

P2-0: **Text Files**

- All files are represented as binary digits – including text files
- Each character is represented by an integer code
 - ASCII – American Standard Code for Information Interchange
- Text file is a sequence of binary digits which represent the codes for each character.

P2-1: **ASCII**

- Each character can be represented as an 8-bit number
 - ASCII for a = 97 = 01100001
 - ASCII for b = 98 = 01100010
- Text file is a sequence of 1's and 0's which represent ASCII codes for characters in the file
 - File “aba” is 97, 97, 98
 - 011000010110001001100001

P2-2: **ASCII**

- Each character in ASCII is represented as 8 bits
 - We need 8 bits to represent all possible character combinations
 - (including control characters, and unprintable characters)
 - Breaking up file into individual characters is easy
 - Finding the kth character in a file is easy

P2-3: **ASCII**

- ASCII is not terribly efficient
 - All characters require 8 bits
 - Frequently used characters require the same number of bits as infrequently used characters
 - We could be more efficient if frequently used characters required fewer than 8 bits, and less frequently used characters required more bits

P2-4: **Representing Codes as Trees**

- Want to encode 4 only characters: a, b, c, d (instead of 256 characters)
 - How many bits are required for each code, if each code has the same length?

P2-5: **Representing Codes as Trees**

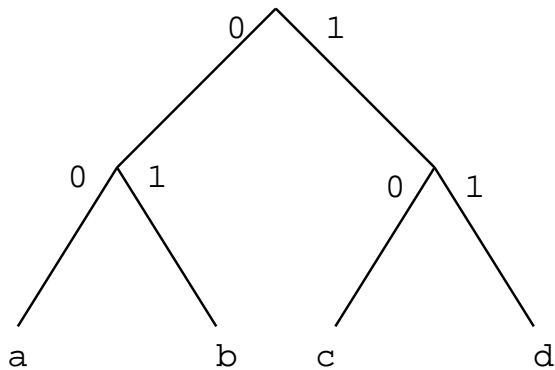
- Want to encode 4 only characters: a, b, c, d (instead of 256 characters)
 - How many bits are required for each code, if each code has the same length?
 - 2 bits are required, since there are 4 possible options to distinguish

P2-6: **Representing Codes as Trees**

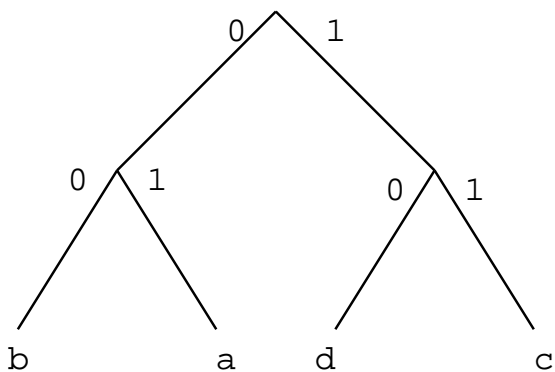
- Want to encode 4 only characters: a, b, c, d
- Pick the following codes:
 - a: 00
 - b: 01
 - c: 10
 - d: 11
- We can represent these codes as a tree
 - Characters are stored at the leaves of the tree
 - Code is represented by path to leaf

P2-7: Representing Codes as Trees

- a: 00, b: 01, c: 10, d:11

**P2-8: Representing Codes as Trees**

- a: 01, b: 00, c: 11, d:10

**P2-9: Prefix Codes**

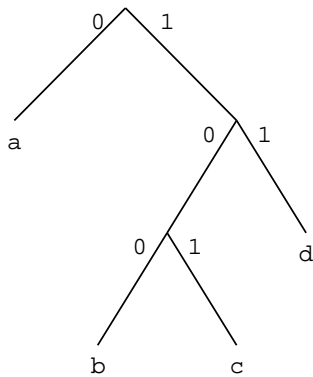
- If no code is a prefix of any other code, then decoding the file is unambiguous.
- If all codes are the same length, then no code will be a prefix of any other code (trivially)
- We can create variable length codes, where no code is a prefix of any other code

P2-10: **Variable Length Codes**

- Variable length code example:
 - a: 0, b: 100, c: 101, d: 11
- Decoding examples:
 - 100
 - 10011
 - 01101010010011

P2-11: **Prefix Codes & Trees**

- Any prefix code can be represented as a tree
- a: 0, b: 100, c: 101, d: 11



P2-12: **File Length**

- If we use the code:
 - a:00, b:01, c:10, d:11

How many bits are required to encode a file of 20 characters?

P2-13: **File Length**

- If we use the code:
 - a:00, b:01, c:10, d:11

How many bits are required to encode a file of 20 characters?

- 20 characters * 2 bits/character = 40 bits

P2-14: **File Length**

- If we use the code:
 - a:0, b:100, c:101, d:11

How many bits are required to encode a file of 20 characters?

P2-15: **File Length**

- If we use the code:
 - a:0, b:100, c:101, d:11

How many bits are required to encode a file of 20 characters?
- It depends upon the number of a's, b's, c's and d's in the file

P2-16: **File Length**

- If we use the code:
 - a:0, b:100, c:101, d:11

How many bits are required to encode a file of:

 - 11 a's, 2 b's, 2 c's, and 5 d's?

P2-17: **File Length**

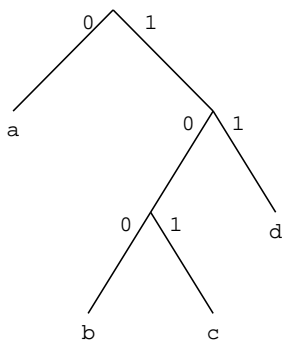
- If we use the code:
 - a:0, b:100, c:101, d:11

How many bits are required to encode a file of:

 - 11 a's, 2 b's, 2 c's, and 5 d's?
- $11*1 + 2*3 + 2*3 + 5*2 = 33 ; 40$

P2-18: **Decoding Files**

- We can use variable length keys to encode a text file
- Given the encoded file, and the tree representation of the codes, it is easy to decode the file



- 0111001010011

P2-19: **Decoding Files**

- We can use variable length keys to encode a text file
- Given the encoded file, and the tree representation of the codes, it is easy to decode the file
- Finding the kth character in the file is more tricky

P2-20: Decoding Files

- We can use variable length keys to encode a text file
- Given the encoded file, and the tree representation of the codes, it is easy to decode the file
- Finding the kth character in the file is more tricky
 - Need to decode the first (k-1) characters in the file, to determine where the kth character is in the file

P2-21: File Compression

- We can use variable length codes to compress files
 - Select an encoding such that frequently used characters have short codes, less frequently used characters have longer codes
 - Write out the file using these codes
 - (If the codes are dependent upon the contents of the file itself, we will also need to write out the codes at the beginning of the file for decoding)

P2-22: File Compression

- We need a method for building codes such that:
 - Frequently used characters are represented by leaves high in the code tree
 - Less Frequently used characters are represented by leaves low in the code tree
 - Characters of equal frequency have equal depths in the code tree

P2-23: Huffman Coding

- For each code tree, we keep track of the total number of times the characters in that tree appear in the input file
- We start with one code tree for each character that appears in the input file
- We combine the two trees with the lowest frequency, until all trees have been combined into one tree

P2-24: Huffman Coding

- Example: If the letters a-e have the frequencies:
 - a: 100, b: 20, c:15, d: 30, e: 1

a:100

b:20

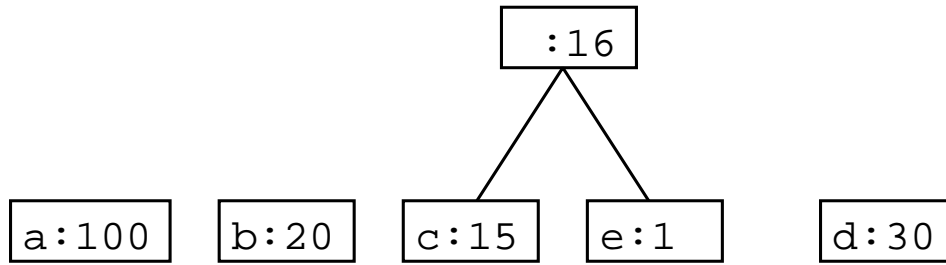
c:15

d:30

e:1

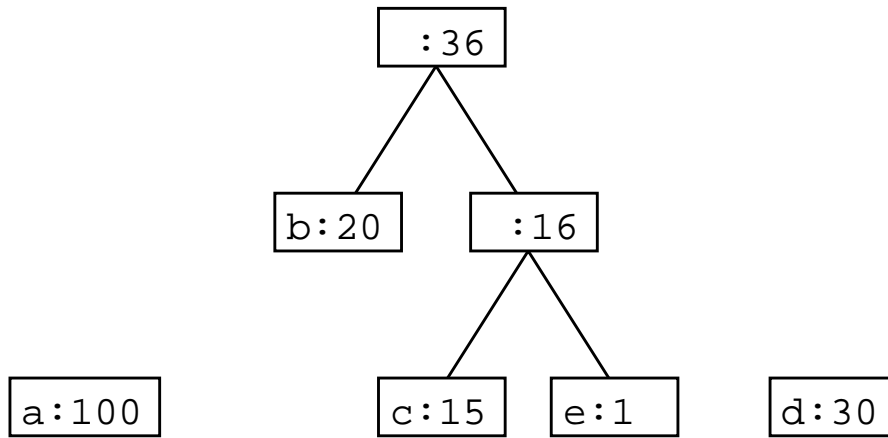
P2-25: Huffman Coding

- Example: If the letters a-e have the frequencies:
 - a: 100, b: 20, c:15, d: 30, e: 1



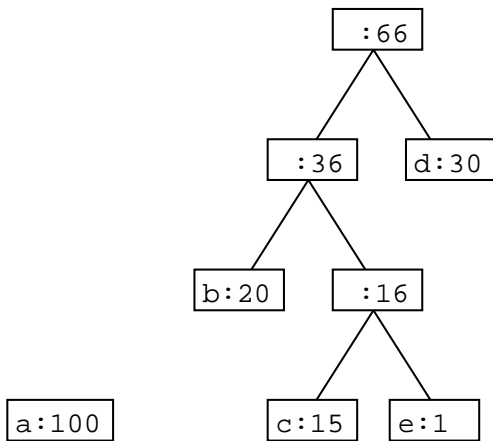
P2-26: **Huffman Coding**

- Example: If the letters a-e have the frequencies:
 - a: 100, b: 20, c:15, d: 30, e: 1



P2-27: **Huffman Coding**

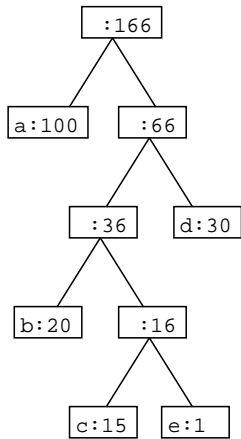
- Example: If the letters a-e have the frequencies:
 - a: 100, b: 20, c:15, d: 30, e: 1



P2-28: **Huffman Coding**

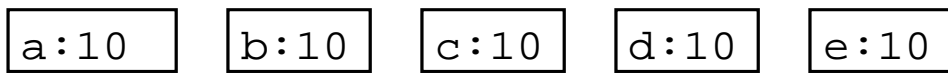
- Example: If the letters a-e have the frequencies:

- a: 100, b: 20, c:15, d: 30, e: 1



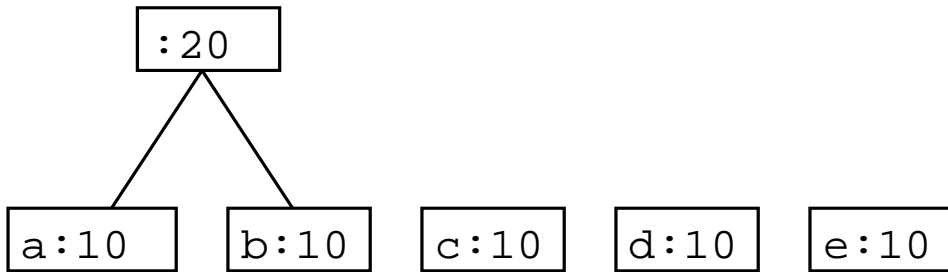
P2-29: **Huffman Coding**

- Example: If the letters a-e have the frequencies:
 - a: 10, b: 10, c:10, d: 10, e: 10



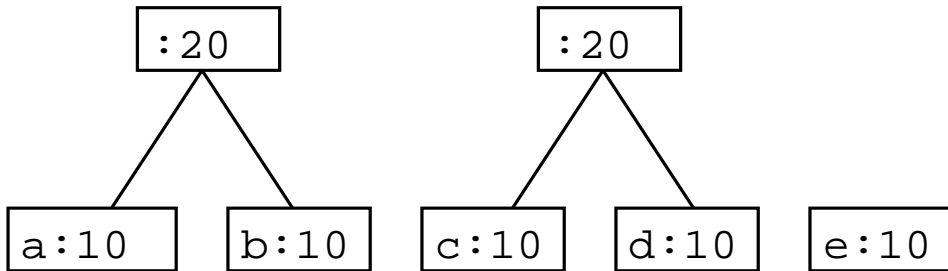
P2-30: **Huffman Coding**

- Example: If the letters a-e have the frequencies:
 - a: 10, b: 10, c:10, d: 10, e: 10



P2-31: **Huffman Coding**

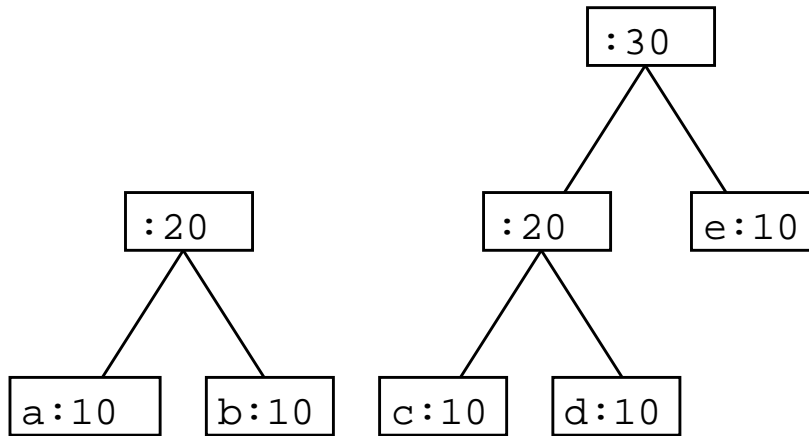
- Example: If the letters a-e have the frequencies:
 - a: 10, b: 10, c:10, d: 10, e: 10



P2-32: **Huffman Coding**

- Example: If the letters a-e have the frequencies:

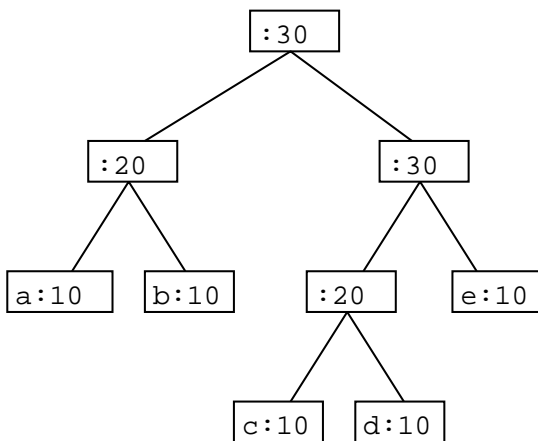
- a: 10, b: 10, c:10, d: 10, e: 10



P2-33: **Huffman Coding**

- Example: If the letters a-e have the frequencies:

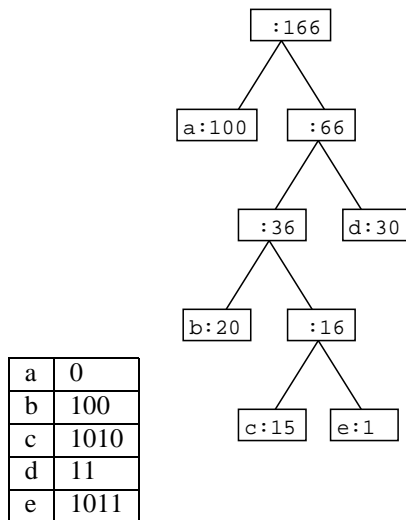
- a: 10, b: 10, c:10, d: 10, e: 10



P2-34: **Huffman Trees & Tables**

- Once we have a Huffman tree, decoding a file is straightforward – but *encoding* a tree requires a bit more information.
- Given just the tree, finding an encoding can be difficult
- ... What would we like to have, to help with encoding?

P2-35: **Encoding Tables**

P2-36: **Creating Encoding Table**

- Traverse the tree
 - Keep track of the path during the traversal
- When a leaf is reached, store the path in the table

P2-37: **Huffman Coding**

- To compress a file using huffman coding:
 - Read in the file, and count the occurrence of each character, and built a frequency table
 - Build the Huffman tree from the frequencies
 - Build the Huffman codes from the tree
 - Print the Huffman tree to the output file (for use in decompression)
 - Print out the codes for each character

P2-38: **Huffman Coding**

- To uncompress a file using huffman coding:
 - Read in the Huffman tree from the input file
 - Read the input file bit by bit, traversing the Huffman tree as you go
 - When a leaf is read, write the appropriate file to an output file

P2-39: **Magic Numbers**

- We only want to uncompress files that we actually compressed ourselves
- Write a “Magic Number” at the beginning of compressed file
 - For this project, use 0x4846 (ASCII 'HF')
- When uncompressing, first check to see that magic number is correct

P2-40: **Binary Files**

```
public BinaryFile(String filename,  
                  char readOrWrite)
```

```
public boolean EndOfFile()  
public char readChar()  
public void writeChar(char c)
```

```
public boolean readBit()  
public void writeBit(boolean bit)
```

```
public void close()
```

P2-41: **Binary Files**

- readBit
 - Read a single bit
- readChar
 - Read a single character (8 bits)

P2-42: **Binary Files**

- writeBit
 - Writes out a single bit
- writeChar
 - Writes out a single (8 bit) character

P2-43: **Binary Files**

- If we write to a binary file:
 - bit, bit, char, bit, int
- And then read from the file:
 - bit, char, bit, int, bit
- What will we get out?

P2-44: **Binary Files**

- If we write to a binary file:
 - bit, bit, char, bit, int
- And then read from the file:
 - bit, char, bit, int, bit
- What will we get out?
- Garbage! (except for the first bit)

P2-45: **Printing out Trees**

- To print out Huffman trees:
 - Print out nodes in pre-order traversal
 - Need a way of denoting which nodes are leaves and which nodes are interior nodes
 - (Huffman trees are full – every node has 0 or 2 children)
 - For each interior node, print out a 0 (single bit). For each leaf, print out a 1, followed by 8 bits for the character at the leaf

P2-46: **Compression?**

- Is it possible that Huffman compression would not compress the file?
- Is it possible that Huffman compression could actually make the file larger?
- How?

P2-47: **Compression?**

- What happens if all the characters have the same frequency?
 - What does the tree look like?
 - What can we say about the lengths of the codes for each character?
 - What does that mean for the file size?

P2-48: **Compression?**

- What happens if all the characters have the same frequency?
 - All nodes are at the same depth in the tree (that is, 8)
 - Each code will have a length of 8
 - The encoded file will be the same size as the original file – *plus the size required to encode the tree*

P2-49: **Compression!**

- What to do?
 - Calculate the size of the input file
 - Calculate the size that the compressed file would be
 - If the compressed file is larger than the input file, don't compress

P2-50: **Compression!**

- Given the frequency array, how large is the input file?

P2-51: **Compression!**

- Given the frequency array, how large is the input file?
 - $\sum_c freq(c) * len(c)$
 - (# of characters in the input file) * 8

P2-52: Compression!

- Given the frequency array & codes for each compressed element, how large is the compressed file?

P2-53: Compression!

- Given the frequency array & codes for each compressed element, how large is the compressed file?
 - $\sum_c freq(c) * len(c) +$
 - Size of tree +
 - + 4 bytes (32 bits) (file overhead) + 2 (16 bits) bytes (Magic Number)
 - Implementation detail of BinaryFile class
- Note file sizes need to be a multiple of 8 bits ...

P2-54: Command Line Arguments

```
public static void main(String args[])
```

- The `args` parameter holds the input parameters
- `java MyProgram arg1 arg2 arg3`
 - `args.length = 3`
 - `args[0] = "arg1"`
 - `args[1] = "arg2"`
 - `args[2] = "arg3"`

P2-55: Calling Huffman

```
java Huffman (-c|-u) [-v] [-f] infile outfile
```

- `(-c|-u)` stands for either `"-c"` (for compress), or `"-u"` (for uncompress)
- `[-v]` stands for an optional `"-v"` flag (for verbose)
- `[-f]` stands for an optional `"-f"` flag (for force compress)
- `infile` is the input file
- `outfile` is the output file