07-0: **Binary Search Trees**

- Binary Trees
  - For each node \( n \), (value stored at node \( n \)) > (value stored in left subtree)
  - For each node \( n \), (value stored at node \( n \)) < (value stored in right subtree)

07-1: **Example Binary Search Trees**

```
          A
         / \  
        B   D
       / \  /  
      D   E F
```

07-2: **Example Binary Search Trees**

- Examples:
  - Finding an element
  - Inserting an element
  - Deleting an element

07-3: **Running Times**

- Best-Case upper limit on the time for insert/delete/find of an element for a BST with \( n \) elements?
  - \( O(\lg n) \), if the tree is balanced
- Worst-Case upper limit on the time for insert/delete/find for a BST with \( n \) elements?
  - \( O(n) \), if the tree is a list
- Expected upper limit on the time for insert/delete/find for a BST with \( n \) elements?
  - What would we mean by “expected” in this instance?

07-4: **Running Times**

- Best-Case upper limit on the time for insert/delete/find of an element for a BST with \( n \) elements?
  - \( O(\lg n) \), if the tree is balanced
- Worst-Case upper limit on the time for insert/delete/find for a BST with \( n \) elements?
  - \( O(n) \), if the tree is a list
- Expected upper limit on the time for insert/delete/find for a BST with \( n \) elements?
- $O(lg n)$, if elements are inserted in random order

07-5: Balanced BSTs
- We can guarantee $O(lg n)$ running time for insert/find/delete if we can guarantee the tree is balanced
- Several methods for guaranteeing a balanced tree
  - AVL trees & Red-Black trees are the most common
  - We’ll look at Red-Black Trees

07-6: Red-Black Trees
- Red-Black Trees as Binary Search trees, with “Null Leaves”
  - Examples of BSTs with “Null Leaves”
  - (Null leaves are mostly a notational convenience)

07-7: Red-Black Trees
- Red-Black Trees are Binary Search trees, with “Null Leaves”, and the following properties:
  - Every Node is either Red or Black
  - (Root is Black) <Not strictly required>
  - Each null “leaf” is Black
  - If a node is red, both children are black
  - For each node, all paths from the node to descendant leaves contain the same number of black nodes

07-8: Red-Black Trees

07-9: Red-Black Trees
07-10: Red-Black Trees

- In a Red-Black tree, what is the greatest possible difference in the length of the path from the root to two different leaves?
- What is the largest height of a Red-Black tree that contains \( n \) elements?

07-11: Red-Black Trees

- Let \( bh(X) \) be the "Black Height" of a node – the number of black nodes on a path from that node to a leaf (not including the node itself)
- The subtree rooted at any node \( X \) has at least \( 2^{bh(X)} - 1 \) internal (non-leaf) nodes
  - Proof by induction (on board)

07-12: Tree Rotations

- Example Red-Black tree ("Null Leaves" left out for clarity)

07-13: Tree Insertions

- Always insert red nodes
- Which property would be violated by inserting a red node?

07-14: Tree Insertions

- Always insert red nodes
- Which property would be violated by inserting a red node?
  - Could have a red node with a red child
- Fix using tree rotations

07-15: Tree Insertions

- To fix a red node with red child:
  - Case 1: Uncle is red
  - Case 2: Uncle is black, Inserted node is right child of parent, and parent is a left child of Grandparent (or node is left child, parent is right child)
  - Case 3: Uncle is black, Node is left child of parent, parent is left child of Grandparent (or node is right child, parent is right child)

07-16: Case 1
- Red Uncle

07-17: Case 1
- Red Uncle

07-18: Case 2
- Black Uncle / parent child different handedness

07-19: Case 3
- Black Uncle / parent child same handedness
07-20: **Case 3**

- Black Uncle / parent child same handedness

07-21: **Deleting nodes**

- Deleting nodes
  - Delete nodes just like in standard BST
  - Which properties could be violated by deleting a red node?
    - Each node red or black
    - Black Root
    - Each red node has 2 black children
    - Black path length to leaves same for each node

07-22: **Deleting nodes**

- Deleting nodes
  - Delete nodes just like in standard BST
  - Which properties could be violated by deleting a red node?
    - None!

07-23: **Deleting Nodes**

- Deleting nodes
• Delete nodes just like in standard BST
• Which properties could be violated by deleting a black node?
  • Each node red or black
  • Black Root
  • Each red node has 2 black children
  • Black path length to leaves same for each node

07-24: Deleting Nodes

• Deleting black node
  • If the child of the deleted node is red ... (show example on board)

07-25: Deleting Nodes

• Deleting black node
  • If the child of the deleted node is black
    • Make the child “doubly black”
    • Push “extra blackness” up the tree until it can be removed by a rotation

07-26: Deleting Nodes

• X is “doubly black” node, X is a left child
  • Case 4:
    • X’s sibling W is black, and W’s right child is red
    • Can remove “double-blackness” of X with a single rotation

07-27: Deleting Nodes

07-28: Deleting Nodes
07-29: Deleting Nodes

Doubly Black

07-30: Deleting Nodes
07-31: Deleting Nodes

- X is “doubly black” node, X is a left child

  - Case 3:
    - X’s sibling W is black, and W’s left child is red, and right child is black
    - Single rotation to get to previous case

07-32: Deleting Nodes

Doubly Black

07-33: Deleting Nodes
07-34: Deleting Nodes

Doubly Black

A

\(\alpha\) \(\beta\)

B

C

\(\gamma\)

D

\(\delta\)

E

\(\epsilon\) \(\zeta\)

07-35: Deleting Nodes

Doubly Black

A

\(\alpha\) \(\beta\)

B

C

\(\gamma\) \(\delta\)

D

\(\epsilon\)

E

\(\zeta\)
07-36: Deleting Nodes

- X is “doubly black” node, X is a left child
  - Case 2:
    - X’s sibling W is black, and both of W’s children are black
    - Push “Blackness” of X and W to parent

07-37: Deleting Nodes

07-38: Deleting Nodes
07-39: Deleting Nodes

Doubly Black

07-40: Deleting Nodes
07-41: Deleting Nodes

- X is “doubly black” node, X is a left child
  - Case 2:
    - X’s sibling W is Red
    - Do a rotation, to make W black. Then one of the other cases will apply.

07-42: Deleting Nodes

07-43: Deleting Nodes
07-44: **Deleting Nodes**

- Need to include symmetric cases
  - In all of the previous examples, swap left/right
  - (Go over at least one example)