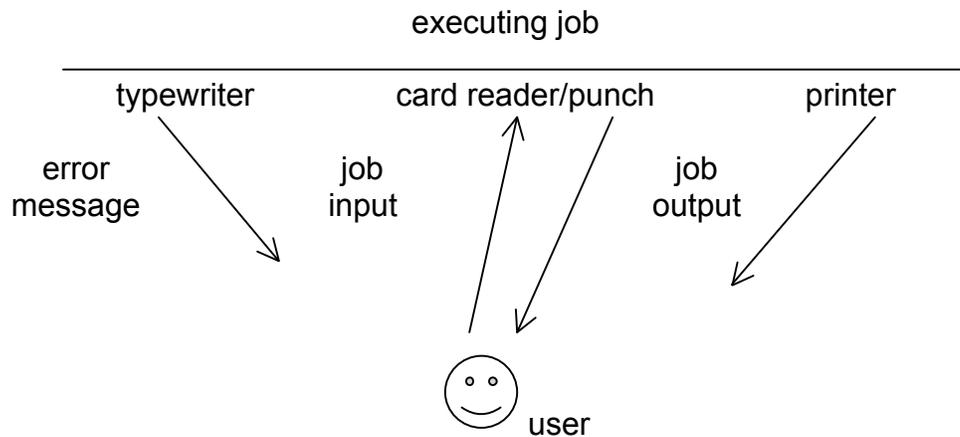


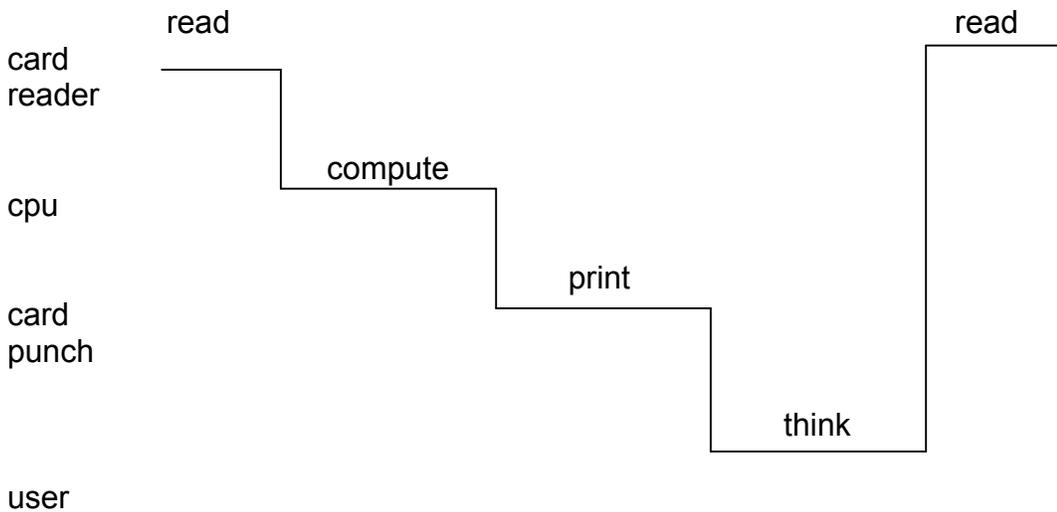
### Open Shop Model

Early computers, e.g., IBM 1620

Single user, hands-on, no operator, one job at a time, no operating system

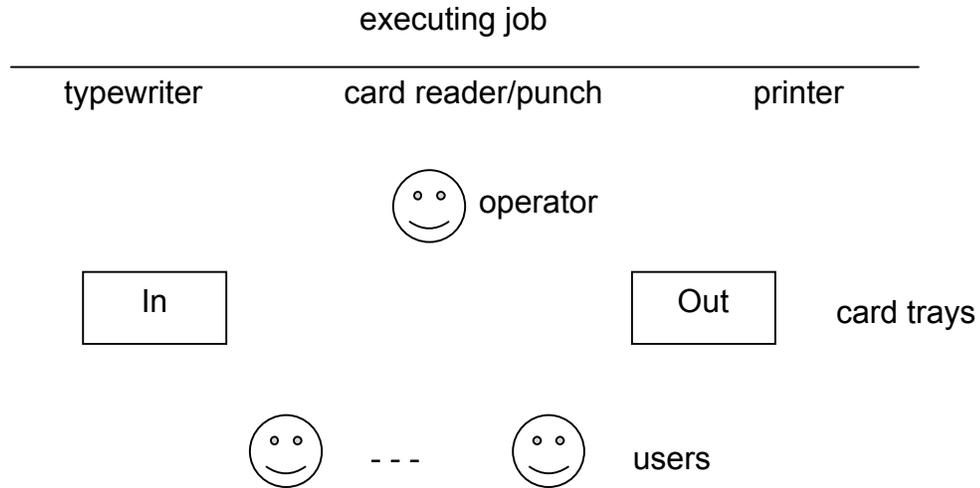


Utilization: Fraction of time used for computation

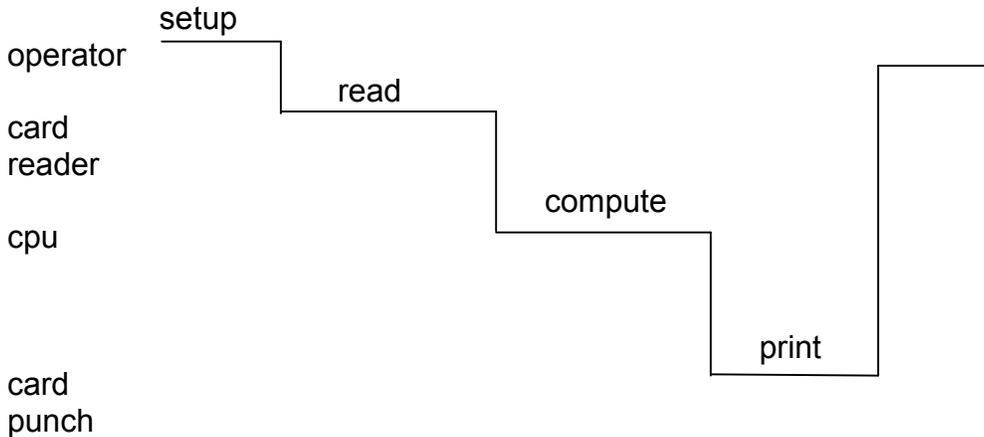


## Operator-Driven Shop

Avoid cpu idle time  
Operator loads jobs



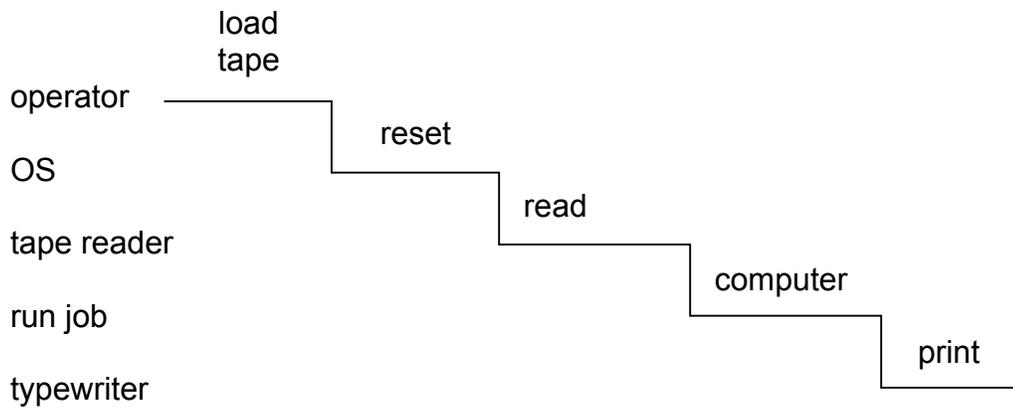
Users charged just for cpu time they use  
Batching of similar job steps  
Priority – users pay more for quick turnaround



## Offline Transport (I/O)

Automate I/O on separate (offline) computer, a.k.a., channel, satellite computer, peripheral processing unit (pppu)

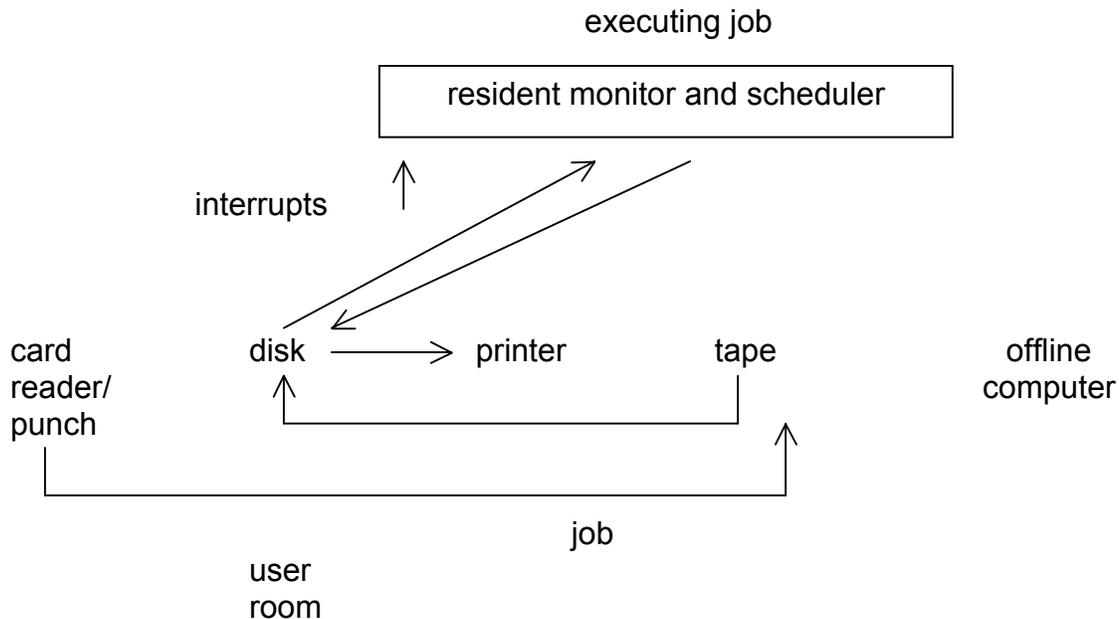
Beginning of a real OS, called resident monitor  
Reset machine after each job  
load next job  
accounting

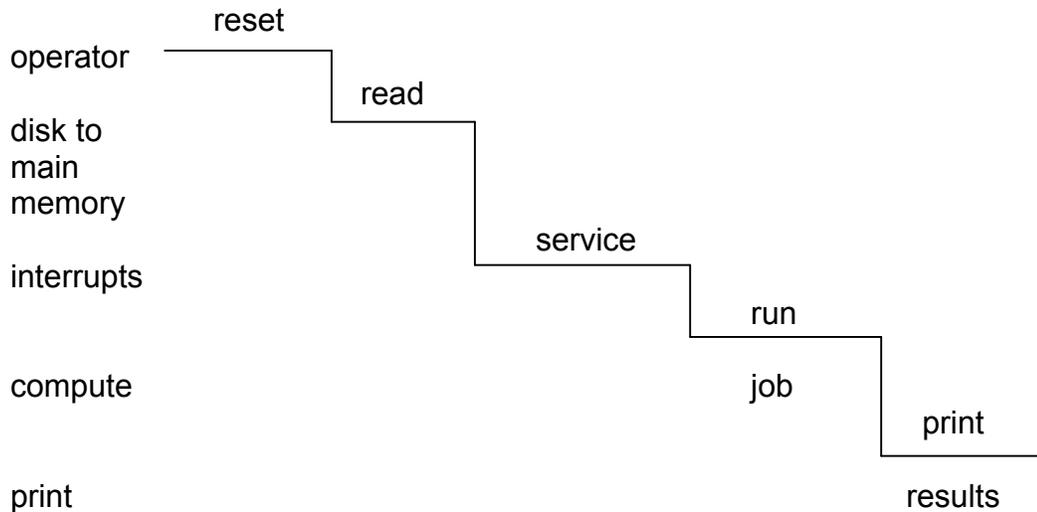


### Spooling

I/O concurrently with running jobs  
 I/O devices generate interrupts when I/O requests finished  
 Include disks in I/O device set, hold jobs  
 Add scheduler to OS

Utilization for spooling  
 Sufficient number of card readers and printers so always jobs ready to run  
 read/write to disks is faster  
 Computer spends a certain percent of time servicing interrupts for transport; not counted as useful computation time  
 OS resets machine between jobs

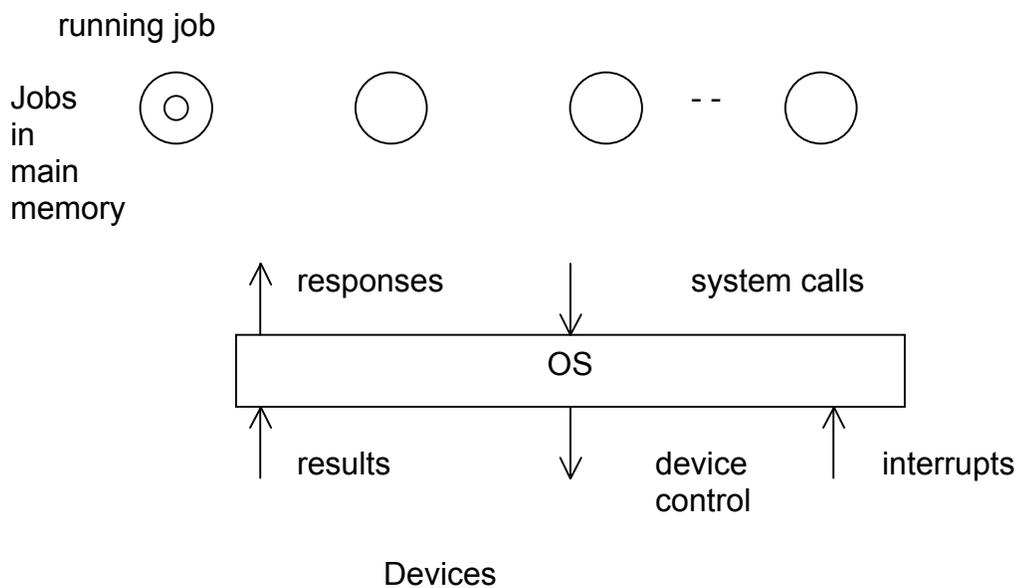




Note: Interrupts can occur at any time. I lumped the interrupt service activity into one interval.

### Batch Multiprocessing/Multiprogramming

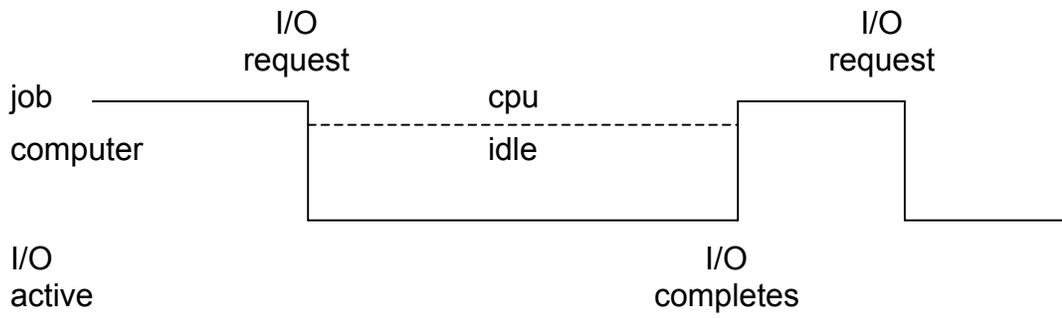
More than one job in main memory ready to run when current job performs an I/O request  
 Now a real OS that offers system calls



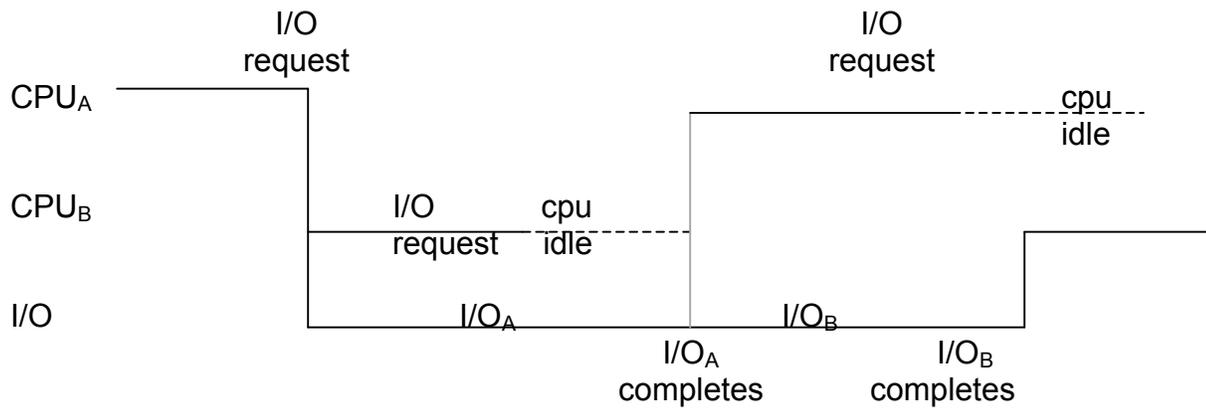
Life of a process:

Compute; I/O request; compute; I/O request; . . .

With only one job/process ready to run, CPU is often idle when I/O requests are serviced



CPU is better utilized when computation can proceed concurrently with I/O



CPU idle time is reduced with more jobs in main memory; but how many jobs can main memory hold?