

CS 683

Computer Security and Privacy

EJ Jung

Announcements

- Step 1 of the lab is due on Tuesday, August 28
- Details at
 - <http://www.cs.usfca.edu/~ejung/courses/683/assignments.html>
- Readings
 - [TCP/IP Illustrated, Volume 1: The Protocols, Second Edition](#)
 - Chapter 2, Chapter 7

Extra credit opportunity

- Teach Internet Security at Tenderloin Technology Lab in St. Anthony's Foundation
- First-come first-served
- Replace 1 quiz of your choice (2~2.5%)
- Available dates:
 - One of September Wed 19th, Thur 20th, Mon 24th or Tues 25th, 10am-12noon
 - One of September Monday 17th or Tuesday the 18th, 2-4pm
 - More dates in October and November

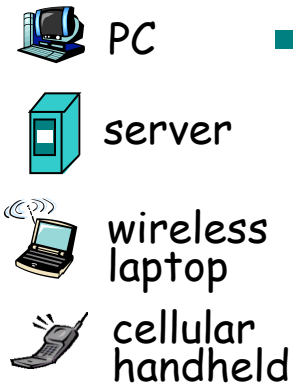
Academic Honesty

- Use any sources that you can learn from
- Do NOT use any sources that you copy from
 - cite if you used a reference
- DO NOT CHEAT
 - 1st time: 0 and report to the Dean's office
 - 2nd time: F

Today's goal: networking

- Half of you might know this already
 - 7 out of 14 have taken Computer Networks
 - fast forward or rewind as you see fit!
- TCP/IP model
 - Link layer/IP/TCP/application

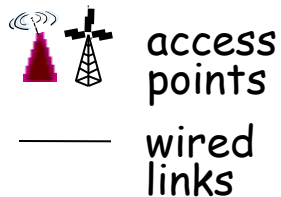
What's the Internet: "nuts and bolts" view



- millions of connected computing devices: *hosts = end systems*

- running *network apps*

- *communication links*

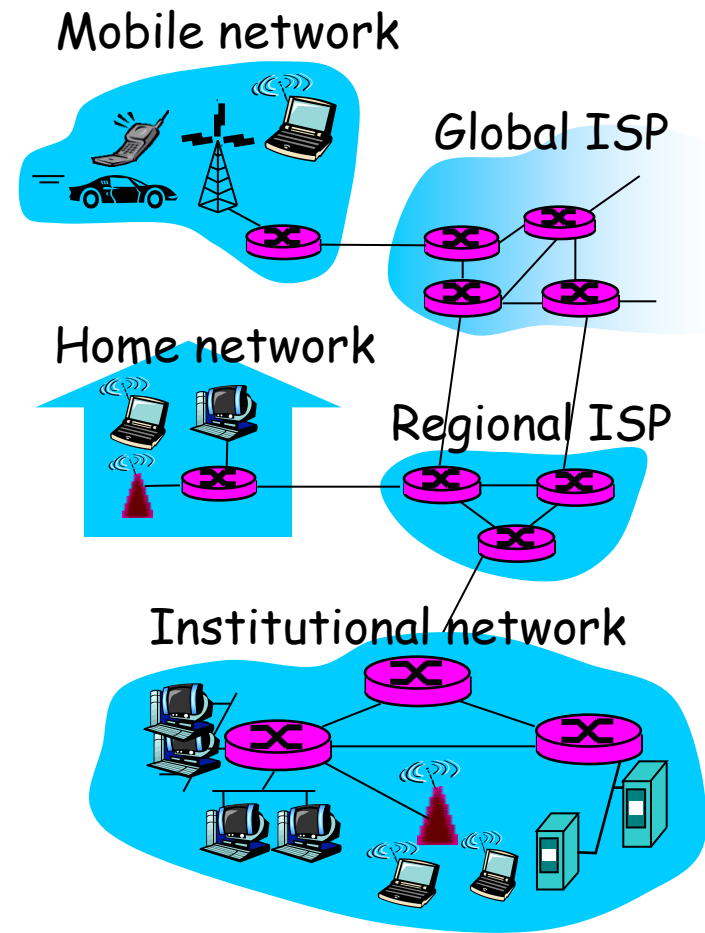


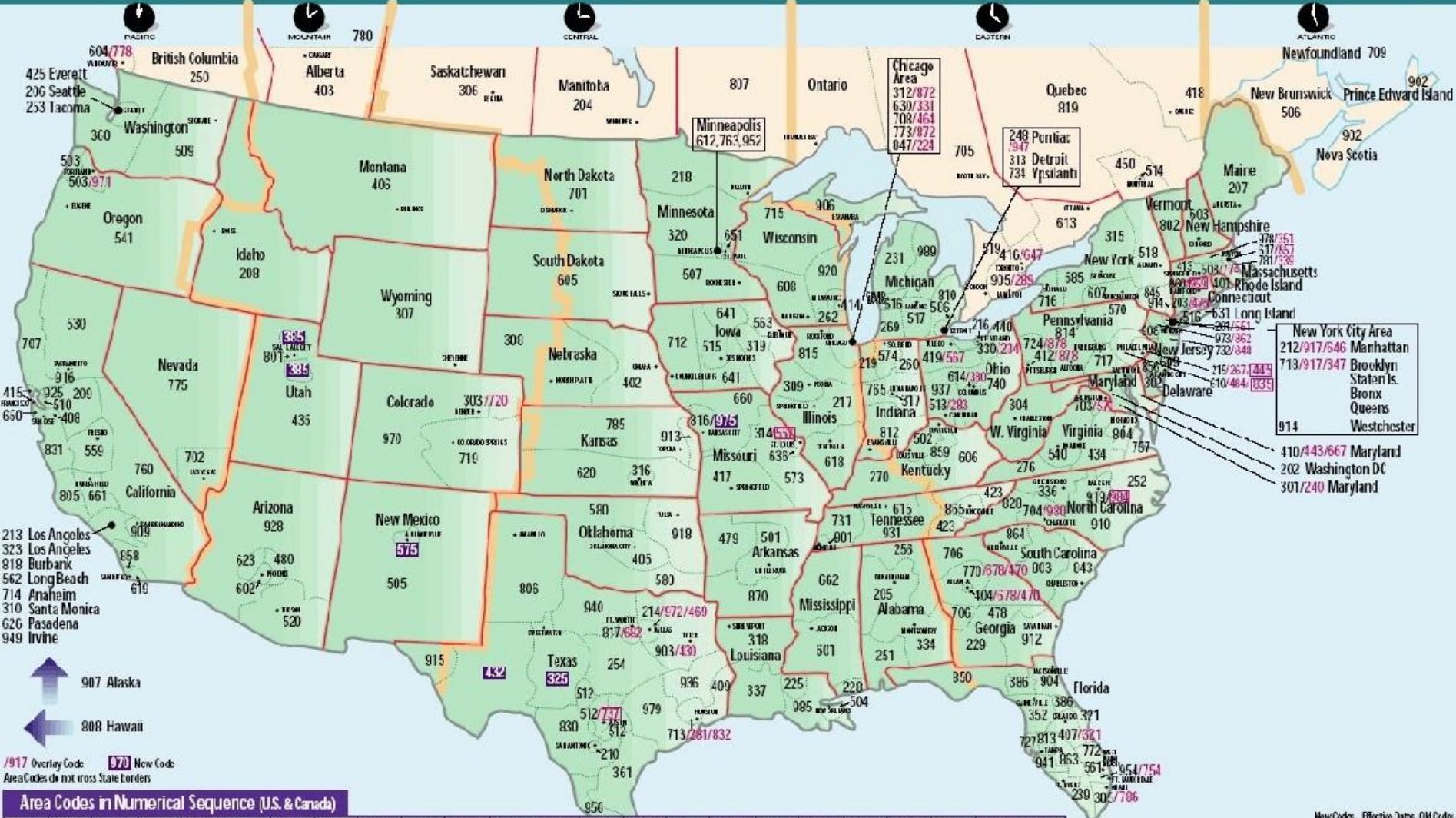
- fiber, copper, radio, satellite

- transmission rate = *bandwidth*



- *routers*: forward packets (chunks of data)

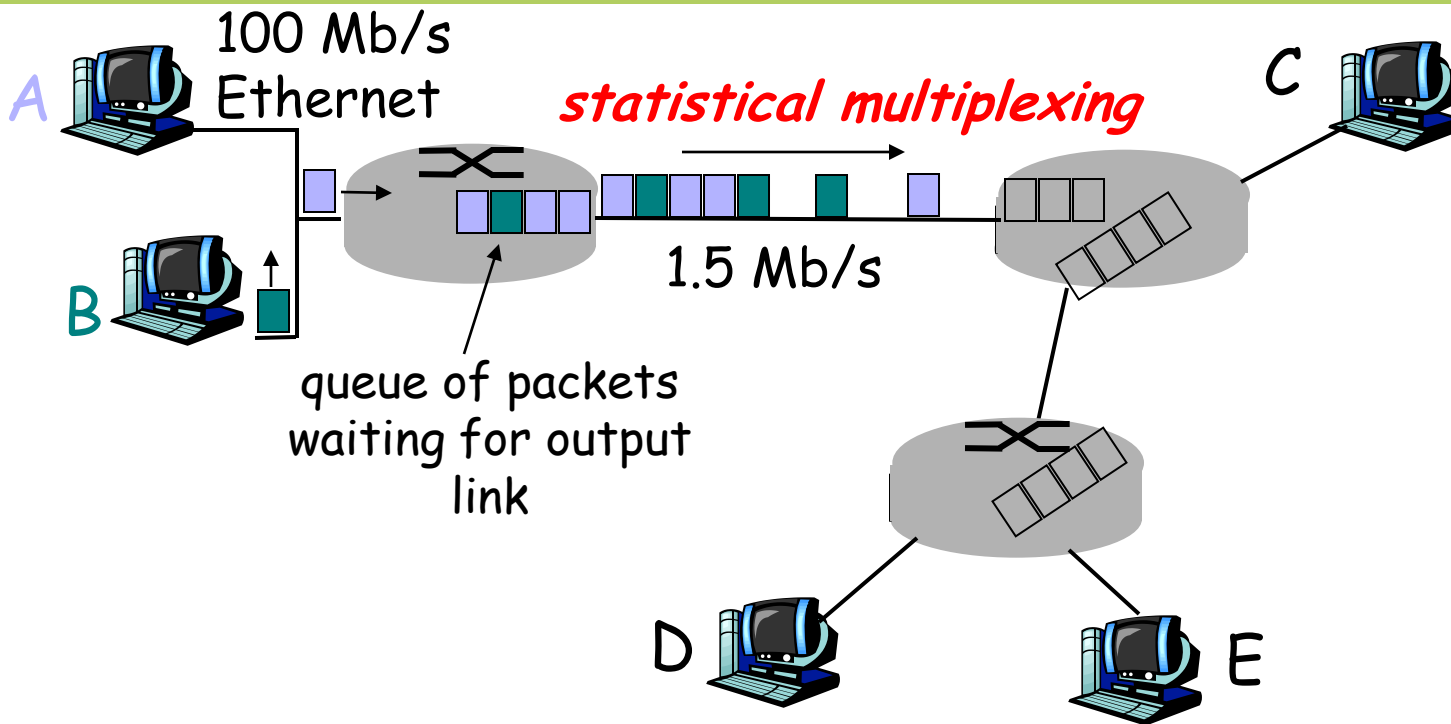




201 NJ	215 PA	240 MI	201 TX	312 IL	331 IL	401 RI	415 GA	445 PA	504 LA	518 NY	570 VA	606 KY	620 KS	662 MS	712 IA	734 MI	778 DC	800 II	832 TX	863 FL	907 AK	925 CA	959 CT
202 DC	216 OH	250 BC	283 OH	313 MI	334 AL	402 NE	410 ON	450 PO	505 NM	519 ON	571 VA	607 NY	623 AZ	667 MN	713 IL	737 TX	780 AR	810 MI	835 PA	864 SC	908 NJ	928 AZ	940 CO
203 CT	217 IL	251 AL	289 ON	314 MO	336 NC	403 AD	417 MO	464 IL	506 MD	520 AZ	573 MO	600 WI	626 CA	670 GA	714 CA	740 OH	781 MA	812 IN	843 SC	855 TN	909 CA	931 TN	971 OR
204 ME	218 MN	252 NC	301 MD	315 NY	337 LA	404 GA	413 PU	469 TX	507 MN	530 CA	574 IN	609 NJ	630 IL	682 TX	715 WI	754 FL	785 KS	813 FL	845 NY	867 MI	910 NC	936 TX	972 IL
205 AL	219 IN	253 WA	302 DC	316 KS	339 MA	405 OK	419 OH	470 GA	508 MA	540 VA	575 NM	610 PA	631 NY	701 ND	716 NY	757 VA	786 FL	814 PA	847 L	870 AR	912 CA	937 OH	973 NJ
206 WA	224 IL	254 TX	303 CO	317 IN	347 NY	406 MI	423 TX	475 CT	509 WA	541 OR	580 OK	612 MN	636 MO	702 NV	717 PA	763 CA	787 PR	815 IL	848 NJ	872 IL	913 KS	939 PR	975 MO
207 ME	225 LA	256 AL	304 WV	318 LA	351 MA	407 FL	425 WA	478 GA	510 CA	551 NJ	585 NY	613 ON	641 IA	703 VA	718 NY	763 MN	801 UT	816 MO	850 FL	878 PA	914 NY	940 TX	978 MA
208 ID	226 MS	250 IN	305 FL	319 IA	352 FL	408 CA	430 TX	479 AR	512 TX	557 MO	586 MI	614 OH	646 NY	704 NC	719 CO	765 IN	802 VT	817 TX	856 WI	901 TN	915 TX	941 F	970 TX
209 CA	229 GA	252 WI	306 SK	320 MN	360 WA	409 TX	432 TX	480 AZ	513 OH	559 CA	601 MS	615 TN	647 ON	705 ON	720 CO	770 GA	803 SC	818 CA	857 MA	902 NS	916 CA	947 MI	980 NC
210 CA	231 MI	257 PA	307 WV	321 FL	361 TX	410 MD	434 VA	484 PA	514 PO	561 FL	602 AZ	616 MI	650 CA	706 GA	724 CO	772 FL	804 VA	819 PQ	858 CA	903 TX	917 NY	949 CA	984 NC
212 NY	234 OH	259 MI	308 NE	323 CA	380 OH	412 PA	435 UT	501 AR	515 IA	562 CA	603 NH	617 MA	651 MN	707 CA	727 FL	773 IL	805 CA	828 NC	859 KY	904 FL	918 OK	952 MN	985 LA
213 CA	236 FL	270 KY	309 IL	325 TX	385 UT	413 MA	440 OH	502 KY	516 NY	563 IA	604 EC	618 IL	660 MO	708 IL	731 TN	774 MA	806 TX	830 TX	860 CT	905 ON	919 NC	954 FL	989 MI
214 TX	240 MD	276 VA	310 CA	330 OH	366 FL	414 WI	443 MD	503 OR	517 MI	567 OH	605 SD	619 CA	661 CA	709 NF	732 NJ	775 NY	807 ON	831 CA	862 NJ	906 MI	920 WI	956 TX	

New Codes			Effective Dates	Old Codes
Caribbean & Other Islands				
242	Bahamas	670	C.N.M.I.	
246	Barbados	671	Guam	
264	Anguilla	758	St. Lucia	
268	Antigua/Barbuda	767	Dominica	
284	B.V.I.	784	St. Vincent	
340	U.S.V.I.	787/938	Puerto Rico	
345	Cayman Islands	809	Dominican Republic	
441	Bermuda	868	Trinidad & Tobago	
473	Grenada	869	St. Kitts & Nevis	
649	Turks & Caicos	876	Jamaica	
664	Montserrat			
325	TX	4/5/03	915	
432	TX	4/5/03	915	
385	UT	3/30/03	801	
575	NM	6/1/03	505	
984	NC	TBD	919	
975	MO	TBD	816	
657	MO	TBD	314	
445	PA	TBD	215/267	
835	PA	TBD	610/484	
475	CT	TBD	203	
959	CT	TBD	850	
737	TX	TBD	512	

Packet Switching: Statistical Multiplexing

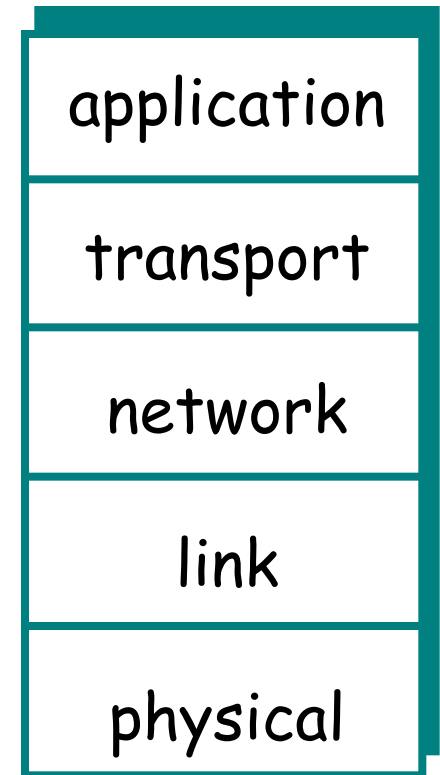


Sequence of A & B packets does not have fixed pattern,
bandwidth shared on demand → **statistical multiplexing**.

TDM: each host gets same slot in revolving TDM frame.

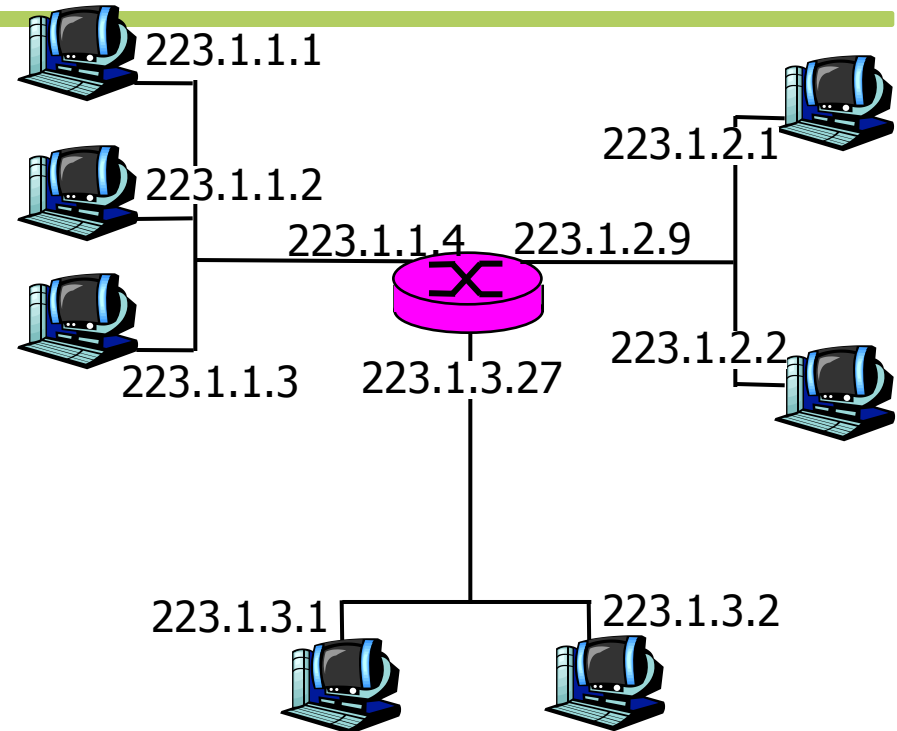
Internet protocol stack

- **application:** supporting network applications
 - FTP, SMTP, HTTP
- **transport:** process-process data transfer
 - TCP, UDP
- **network:** routing of datagrams from source to destination
 - IP, routing protocols
- **link:** data transfer between neighboring network elements
 - PPP, Ethernet
- **physical:** bits “on the wire”



IP Addressing: introduction

- IP address: 32-bit identifier for host, router *interface*
- *interface*: connection between host/router and physical link
 - router's typically have multiple interfaces
 - host typically has one interface
 - IP addresses associated with each interface



$$223.1.1.1 = \underbrace{11011111}_{223} \underbrace{00000001}_1 \underbrace{00000001}_1 \underbrace{00000001}_1$$

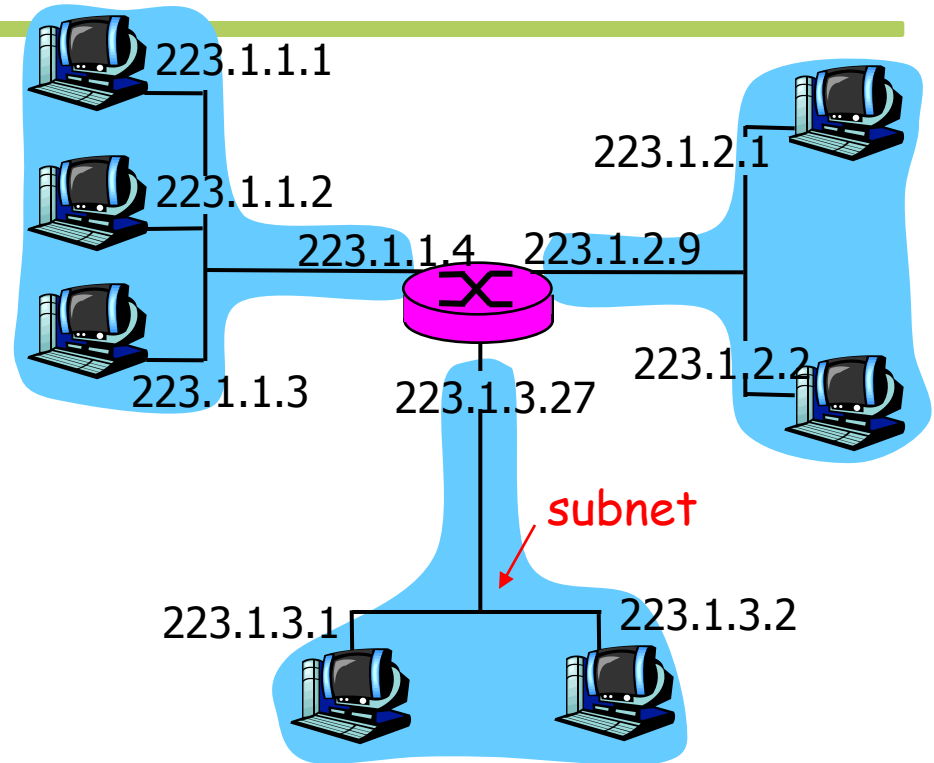
Subnets

➤ IP address:

- subnet part (high order bits)
- host part (low order bits)

➤ *What's a subnet ?*

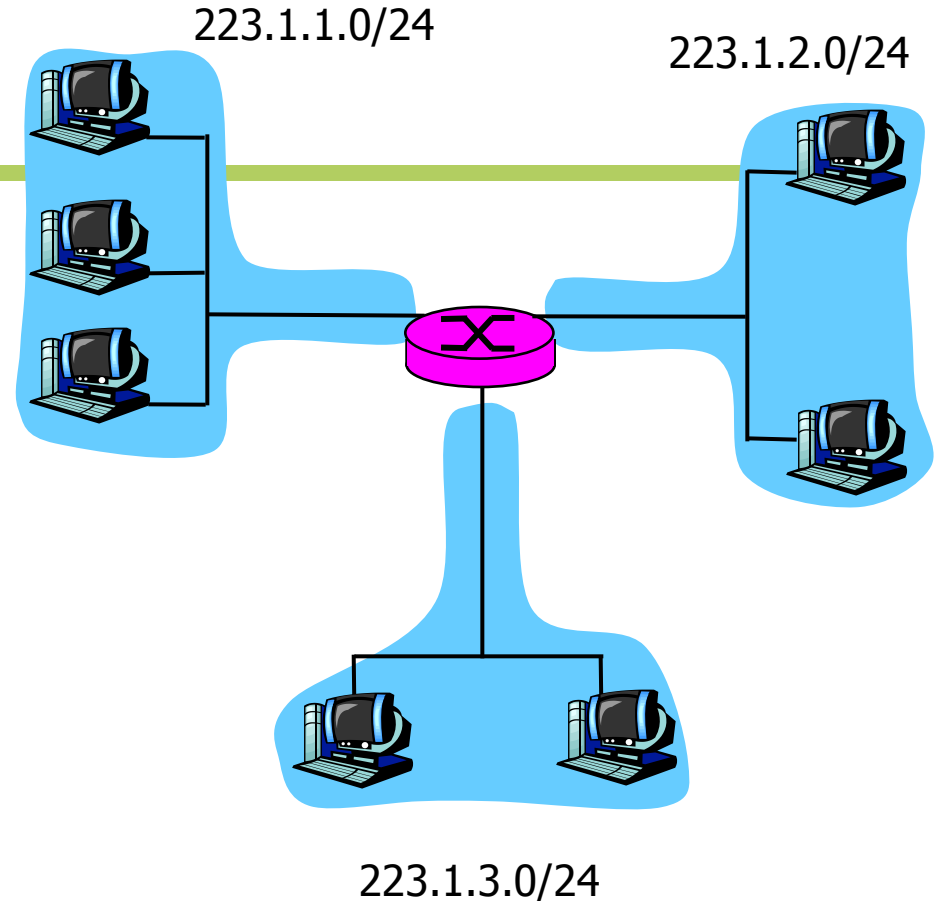
- device interfaces with same subnet part of IP address
- can physically reach each other without intervening router



network consisting of 3 subnets

Recipe

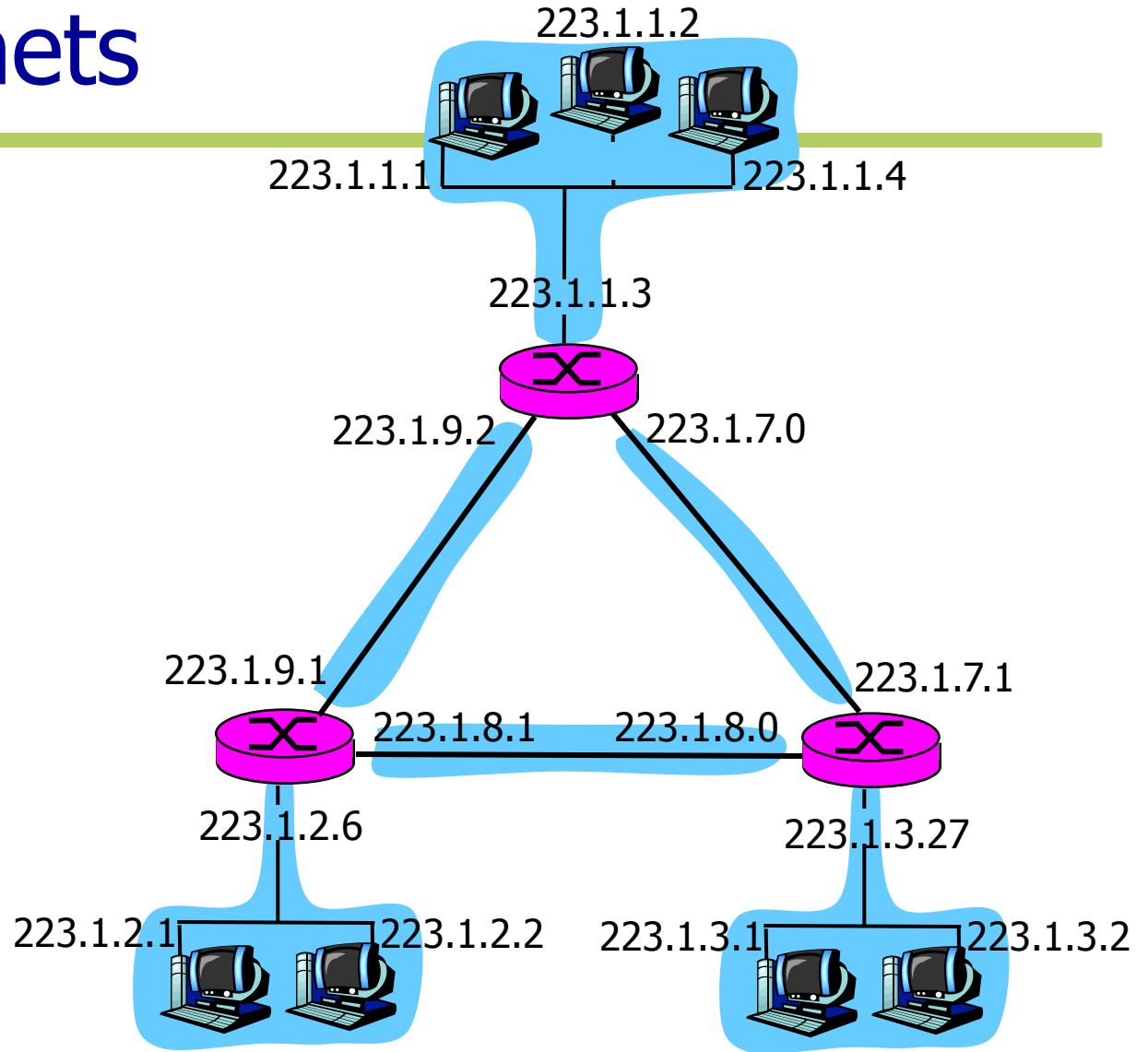
- To determine the subnets, detach each interface from its host or router, creating islands of isolated networks. Each isolated network is called a **subnet**.



Subnet mask: /24

Subnets

How many?



IP Addressing

given notion of “network”, let’s re-examine IP addresses:

“class-ful” addressing in the original IP design:

class

A	0	network		host		1.0.0.0 to 127.255.255.255
B	10		network		host	128.0.0.0 to 191.255.255.255
C	110		network		host	192.0.0.0 to 223.255.255.255
D	1110			multicast address		224.0.0.0 to 239.255.255.255

← 32 bits →

IP Addressing: CIDR

➤ (Static)classful addressing:

- inefficient use of address space, address space exhaustion
 - e.g., a class A net allocated enough addresses for 16 million hosts, even if only 2K hosts in that network
- not flexible for aggregation

➤ CIDR: Classless InterDomain Routing

- network portion of address of arbitrary length
- address format: a.b.c.d/x, where x is # bits in network portion of address

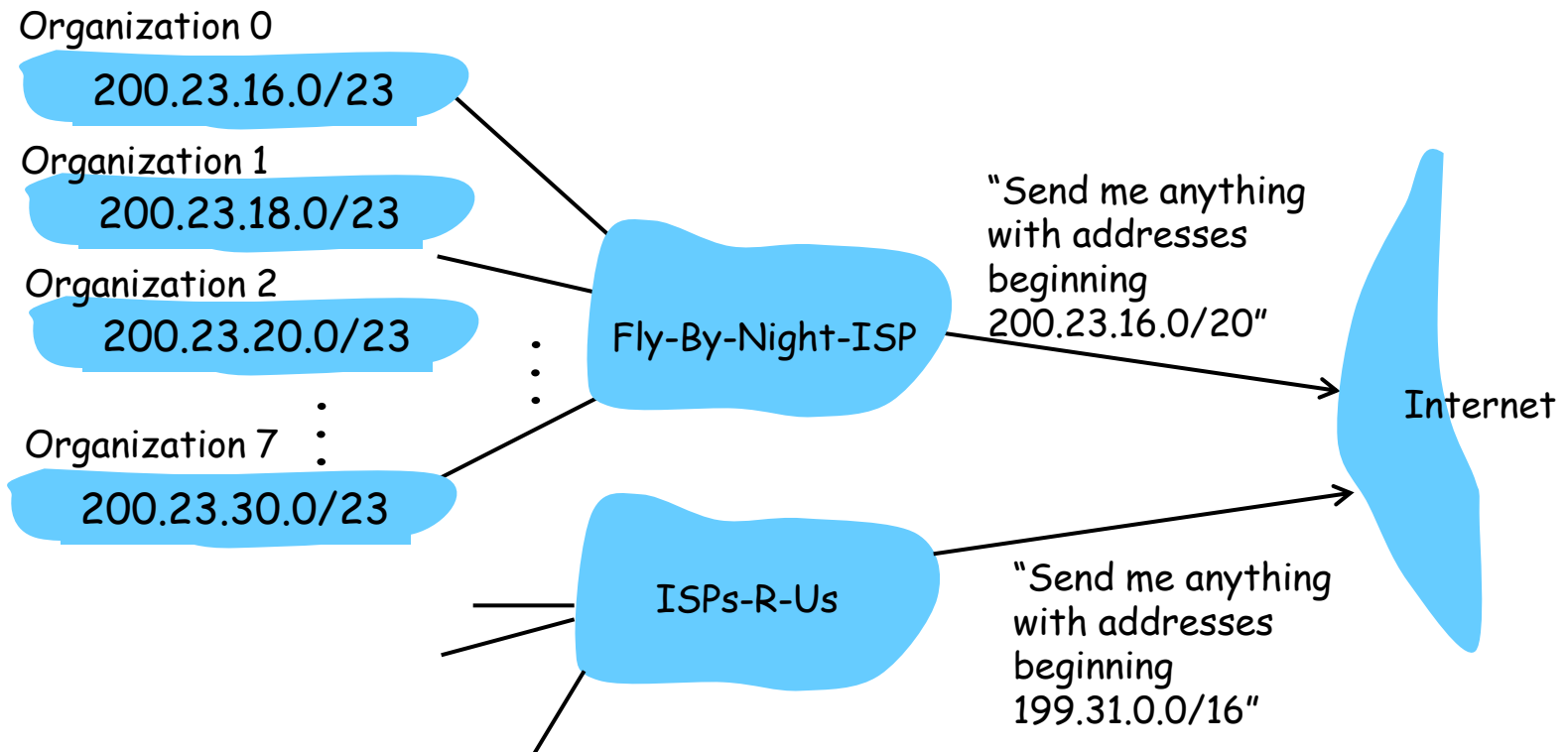


200.23.16.0/23

Some systems use mask, instead of the /x format

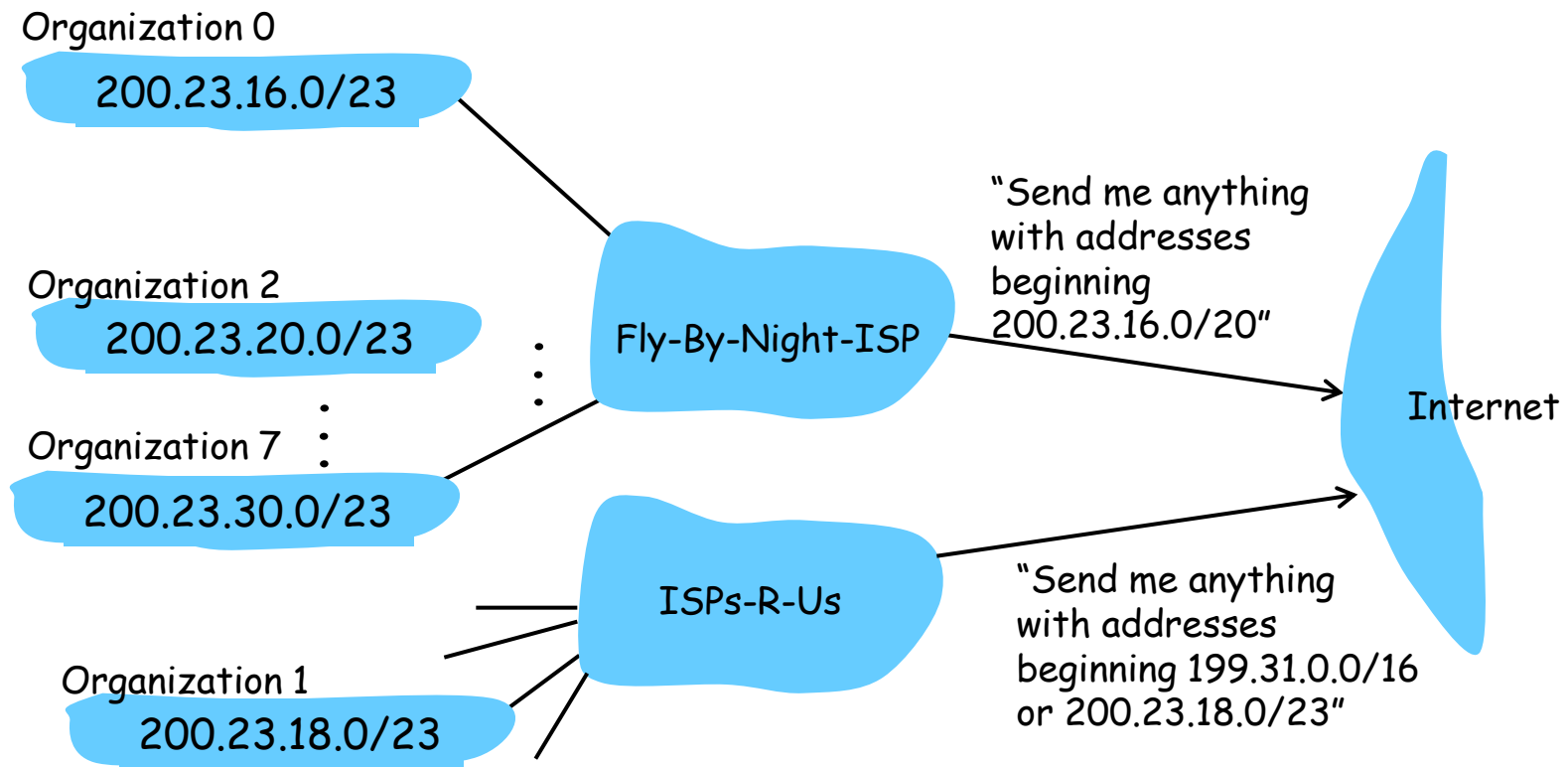
Hierarchical addressing: route aggregation

Hierarchical addressing allows efficient advertisement of routing information:

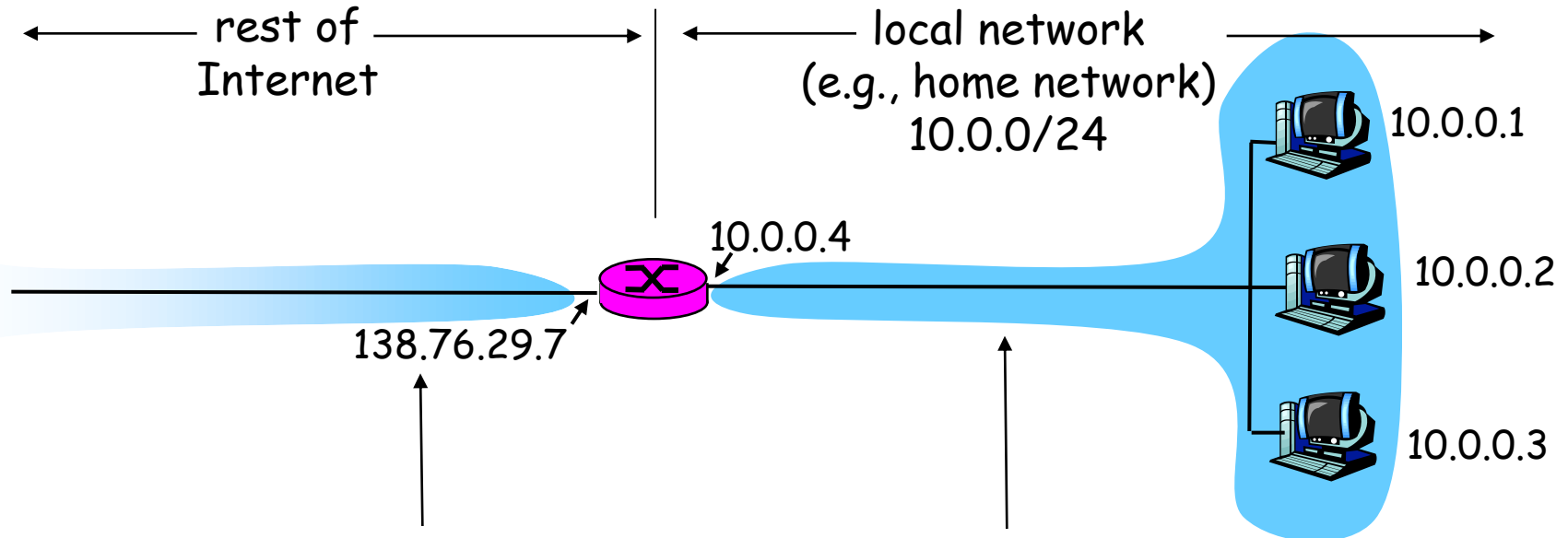


Hierarchical addressing: more specific routes

ISPs-R-Us has a more specific route to Organization 1



NAT: Network Address Translation



All datagrams *leaving* local network have *same* single source NAT IP address: 138.76.29.7, different source port numbers

Datagrams with source or destination in this network have 10.0.0/24 address for source, destination (as usual)

NAT: Network Address Translation

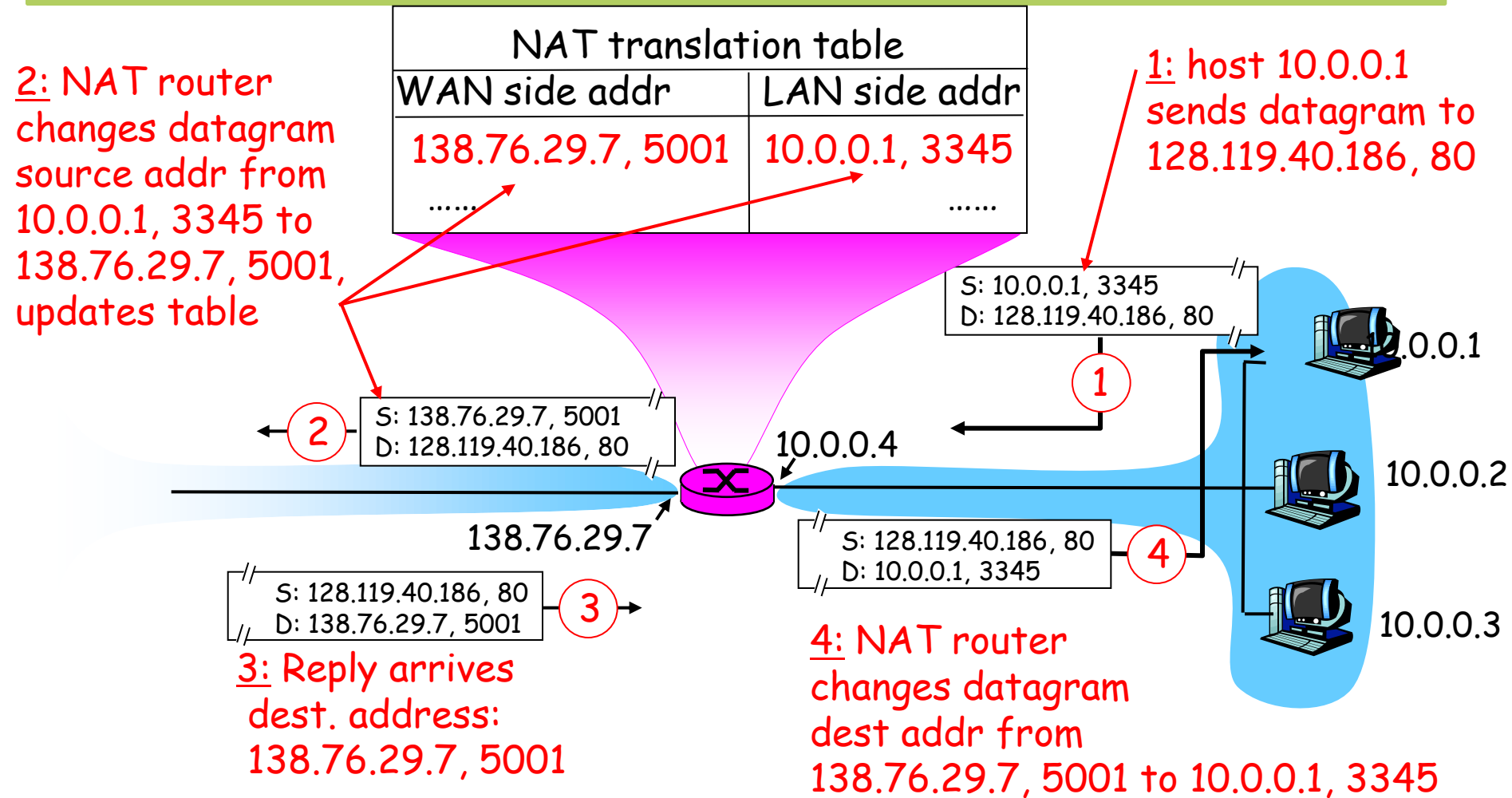
- **Motivation:** local network uses just one IP address as far as outside world is concerned:
- range of addresses not needed from ISP: just one IP address for all devices
 - can change addresses of devices in local network without notifying outside world
 - can change ISP without changing addresses of devices in local network
 - devices inside local net not explicitly addressable, visible by outside world (a security plus).

NAT: Network Address Translation

Implementation: NAT router must:

- *outgoing datagrams: replace* (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
 - . . . remote clients/servers will respond using (NAT IP address, new port #) as destination addr.
- *remember (in NAT translation table)* every (source IP address, port #) to (NAT IP address, new port #) translation pair
- *incoming datagrams: replace* (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table

NAT: Network Address Translation

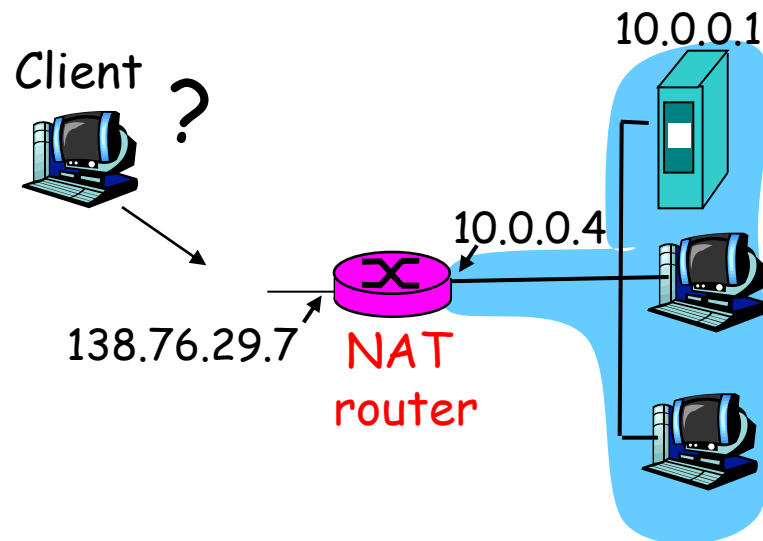


NAT: Network Address Translation

- 16-bit port-number field:
 - 60,000 simultaneous connections with a single LAN-side address!
- NAT is controversial:
 - routers should only process up to Network layer
 - violates end-to-end argument
 - NAT possibility must be taken into account by app designers, eg, P2P applications
 - address shortage should instead be solved by IPv6

NAT traversal problem

- client want to connect to server with address 10.0.0.1
 - server address 10.0.0.1 local to LAN (client can't use it as destination addr)
 - only one externally visible NATted address: 138.76.29.7
- solution 1: statically configure NAT to forward incoming connection requests at given port to server
 - e.g., (123.76.29.7, port 2500) always forwarded to 10.0.0.1 port 25000

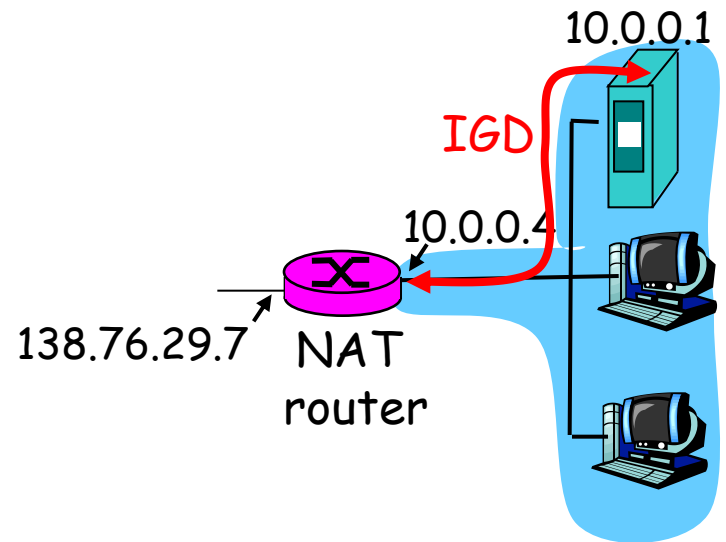


NAT traversal problem

- solution 2: Universal Plug and Play (UPnP) Internet Gateway Device (IGD) Protocol. Allows NATted host to:

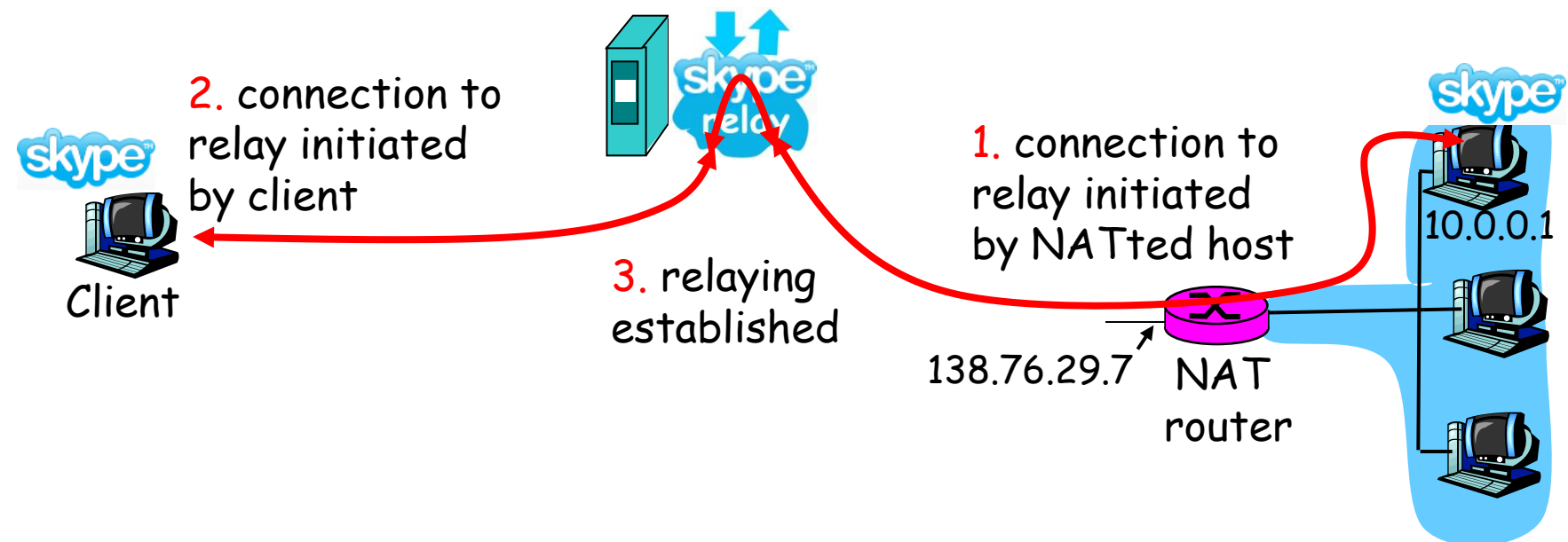
- ❖ learn public IP address (138.76.29.7)
- ❖ enumerate existing port mappings
- ❖ add/remove port mappings (with lease times)

i.e., automate static NAT port map configuration

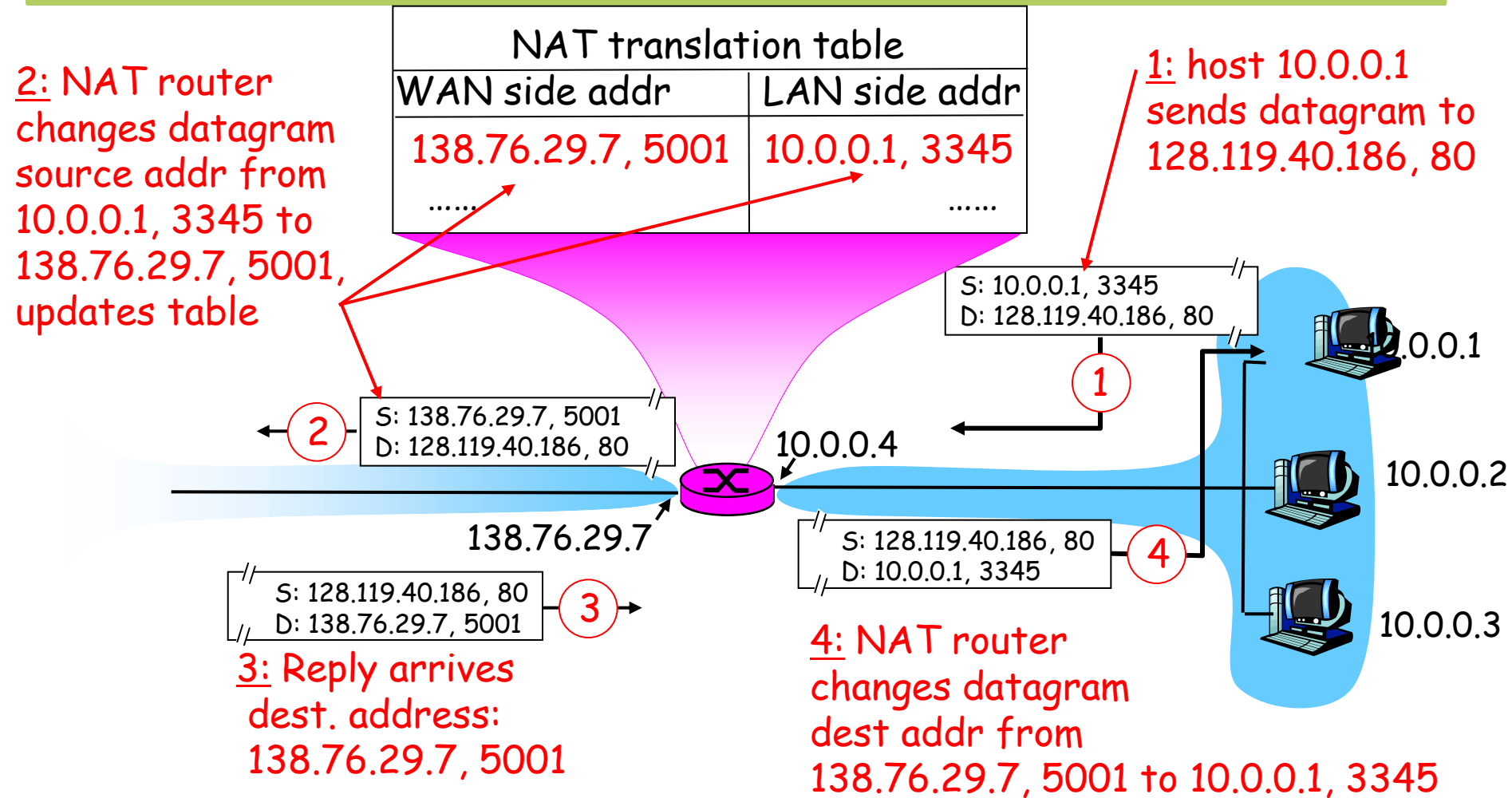


NAT traversal problem

- solution 3: relaying (used in Skype)
 - NATed server establishes connection to relay
 - External client connects to relay
 - relay bridges packets between to connections



NAT: Network Address Translation

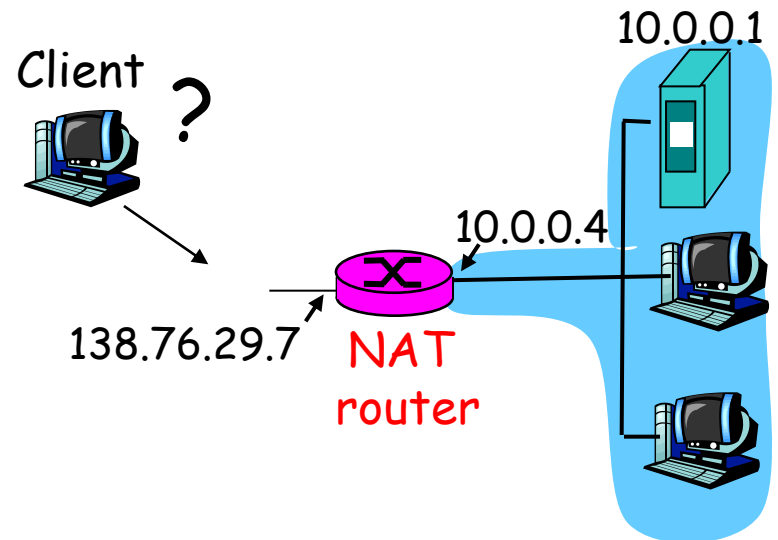


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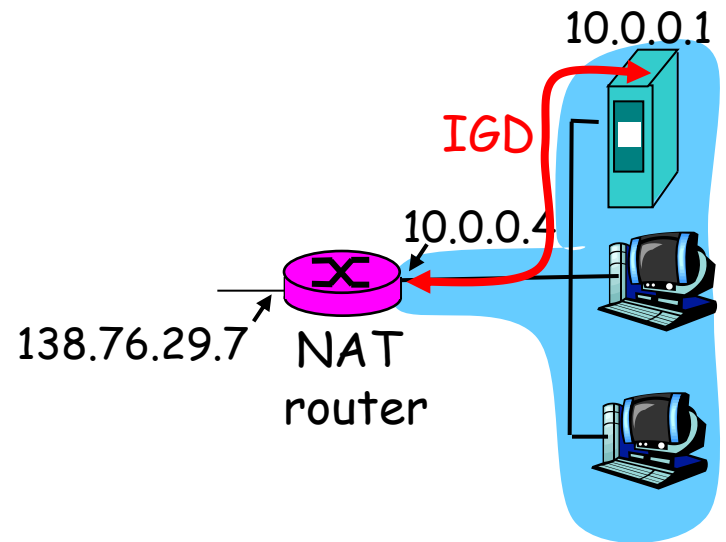


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