Introduction to Computer Science II

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Recursion I

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14-0: Recursion Review

- Factorial from visualization

14-1: Recursion – Power

- Write a method power
  `int power(int x, int n)`
  - Return x^n
- What is the base case?
- How can we make the problem smaller?
- How can we use the solution to the smaller problem to solve the original problem?

14-2: Recursion – Power

```
int power(int x, int n)
```

- What is the base case?
  - What is a version of the problem that is easy, that we can solve immediately?

14-3: Recursion – Power

`int power(int x, int n)`

- What is the base case?
  - Raising a number to the 1st power is easy
  - $x^1 == x$

```
int power(int x, int n)
{
   if (n == 1)
   {
      return x;
   }
}
```

14-4: Recursion – Power

```
int power(int x, int n)
```

- What is the recursive case?
  - How do we make the problem smaller?
  - How do we use the solution to the smaller problem to solve the original problem?

14-5: Recursion – Power

```
int power(int x, int n)
```

- What is the recursive case?
  - How do we make the problem smaller?
    - x^{n-1} is a smaller problem than x^n
  - How do we use the solution to the smaller problem to solve the original problem?
    - $x^n == x^{n-1} * x$

14-6: Recursion – Power

```
int power(int x, int n)
{
   if (n == 1)
   {
      return x;
   }
   else
   {
      return x * power(x, n - 1);
   }
}
```

- What about x^0 ?

14-7: Recursion – Power

```
int power(int x, int n)
{
   if (n == 0)
   {
      return 1;
   }
   else
   {
      return x * power(x, n - 1);
   }
}
```

14-8: Infinite Recursion

- What happens if we forget the base case?

```
int power(int x, int n)
{
   return x * power(x, n - 1);
}
```

14-9: Infinite Recursion

- What happens if we don’t make progress towards the base case?

```
int power(int x, int n)
{
   if (n == 0)
   {
      return 1;
   }
   else
   {
      return power(x, n);
   }
}
```

14-10: Recursion – Reversing Digits

- Function that takes as input an integer
- Writes out the digits in reverse order

```
void printReversed(int n)
{
  ...
}
```

14-11: Recursion – Reversing Digits

- What's a easy number to print reversed?

```
void printReversed(int n)
{
  ...
}
```

14-12: Recursion – Reversing Digits

- What's a easy number to print reversed?

```
void printReversed(int n)
{
   if (n < 10)
   {
      System.out.println(n);
   }
   ...
}
```

14-13: Recursion – Reversing Digits

- How can we make the problem smaller
  - We have to make the problem smaller such that a solution to the smaller problem helps us solve the original problem

14-14: Recursion – Reversing Digits

- How can we make the problem smaller
  - Remove the last digit (dividing by 10)
  - How can this help?

14-15: Recursion – Reversing Digits

```
void printReversed(int n)
{
   if (n < 10)
   {
      System.out.println(n);
   }
   else
   {
      System.out.print(n % 10);
      printReversed(n / 10);
   }
}
```

14-16: Tail Recursion

- A method is *tail recursive* if no more work needs to be done after the recursive call
- We return the value of the recursve call unchanged
- None of the functions that we have seen so far have been tail-recursive

14-17: Tail Recursion

- Typically, when a function is tail recursive, we have an extra parameter
  - Extra parameter builds up the solution to the problem
  - Each recursive call adds to the solution
  - Base case returns this solution
  - Solution is returned all the way to the end

14-18: Tail Recursion

```
int factorialTR(int n, int result)
{
   if (n == 0)
      return result;
   return factorialTR(n - 1, result * n);
}

int factorial(int n)
{
   return factorialTR(n, 1);
}
```

14-19: Tail Recursion

- Tail recursion is a little easier to see with reversing a string:
  - Start with an empty result
  - Remove first character from input, push it on to result
  - repeat until the input is empty

14-20: Tail Recursion

```
 public static String reverseTR(String s, String reversed)
 {
     if (s.length() == 0)
         return reversed;
     return reverseTR(s.substring(1), s.charAt(0) + reversed);
 }

public static String reverse(String s)
{
     return reverseTR(s, "");
}
```

14-21: Tail Recursion

- Why is tail recursion useful?
  - “Standard” recursive functions require an activation record on the stack for each recursive call
    - We need to do some work after the recursive call is done
    - We need the information stored on the stack
    - Examples: factorial / reversing

14-22: Tail Recursion

- Why is tail recursion useful?
  - Tail recursive functions don’t need to maintain the activation record after the function is called
    - Just return the value returned by the recursive call
    - We could reuse the same activation record
    - Could even change the recursive call to a loop (scheme)

14-23: Searching a String

- Anywhere you use a loop, you *could* use recursion instead (and vice-versa)
  - Though there are some problems that are *easier* to solve recursively, and some that ar *easier* to solve iteratively
- Create a recurive function countLetters, which takes as input a String s and a characeter c, and returns the number of times that c occurs in s.

```
int occurs(String s, char c)
```

14-24: Searching a String

```
public static int occurances(String s, char c)
{
   if (s.equals(""))
   {
      return 0;
   }
   else if (s.charAt(0) == c)
   {
      return 1 + occurances(s.substring(1), c);
   }
   else
   {
      return occurances(s.substring(1), c);
   }
}
```

14-25: Searching an array

- Previous string searching code called substring over and over
  - A little inefficient, creating lots of new strings
- For searching a list, we *really* don't want to make extra copies
- Instead, we will write a function that searches a *range* of indicies in a list, instead of an entire list

14-26: Searching an array

```
boolean search(int A[], int elem, int lowIndex,
```

- return true if elem is in the list, between lowIndex and highIndex (inclusive)
- First up: Iterative solution

14-27: Searching an array

```
public static boolean find(int A[], int elem, int lowIndex, int highIndex)
{
  for (int i = lowIndex; i <= highIndex; i++)
  {
    if (A[i] == elem)
      return true;
  }
  return false;
}
```

14-28: Searching an array

```
boolean search(int A[], int elem, int lowIndex,
```

- return true if elem is in the list, between lowIndex and highIndex (inclusive)
- Next up: Recursive solution.

14-29: Searching an array

```
public static boolean findR(int A[], int elem, int lowIndex, int highIndex)
{
  if (lowIndex > highIndex)
  {
    return false;
  }
  else if (A[lowIndex] == elem)
  {
    return true;
  }
  else
  {
    return findR(A, elem, lowIndex+1, highIndex);
  }
}
```

14-30: Binary Search

- We have a sorted list of integers
- Want to determine if a given integer is in the list
- We could do a linear search (start from beginning, search to the end)
- Is there a better way?

14-31: Binary Search

- If we are looking for an element in an empy list, return false
- If the element in the center of the list is what we are looking for, return true
- If the element in the center of the list is less that what we are looking for, discard the left half of the list, continue looking
- If the element in the center of the list is greater than what we are looking for, discard the right half of the list, continue looking

14-32: Binary Search

- As before, actually throwing away half the list (creating a new list half as large) is not efficient
- We can have our binary search take as input parameters the range in which we are searching
- `boolean search(int A[], int lowIndex, int highIndex)`

14-33: Binary Search

```
public static boolean search(int A[], int elem, int lowIndex, int highIndex)
{
    if (lowIndex > highIndex)
    {
        return false;
    }
    int midIndex = (lowIndex + highIndex) / 2;
    if (A[midIndex] == elem)
        return true;
    else if (A[midIndex] < elem)
        return search(A, elem, lowIndex, midIndex - 1);
    else
        return search(A, elem, midIndex+1, highIndex);
}
```

14-34: Binary Search

- Our implementation of Binary Search is tail recursive
- How could we modify binary search to return the $index$ of the element in the array, if it exists, or -1 if it does not?

14-35: Binary Search

```
public static int search(int A[], int elem, int lowIndex, int highIndex)
{
    if (lowIndex > highIndex)
    {
        return -1;
    }
    int midIndex = (lowIndex + highIndex) / 2;
    if (A[midIndex] == elem)
        return midIndex;
    else if (A[midIndex] < elem)
        return search(A, elem, lowIndex, midIndex - 1);
    else
        return search(A, elem, midIndex+1, highIndex);
}
```

14-36: Hands On

- Write a recursive version of toUpperCase
  - String toUpperCase(String input)
  - Easier: Use Character.toUpperCase
  - Harder: Use casting of char to int (and back again)
- write a version of occurances that does not create any extra strings (you may add parameters)