Data Structures and Algorithms

CS245-2017S-14

Disjoint Sets

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14-0: Disjoint Sets

- Maintain a collection of sets
- Operations:
  - Determine which set an element is in
  - Union (merge) two sets
- Initially, each element is in its own set
  - # of sets = # of elements

14-1: Disjoint Sets

- Elements will be integers (for now)
- Operations:
  - CreateSets(n) – Create n sets, for integers 0..(n-1)
  - Union(x,y) – merge the set containing x and the set containing y
  - Find(x) – return a representation of x’s set
    - Find(x) = Find(y) iff x,y are in the same set

14-2: Disjoint Sets

- Implementing Disjoint sets
  - How should disjoint sets be implemented?

14-3: Implementing Disjoint Sets

- Implementing Disjoint sets (First Try)
  - Array of set identifiers:
    Set[i] = set containing element i
  - Initially, Set[i] = i

14-4: Implementing Disjoint Sets

- Creating sets:

14-5: Implementing Disjoint Sets

- Creating sets: (pseudo-Java)

```
void CreateSets(n) {
    for (i=0; i<n; i++) {
        Set[i] = i;
    }
}
```

14-6: Implementing Disjoint Sets

- Find:

14-7: Implementing Disjoint Sets

- Find: (pseudo-Java)

```
int Find(x) {
   return Set[x];
}
```

14-8: Implementing Disjoint Sets

- Union:

14-9: Implementing Disjoint Sets

- Union: (pseudo-Java)

```
void Union(x,y) {
   set1 = Set[x];
   set2 = Set[y];

   for (i=0; i < n; i=+)
     if (Set[i] == set2)
        Set[i] = set1;
}
```

14-10: **Disjoint Sets** $\Theta()$

- CreateSets
- Find
- Union

14-11: Disjoint Sets $\Theta()$

- CreateSets: $\Theta(n)$
- Find: $\Theta(1)$
- Union: $\Theta(n)$

14-12: Disjoint Sets $\Theta()$

- CreateSets: $\Theta(n)$
- Find: $\Theta(1)$
- Union: $\Theta(n)$

We can do better! (At least for Union ...)

14-13: Implementing Disjoint Sets II

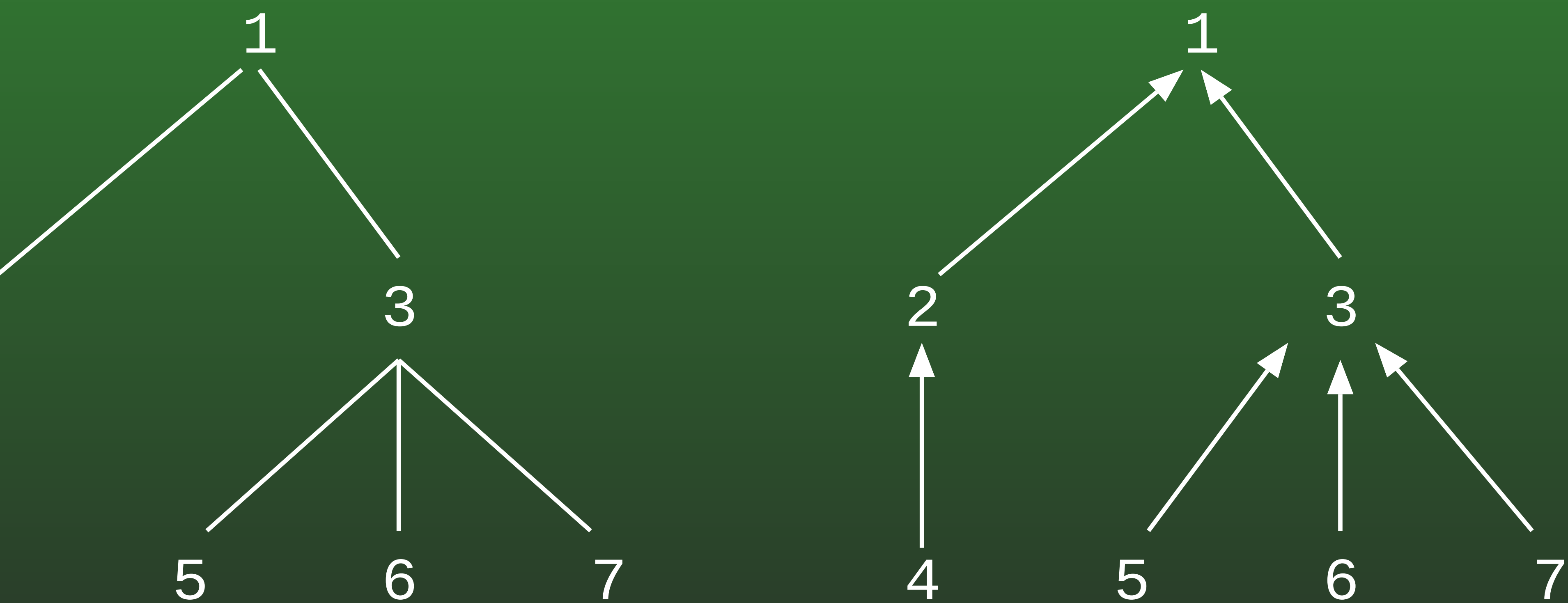
- Store elements in trees
- All elements in the same set will be in the same tree
- Find(x) returns the element at the root of the tree containing x
  - How can we easily find the root of a tree containing x?

14-14: Implementing Disjoint Sets II

- Store elements in trees
- All elements in the same set will be in the same tree
- Find(x) returns the element at the root of the tree containing x
  - How can we easily find the root of a tree containing x?
  - Implement trees using *parent pointers* instead of *children pointers*

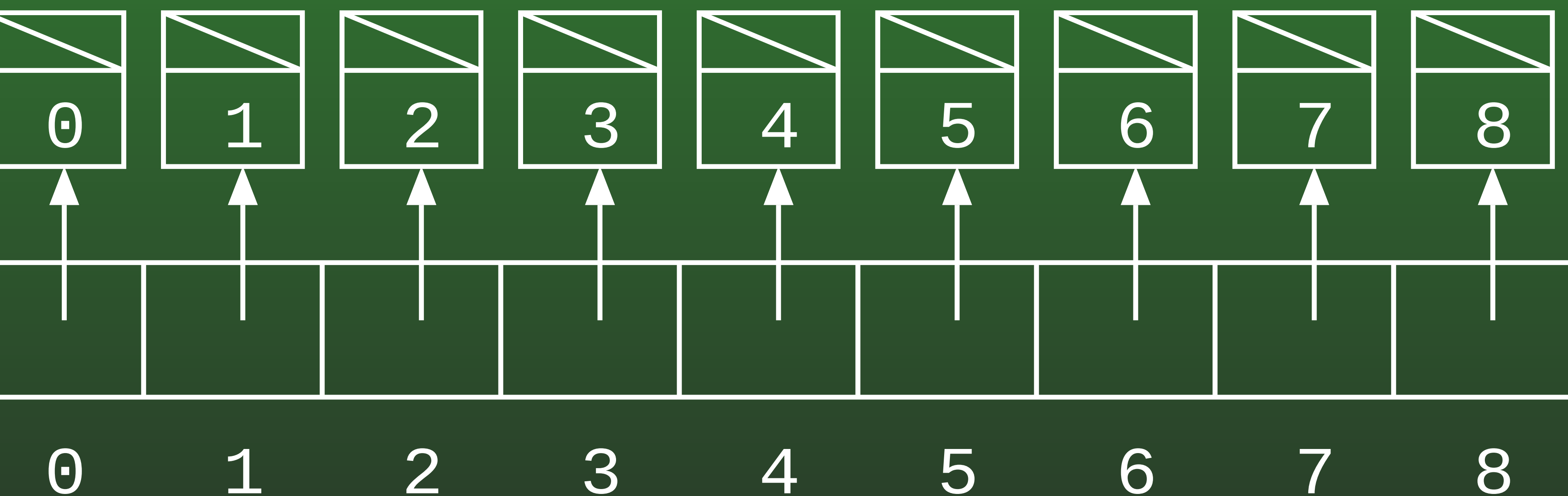
14-15: Trees Using Parent Pointers

- Examples:



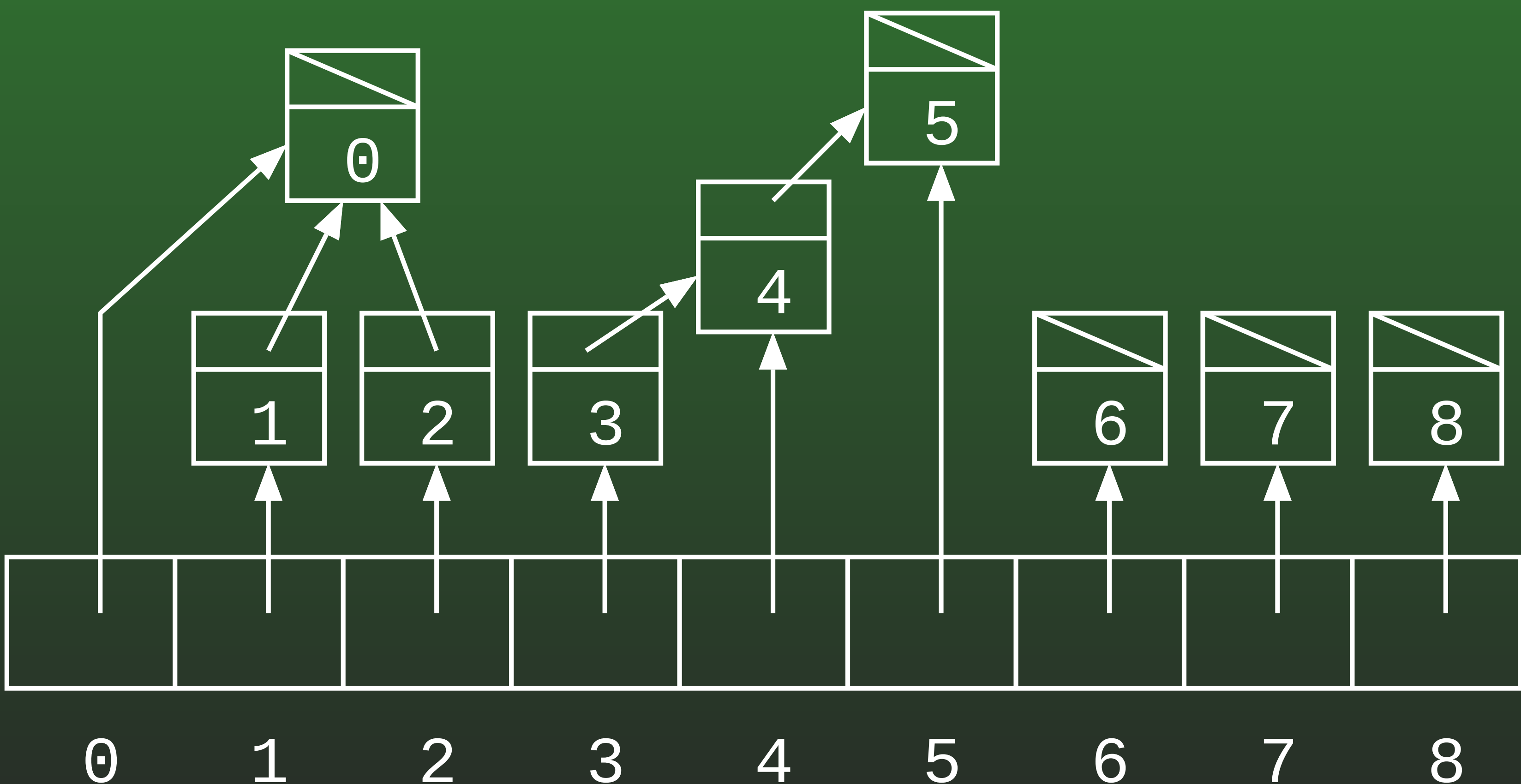
14-16: Implementing Disjoint Sets II

- Each element is represented by a node in a tree
- Maintain an array of pointers to nodes



14-17: Implementing Disjoint Sets II

- Each element is represented by a node in a tree
- Maintain an array of pointers to nodes



14-18: Implementing Disjoint Sets II

- Find:

14-19: Implementing Disjoint Sets II

- Find:
  - Follow parent pointers, until root is reached.
    - Root is node with `null` parent pointer.
  - Return element at root

14-20: Implementing Disjoint Sets II

- Find: (pseudo-Java)

```
int Find(x) {
   Node tmp = Sets[x];
   while (tmp.parent != null)
      tmp = tmp.parent;
   return tmp.element;
}
```

14-21: Implementing Disjoint Sets II

- Union(x,y)

Implementing Disjoint Sets II

- Union(x,y)
  - Calculate:
    - Root of x’s tree, rootx
    - Root of y’s tree, rooty
  - Set parent(rootx) = rooty

14-23: Implementing Disjoint Sets II

- Union(x,y) (pseudo-Java)

```
void Union(x,y) {
   rootx = Find(x);
   rooty = Find(y);
   Sets[rootx].parent = Sets[rooty];
}
```

Removing pointers

- We don’t need any pointers
- Instead, use index into set array

-1	-1	-1	-1	-1	-1	-1	-1	-1
0	1	2	3	4	5	6	7	8

Removing pointers

-1	-1	-1	-1	-1	-1	-1	-1	-1
0	1	2	3	4	5	6	7	8

- Union(2,3), Union(6,8), Union(0,2), Union(2,6)

Removing pointers

- Union(2,3), Union(6,8), Union(0,2), Union(2,8)

3	-1	3	8	-1	-1	8	-1	-1
0	1	2	3	4	5	6	7	8

14-27: Implementing Disjoint Sets III

- Find: (pseudo-Java)

```
int Find(x) {
   while (Parent[x] != -1)
      x = Parent[x]
   return x
}
```

14-28: Implementing Disjoint Sets II

- Union(x,y) (pseudo-Java)

```
void Union(x,y) {
   rootx = Find(x);
   rooty = Find(y);
   Parent[rootx] = rooty;
}
```

14-29: Efficiency of Disjoint Sets II

- So far, we haven't done much to improve the run-time efficiency of Disjoint sets.
- Two improvements will make a huge difference:
  - Union by rank
  - Path compression

14-30: Union by Rank

- When we merge two sets:
  - “Shorter” tree point to the taller tree

14-31: Union by Rank

- We need to keep track of the height of each tree
- How?

14-32: Union by Rank

- We need to keep track of the height of each tree
  - Store the height of the tree at the root
  - If a node x is not a root, $Parent[x]$ = parent of x
  - If a node x is a root, $Parent[x]$ = 0 - # height of tree rooted at x

14-33: Union by Rank

- Examples

14-34: Union by Rank

- When we merge two trees, how do we know which tree to point at the other?

14-35: Union by Rank

- When we merge two trees, how do we know which tree to point at the other?
  - The node with the larger (less negative) Parent[] value points to the node with the smaller (more negative) Parent[] value. Break ties arbitrarily.
- How do we update the height of the new merged tree?

Union by Rank

- When we merge two trees, how do we know which tree to point at the other?
    - The node with the larger (less negative) Parent[] value points to the node with the smaller (more negative) Parent[] value. Break ties arbitrarily.
- How do we update the height of the new merged tree?
    - If trees are different sizes, do nothing
    - If trees are the same size, increase (decrease) new parent by 1.

Union by Rank

- Union(x,y) (pseudo-Java)

```
void Union(x,y) {
   rootx = Find(x);
   rooty = Find(y);
   if (Parent[rootx] < Parent[rooty]) {
       Parent[rooty] = x;
   } else {
      if Parent[rootx] == Parent[rooty]
         Parent[rooty]--;
      Parent[rootx] = rooty;
   }
}
```

14-38: Path Compression

- After each call to `Find(x)`, change `x`'s parent pointer to point directly at root
- Also, change all parent pointers on path from `x` to root

14-39: Implementing Disjoint Sets III

- Find: (pseudo-Java)

```
int Find(x) {
   if (Parent[x] < 0)
      return x;
   else {
       Parent[x] = Find(Parent[x]);
       return Parent[x];
   }
}
```

14-40: Disjoint Set Θ

- Time to do a Find / Union proportional to the depth of the trees
- "Union by Rank" tends to keep tree sizes down
- "Path compression" makes Find and Union causes trees to flatten
- Union / Find take roughly time O(1) on average

14-41: Disjoint Set Θ

- Technically, n Find/Unions take time $O(n \lg^* n)$
- $\lg^* n$ is the number of times we need to take $\lg$ of n to get to 1.
  - $\lg 2 = 1$, $\lg^* 2 = 1$
  - $\lg(\lg 4) = 1$, $\lg^* 4 = 2$
  - $\lg(\lg(\lg 16)) = 1$, $\lg^* 16 = 3$
  - $\lg(\lg(\lg(\lg 65536))) = 1$, $\lg^* 65536 = 4$
  - $\ldots$
  - $\lg^* 2^{65536} = 5$
- # of atoms in the universe $\approx 10^{80} \ll 2^{65536}$
- $\lg^* n <= 5$ for all practical values of n