# Data Structures and Algorithms CS245-2017S-10

#### Sorting

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### 10-0: Main Memory Sorting

- All data elements can be stored in memory at the same time
- Data stored in an array, indexed from  $0 \dots n 1$ , where n is the number of elements
- Each element has a key value (accessed with a key() method)
- We can compare keys for <, >, =
- For illustration, we will use arrays of integers though often keys will be strings, other Comparable types

### 10-1: Stable Sorting

- A sorting algorithm is *Stable* if the relative order of duplicates is preserved
- The order of duplicates matters if the *keys* are duplicated, but the *records* are not.

3	1	2	1	1	2	3	Key
B O b	J O e	E d	A m Y	S u e	A l	B u d	Data

1	1	1	2	2	3	3	Key
A	J	S	Е	A	В	В	Data
m	0	u	d	1	0	u	
У	e	е			b	d	

A non-Stable sort

#### 10-2: Insertion Sort

- Separate list into sorted portion, and unsorted portion
- Initially, sorted portion contains first element in the list, unsorted portion is the rest of the list
  - (A list of one element is always sorted)
- Repeatedly insert an element from the unsorted list into the sorted list, until the list is sorted

## 10-3: $\Theta()$ For Insertion Sort

- Running time  $\propto$  # of comparisons
- Worst Case:

## 10-4: $\Theta()$ For Insertion Sort

- Running time  $\propto$  # of comparisons
- Worst Case: Inverse sorted list
- # of comparisons:

## 10-5: $\Theta()$ For Insertion Sort

- Running time  $\propto$  # of comparisons
- Worst Case: Inverse sorted list

*#* of comparisons:

$$\sum_{i=1}^{n-1} i \in \Theta(n^2)$$

## 10-6: $\Theta()$ For Insertion Sort

- Running time  $\propto$  # of comparisons
- Best Case:

## 10-7: $\Theta()$ For Insertion Sort

- Running time  $\propto$  # of comparisons
- Best Case: Sorted List
- *#* of comparisons:

## 10-8: $\Theta()$ For Insertion Sort

- $\bullet\,$  Running time  $\propto$  # of comparisons
- Best Case: Sorted List

*#* of comparisons:

n-1

#### 10-9: Bubble Sort

- Scan list from the last index to index 0, swapping the smallest element to the front of the list
- Scan the list from the last index to index 1, swapping the second smallest element to index 1
- Scan the list from the last index to index 2, swapping the third smallest element to index 2
- Swap the second largest element into position (n-2)

### 10-10: $\Theta()$ for Bubble Sort

- $\bullet\,$  Running time  $\propto$  # of comparisons
- Number of Comparisons:

## 10-11: $\Theta()$ for Bubble Sort

- $\bullet\,$  Running time  $\propto$  # of comparisons
- Number of Comparisons:

$$\sum_{i=1}^{n-1} i \in \Theta(n^2)$$

#### 10-12: Selection Sort

- Scan through the list, and find the smallest element
- Swap smallest element into position 0
- Scan through the list, and find the second smallest element
- Swap second smallest element into position 1
  ...
- Scan through the list, and find the second largest element
- Swap smallest largest into position n-2

### 10-13: $\Theta()$ for Selection Sort

- Running time  $\propto$  # of comparisons
- Number of Comparisons:

### 10-14: $\Theta()$ for Selection Sort

- $\bullet\,$  Running time  $\propto$  # of comparisons
- Number of Comparisons:

$$\sum_{i=1}^{n-1} i \in \Theta(n^2)$$

### 10-15: Improving Insertion Sort

- Insertion sort is fast if a list is "almost sorted"
- How can we use this?
  - Do some work to make the list "almost sorted"
  - Run insertion sort to finish sorting the list
- Only helps if work required to make list "almost sorted" is less than  $n^2$

- Sort n/2 sublists of length 2, using insertion sort
- Sort n/4 sublists of length 4, using insertion sort
- Sort n/8 sublists of length 8, using insertion sort  $\dots$
- Sort 2 sublists of length n/2, using insertion sort
- Sort 1 sublist of length n, using insertion sort

#### 10-17: Shell's Increments

- Shell sort runs several insertion sorts, using increments
  - Code on monitor uses "Shell's Increments":  $\{n/2, n/4, \dots 4, 2, 1\}$
- Problem with Shell's Increments:
  - Various sorts do not interact much
  - If all large elements are stored in large indices, and small elements are stored in even indices, what happens?

### 10-18: Other Increments

- Shell's Increments:  $\{n/2, n/4, \dots, 4, 2, 1\}$ • Running time:  $O(n^2)$
- "/3" increments:  $\{n/3, n/9, \dots, 9, 3, 1\}$ 
  - Running time:  $O(n^{\frac{3}{2}})$
- Hibbard's Increments:  $\{2^k 1, 2^{k-1} 1, \dots, 7, 3, 1\}$ 
  - Running time:  $O(n^{\frac{3}{2}})$

#### 10-19: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?

#### 10-20: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?

#### 10-21: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?
    - $\Theta(n)$
  - How Many Passes?

#### 10-22: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?
    - $\Theta(n)$
  - How Many Passes?
    - $\lg n$
  - Total running time?

#### 10-23: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?
    - $\Theta(n)$
  - How Many Passes?
    - $\lg n$
  - Total running time?
    - $\Theta(n \lg n)$

### 10-24: Stability

- Is Insertion sort stable?
- Is Bubble Sort stable?
- Is Selection Sort stable?
- Is Shell Sort stable?

#### 10-25: Stability

- Is Insertion sort stable? Yes!
- Is Bubble Sort stable? Yes!
- Is Selection Sort stable? No!
- Is Shell Sort stable? No!

Note that minor changes to the stable sorting algorithms will make them unstable (for instance, swaping A[i] and A[i + 1] when  $A[i] \ge A[i + 1]$ , not just when A[i] > A[i + 1]