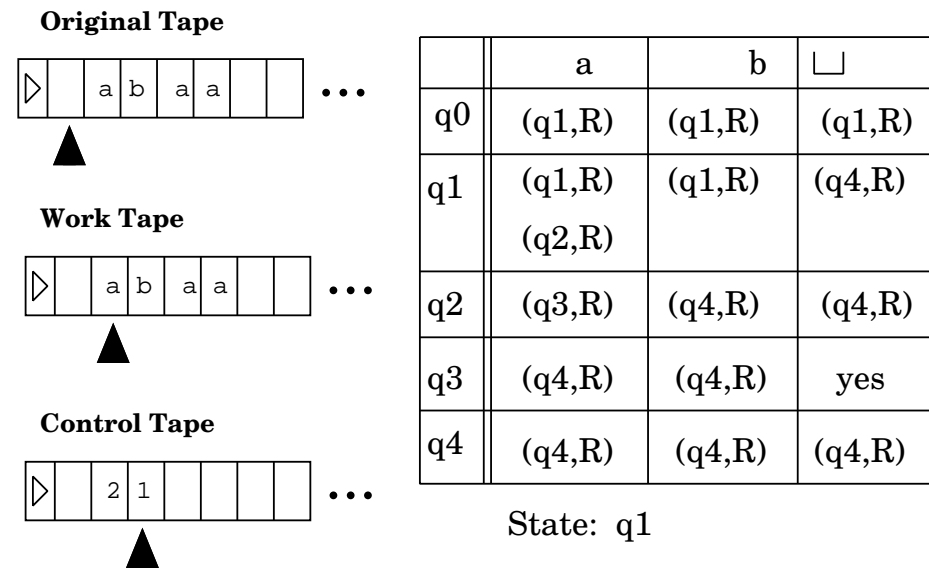
12-60: Non-Determinism



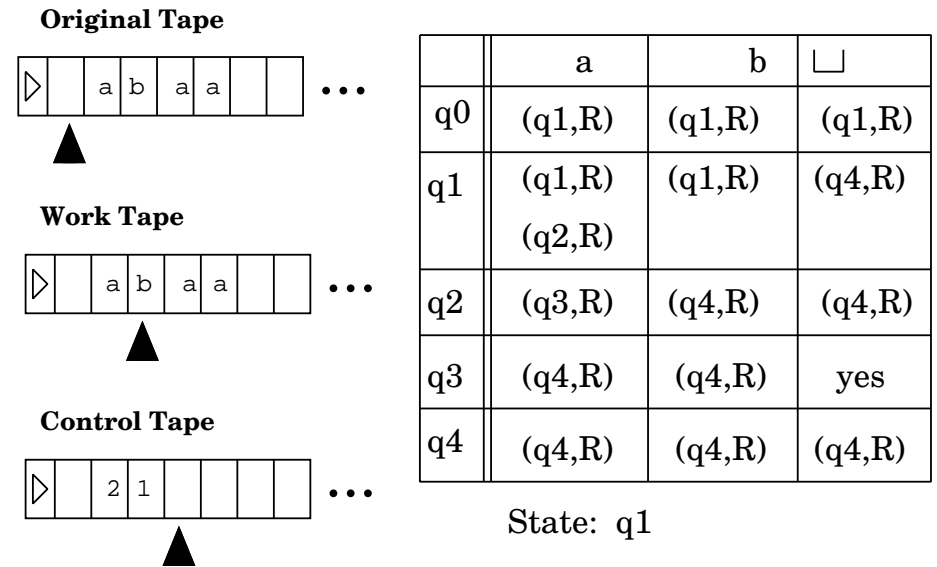
12-61: Non-Determinism



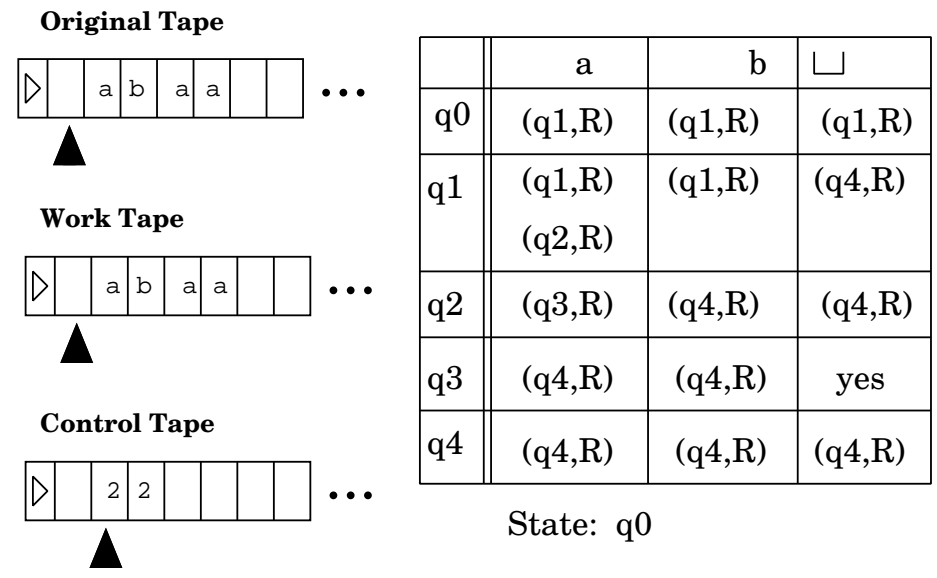
12-62: Non-Determinism



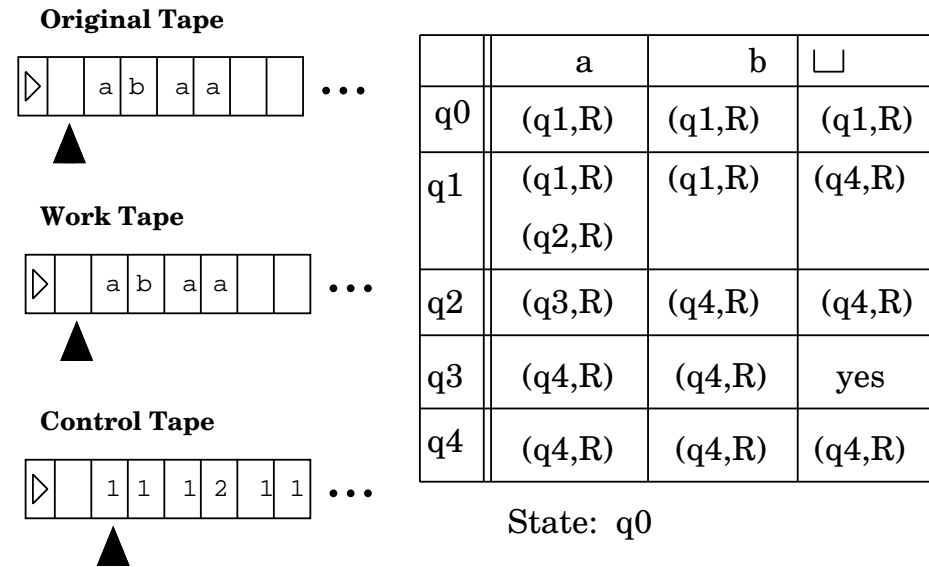
12-63: Non-Determinism



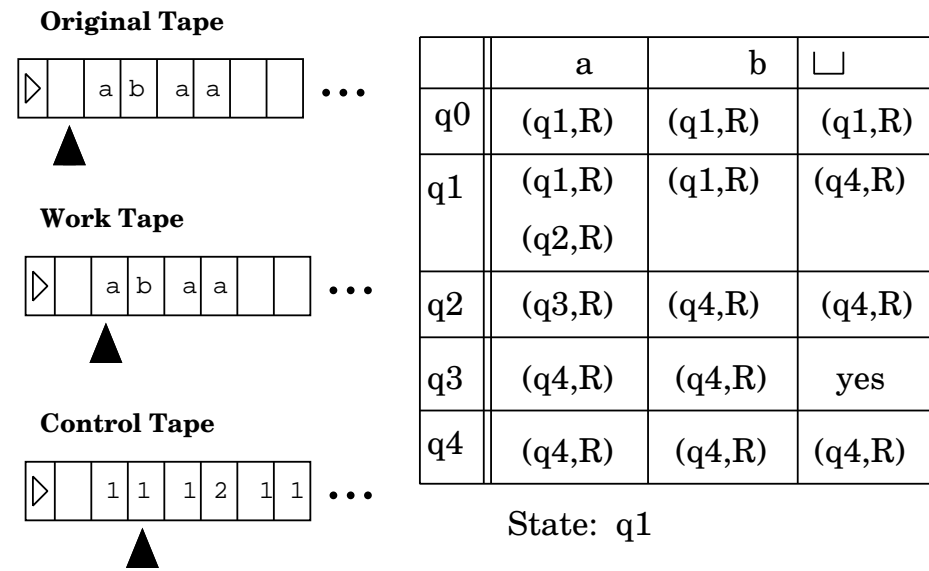
12-64: Non-Determinism



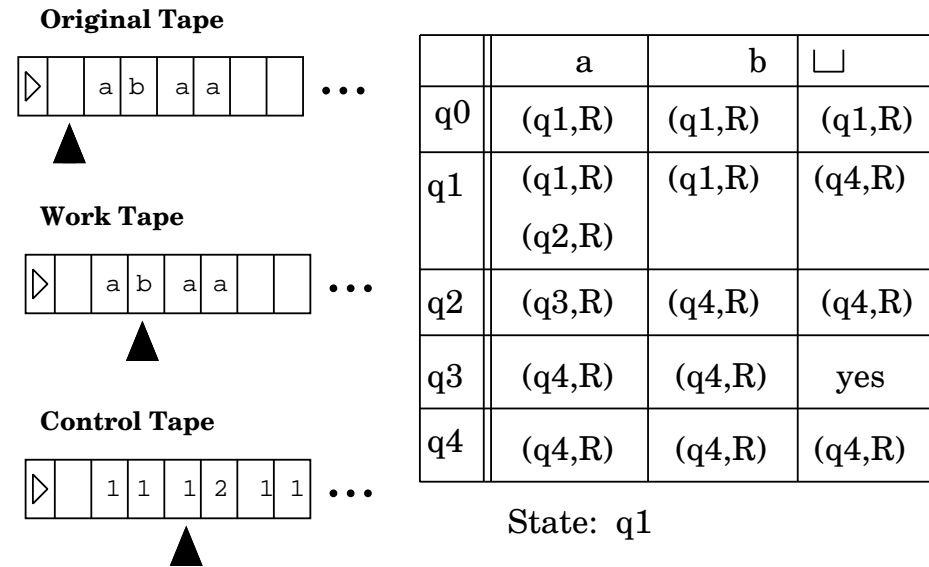
12-65: Non-Determinism



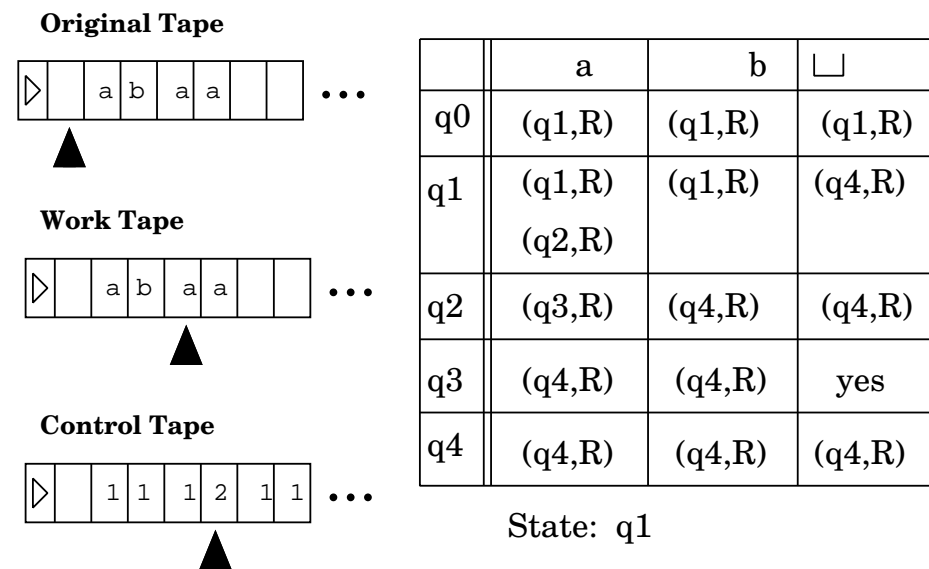
12-66: Non-Determinism



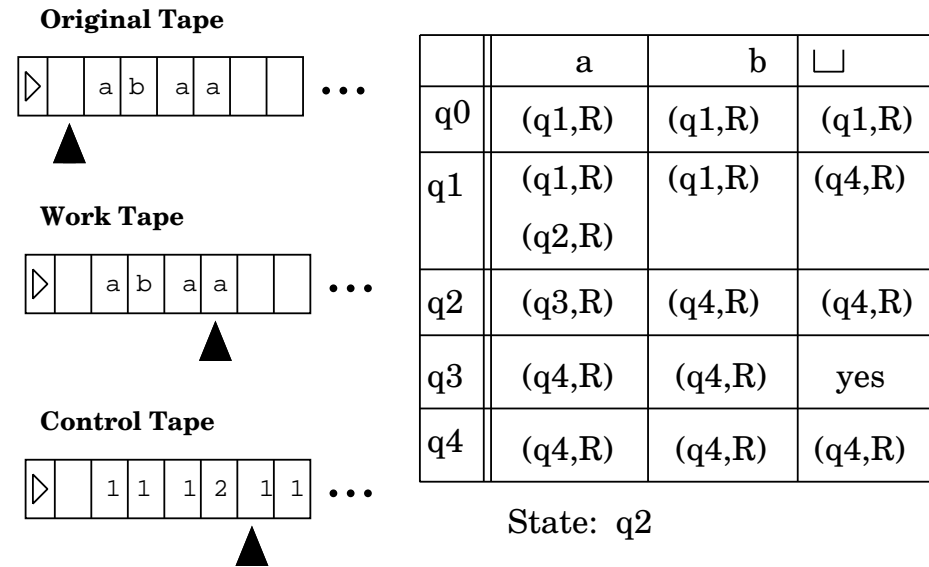
12-67: Non-Determinism



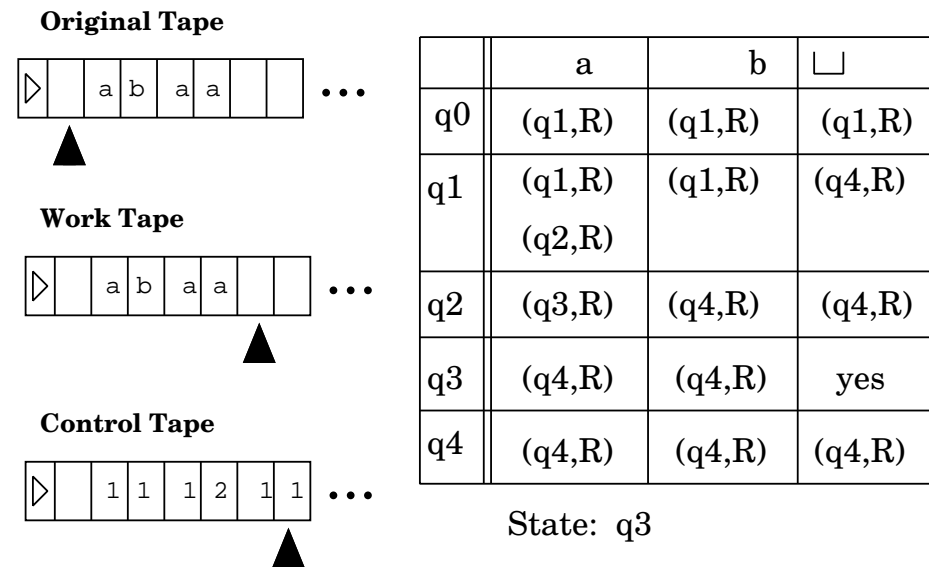
12-68: Non-Determinism



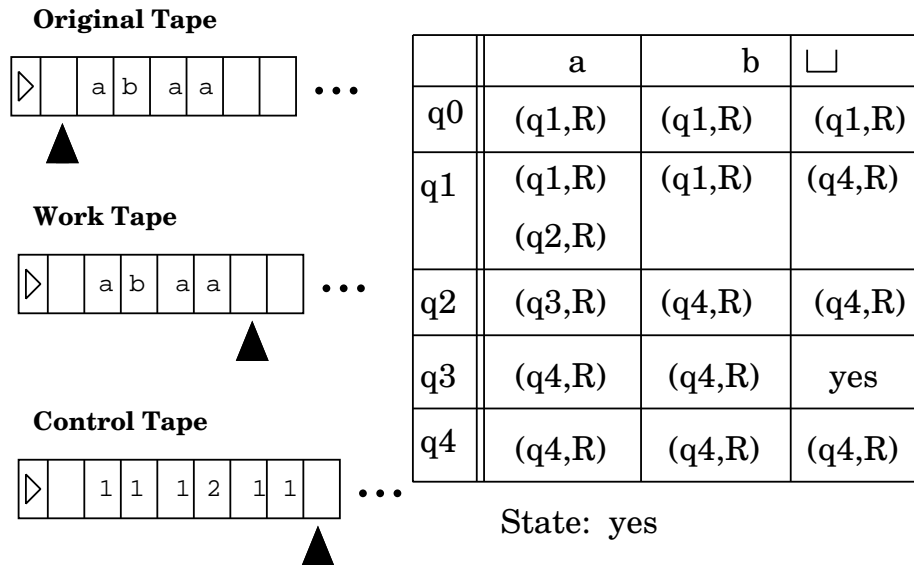
12-69: Non-Determinism



12-70: Non-Determinism



12-71: Non-Determinism

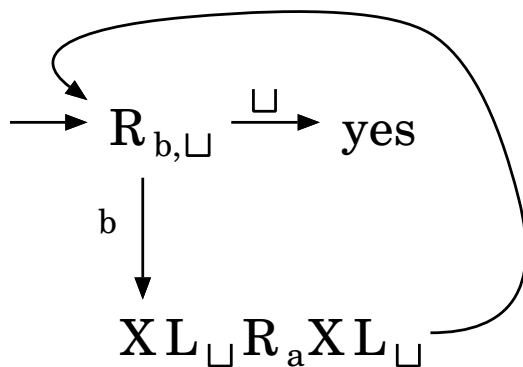


12-72: Turing Machines

- Some Turing Machine review problems:
  - Create a Turing Machine that *semi-decides* the language  $L =$  all strings over  $\{a, b\}$  with at least as many a's as b's

12-73: Turing Machines

- Create a Turing Machine that *semi-decides* the language  $L =$  all strings over  $\{a, b\}$  with at least as many a's as b's



12-74: Turing Machines

- Some Turing Machine review problems:
  - Create a Turing Machine that computes the function  $\lceil \lg x \rceil$ , where  $x$  is a binary number

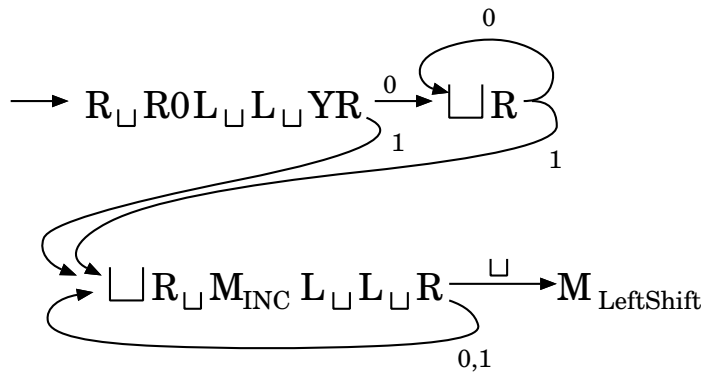
12-75: Turing Machines

- Some Turing Machine review problems:
  - Create a Turing Machine that computes the function  $\lceil \lg x \rceil$ , where  $x$  is a binary number

- Set result to 0
- While  $x \leq 2$ , divide  $x$  by 2, and add one to the result

12-76: Turing Machines

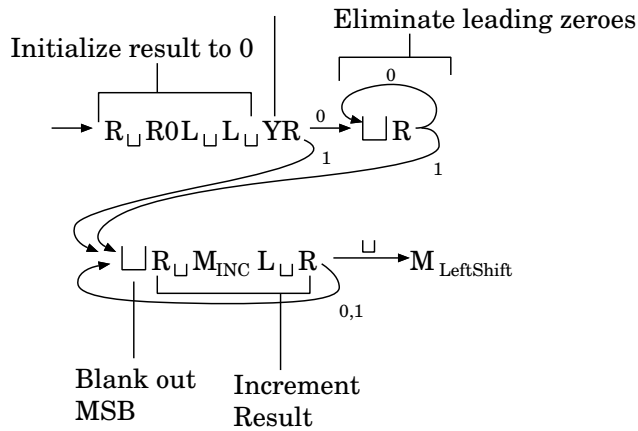
- Create a Turing Machine that computes the function  $\lceil \lg x \rceil$ , where  $x$  is a binary number



12-77: Turing Machines

- Create a Turing Machine that computes the function  $\lceil \lg x \rceil$ , where  $x$  is a binary number

Set marker for shifting at end of computation



12-78: Turing Machines

$M_{LeftShift}$

