08-0: **Priority Queue ADT**

Operations

- Add an element / key pair
- Return (and remove) element with smallest key

Keys are “priorities”, with smaller keys having a “better” priority

08-1: **Priority Queue ADT**

Operations

- Add an element / key pair
- Return (and remove) element with smallest key

Implementation:

- Sorted Array
  - Add Element
  - Remove Smallest Key

08-2: **Priority Queue ADT**

Operations

- Add an element / key pair
- Return (and remove) element with smallest key

Implementation:

- Sorted Array
  - Add Element $O(n)$
  - Remove Smallest Key $O(1)$ (using circular array)

08-3: **Priority Queue ADT**

Operations

- Add an element / key pair
- Return (and remove) element with smallest key

Implementation:

- Binary Search Tree
  - Add Element
  - Remove Smallest Key

08-4: **Priority Queue ADT**

Operations

- Add an element / key pair
- Return (and remove) element with smallest key
Implementation:

- Binary Search Tree
  - Add Element \( O(lg\ n) \)
  - Remove Smallest Key \( O(lg\ n) \)

If the tree is balanced

08-5: **Priority Queue ADT**

Operations

- Add an element / key pair
- Return (and remove) element with smallest key

Implementation:

- Binary Search Tree
  - Add Element \( O(n) \)
  - Remove Smallest Key \( O(n) \)

Computer Scientists are Pessimists

(Murphy was right)

08-6: **Heap Definition**

- Complete Binary Tree
- Heap Property
  - For every subtree in a tree, each value in the subtree is \( \geq \) value stored at the root of the subtree

08-7: **Heap Examples**

```
         1
        / \  
       2   4
      /   /  
     7   3   14
    /   /     
   9   8   5
```

Valid Heap

08-8: **Heap Examples**
08-9: **Heap Insert**

- What is the only place we can insert an element in a heap, and maintain the complete binary tree property?

08-10: **Heap Insert**

- What is the only place we can insert an element in a heap, and maintain the complete binary tree property?
  - “End” of the tree – as a child of the shallowest leaf that is farthest to the left
  - Will the resulting tree still be a heap?

08-11: **Heap Insert**

- What is the only place we can insert an element in a heap, and maintain the complete binary tree property?
  - “End” of the tree – as a child of the shallowest leaf that is farthest to the left
- Inserting an element at the “end” of the heap may break the heap property
  - Swap the value up the tree (examples)

08-12: **Heap Insert**

- Running time for Insert?

08-13: **Heap Insert**

- Running time for Insert?
  - Place element at end of tree: $O(1)$ (We’ll see a clever way to find the “end” of the tree in a bit)
  - Swap element up the tree: $O(\text{height of tree})$ (Worst case, swap all the way up to the root)
    - Height of a Complete Binary Tree with $n$ nodes?

08-14: **Heap Insert**

- Running time for Insert?
  - Place element at end of tree: $O(1)$ (We’ll see a clever way to find the “end” of the tree in a bit
  - Swap element up the tree: $O(\text{height of tree})$ (Worst case, swap all the way up to the root)
• Height of a Complete Binary Tree with \( n \) nodes = \( \Theta(lg n) \)

• Total running time: \( \Theta(lg n) \) in the worst case

08-15: **Heap Remove Smallest**

• Finding the smallest element is easy – at the root of the tree

• Removing the Root of the heap is hard

• What element is easy to remove? How could this help us?

08-16: **Heap Remove Smallest**

• Finding the smallest element is easy – at the root of the tree

• Removing the Root of the heap is hard

• Removing the element at the “end” of the heap is easy
  
  • Copy last element of heap into root
  
  • Remove the last element
  
  • Problem?

08-17: **Heap Remove Smallest**

• Finding the smallest element is easy – at the root of the tree

• Removing the Root of the heap is hard

• Removing the element at the “end” of the heap is easy
  
  • Copy last element of heap into root
  
  • Remove the last element
  
  • May break the heap property

08-18: **Heap Remove Smallest**

• Finding the smallest element is easy – at the root of the tree

• Removing the Root of the heap is hard

• Removing the element at the “end” of the heap is easy
  
  • Copy last element of heap into root
  
  • Remove the last element
  
  • Push the root down, until heap property is satisfied

08-19: **Heap Remove Smallest**

• Running time for remove smallest?

08-20: **Heap Remove Smallest**

• Running time for remove smallest?
• Copy last element into root, remove last element: $O(1)$, given a $O(1)$ time method to find the last element
• Push the root down: $O(\text{height of the tree})$ (Worst case, push element all the way down)
  • As before, Complete Binary Tree with $n$ elements has height $\Theta(\lg n)$
• Total time: $\Theta(\lg n)$ in the worst case

08-21: Representing Heaps

• Represent heaps using pointers, much like BSTs
  • Need to add parent pointers for insert to work correctly
  • Need to maintain a pointer to the location to insert the next element (this could be hard to update & maintain)
  • Space needed to store pointers – 3 per node – could be greater than the space need to store the data in the heap!
  • Memory allocation and deallocation is slow
• There is a better way!

08-22: Representing Heaps

A Complete Binary Tree can be stored in an array:

```
1 2 14 5 3 16 15 7 6 8 9
```

08-23: CBTs as Arrays

• The root is stored at index 1
• For the node stored at index $i$:
  • Left child is stored at index $2 \times i$
  • Right child is stored at index $2 \times i + 1$
  • Parent is stored at index $\lfloor i / 2 \rfloor$

08-24: CBTs as Arrays

Finding the parent of a node

```java
int parent(int n) {
    return (n / 2);
}
```
Finding the left child of a node

```c
int leftchild(int n) {
    return 2 * n;
}
```

Finding the right child of a node

```c
int rightchild(int n) {
    return 2 * n + 1;
}
```

08-25: **Building a Heap**
Build a heap out of \(n\) elements

08-26: **Building a Heap**
Build a heap out of \(n\) elements
- Start with an empty heap
- Do \(n\) insertions into the heap

```java
MinHeap H = new MinHeap();
for(i=0 < i<A.size(); i++)
    H.insert(A[i]);
```

Running time?

08-27: **Building a Heap**
Build a heap out of \(n\) elements
- Start with an empty heap
- Do \(n\) insertions into the heap

```java
MinHeap H = new MinHeap();
for(i=0 < i<A.size(); i++)
    H.insert(A[i]);
```

Running time? \(O(n \log n)\) – is this bound tight?

08-28: **Building a Heap**
Total time: \(c_1 + \sum_{i=1}^{n} c_2 \log i\)

\[
c_1 + \sum_{i=1}^{n} c_2 \log i \geq \sum_{i=n/2}^{n} c_2 \log i \\
\geq \sum_{i=n/2}^{n} c_2 \log(n/2) \\
= (n/2)c_2 \log(n/2) \\
= (n/2)c_2((\log n) - 1) \\
\in \Omega(n \log n)
\]

Running Time: \(\Theta(n \log n)\)

08-29: **Building a Heap**
Build a heap from the bottom up
• Place elements into a heap array
• Each leaf is a legal heap
• First potential problem is at location \( \lfloor i/2 \rfloor \)

08-30: **Building a Heap**
Build a heap from the bottom up
• Place elements into a heap array
• Each leaf is a legal heap
• First potential problem is at location \( \lfloor i/2 \rfloor \)

\[
\text{for}\ (i=n/2; \ i>=0; \ i--) \\
\text{pushdown}(i); \\
\]

08-31: **Building a Heap**
How many swaps, worst case? If every \text{pushdown} has to swap all the way to a leaf:
\[
\begin{align*}
n/4 \text{ elements} & \quad 1 \text{ swap} \\
n/8 \text{ elements} & \quad 2 \text{ swaps} \\
n/16 \text{ elements} & \quad 3 \text{ swaps} \\
n/32 \text{ elements} & \quad 4 \text{ swaps} \\
\ldots \\
\text{Total # of swaps:} & \quad n/4 + 2n/8 + 3n/16 + 4n/32 + \ldots + (\log n)n/n \\
\]

08-32: **Decreasing a Key**
• Given a specific element in a heap, how can we decrease the key of that element, and maintain the heap property?
  • Examples

08-33: **Decreasing a Key**
• Given a specific element in a heap, how can we decrease the key of that element, and maintain the heap property?
  • Examples
• Push the element up the tree, just like after an insert
  • Examples

08-34: **Decreasing a Key**
• Decrease the key of a specific element in a heap:
  • Decrease the key value
  • Push the element up the tree, just like after an insert
• Time required?

08-35: **Decreasing a Key**
• Decrease the key of a specific element in a heap:
• Decrease the key value
• Push the element up the tree, just like after an insert
• Time required: $\Theta(lg \ n)$, in the worst case.

08-36: **Removing an Element**

• Given a specific element in a heap, how can we remove that element, and maintain the heap property?
  • Examples

08-37: **Removing an Element**

• Given a specific element in a heap, how can we remove that element, and maintain the heap property?
  • Examples
  • Decrease key to a value $< \text{root}$
  • Remove smallest element

08-38: **Removing an Element**

• Given a specific element in a heap, how can we remove that element, and maintain the heap property?
  • Examples
  • Decrease key to a value $< \text{root}$. Time $\Theta(lg \ n)$ worst case
  • Remove smallest element. Time $\Theta(lg \ n)$ worst case

08-39: **Java Specifics**

• When inserting an element, push value up until it reaches the root, or it’s $\geq$ its parent
  • Our while statement will have two tests
• We can insert a *sentinel* value at index 0, guaranteed to be $\leq$ any element in the heap
  • Now our while loop only requires a single test