Data Structures and Algorithms

CS245-2015S-10

Sorting

David Galles

Department of Computer Science
University of San Francisco

10-0: Main Memory Sorting

- All data elements can be stored in memory at the same time
- Data stored in an array, indexed from $0 \ldots n-1$, where n is the number of elements
- Each element has a key value (accessed with a `key()` method)
- We can compare keys for <, >, =
- For illustration, we will use arrays of integers – though often keys will be strings, other Comparable types

10-1: Stable Sorting

- A sorting algorithm is *Stable* if the relative order of duplicates is preserved
- The order of duplicates matters if the *keys* are duplicated, but the *records* are not.

3	1	2	1	1	2	3	Key
Bob	Joe	Ed	Amy	Sue	Al	Bud	Data

1	1	1	2	2	3	3	Key
Amy	Joe	Sue	Ed	Al	Bob	Bud	Data

A *non*-Stable sort

10-2: Insertion Sort

- Separate list into sorted portion, and unsorted portion
- Initially, sorted portion contains first element in the list, unsorted portion is the rest of the list
  - (A list of one element is always sorted)
- Repeatedly insert an element from the unsorted list into the sorted list, until the list is sorted

10-3: $\Theta()$ For Insertion Sort

- Running time $\propto$ # of comparisons
- Worst Case:

10-4: $\Theta()$ For Insertion Sort

- Running time $\propto$ # of comparisons
- Worst Case: Inverse sorted list

# of comparisons:

10-5: $\Theta()$ For Insertion Sort

- Running time $\propto$ # of comparisons
- Worst Case: Inverse sorted list

# of comparisons:

$$\sum_{i=1}^{n-1} i \in \Theta(n^2)$$

10-6: $\Theta()$ For Insertion Sort

- Running time $\propto$ # of comparisons
- Best Case:

10-7: $\Theta()$ For Insertion Sort

- Running time $\propto$ # of comparisons
- Best Case: Sorted List

# of comparisons:

10-8: $\Theta()$ For Insertion Sort

- Running time $\propto$ # of comparisons
- Best Case: Sorted List

# of comparisons:

$$n - 1$$

10-9: Bubble Sort

- Scan list from the last index to index 0, swapping the smallest element to the front of the list
- Scan the list from the last index to index 1, swapping the second smallest element to index 1
- Scan the list from the last index to index 2, swapping the third smallest element to index 2

  ...
- Swap the second largest element into position $(n-2)$

10-10: $\Theta()$ for Bubble Sort

- Running time $\propto$ # of comparisons
- Number of Comparisons:

10-11: $\Theta()$ for Bubble Sort

- Running time $\propto$ # of comparisons
- Number of Comparisons:

$$\sum_{i=1}^{n-1} i \in \Theta(n^2)$$

10-12: Selection Sort

- Scan through the list, and find the smallest element
- Swap smallest element into position 0
- Scan through the list, and find the second smallest element
- Swap second smallest element into position 1

  . . .
- Scan through the list, and find the second largest element
- Swap smallest largest into position $n - 2$

10-13: $\Theta()$ for Selection Sort

- Running time $\propto$ # of comparisons
- Number of Comparisons:

10-14: $\Theta()$ for Selection Sort

- Running time $\propto$ # of comparisons
- Number of Comparisons:

$$\sum_{i=1}^{n-1} i \in \Theta(n^2)$$

10-15: Improving Insertion Sort

- Insertion sort is fast if a list is “almost sorted”
- How can we use this?
  - Do some work to make the list “almost sorted”
  - Run insertion sort to finish sorting the list
- Only helps if work required to make list “almost sorted” is less than n^2

10-16: Shell Sort

- Sort $n/2$ sublists of length 2, using insertion sort
- Sort $n/4$ sublists of length 4, using insertion sort
- Sort $n/8$ sublists of length 8, using insertion sort

  $\ldots$
- Sort 2 sublists of length $n/2$, using insertion sort
- Sort 1 sublist of length n, using insertion sort

10-17: Shell's Increments

- Shell sort runs several insertion sorts, using increments
  - Code on monitor uses "Shell's Increments": $\{n/2, n/4, \ldots 4, 2, 1\}$
- Problem with Shell's Increments:
  - Various sorts do not interact much
  - If all large elements are stored in large indices, and small elements are stored in even indices, what happens?

10-18: Other Increments

- Shell's Increments: $\{n/2, n/4, \ldots 4, 2, 1\}$
  - Running time: $O(n^2)$
- "$/3$" increments: $\{n/3, n/9, \ldots, 9, 3, 1\}$
  - Running time: $O(n^{\frac{3}{2}})$
- Hibbard's Increments: $\{2^k - 1, 2^{k-1} - 1, \ldots 7, 3, 1\}$
  - Running time: $O(n^{\frac{3}{2}})$

10-19: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?

10-20: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?

10-21: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?
    - $\Theta(n)$
  - How Many Passes?

10-22: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?
    - $\Theta(n)$
  - How Many Passes?
    - $\lg n$
  - Total running time?

10-23: Shell Sort: Best case

- What is the best case running time for Shell Sort (using Shell's increments)
  - When would the best case occur?
    - When the list was originally sorted
  - How long would each pass through Shell Sort take?
    - $\Theta(n)$
  - How Many Passes?
    - $\lg n$
  - Total running time?
    - $\Theta(n \lg n)$

10-24: Stability

- Is Insertion sort stable?
- Is Bubble Sort stable?
- Is Selection Sort stable?
- Is Shell Sort stable?

10-25: Stability

- Is Insertion sort stable? Yes!
- Is Bubble Sort stable? Yes!
- Is Selection Sort stable? No!
- Is Shell Sort stable? No!

Note that minor changes to the stable sorting algorithms will make them unstable (for instance, swaping $A[i]$ and $A[i+1]$ when $A[i] \geq A[i+1]$, not just when $A[i] > A[i+1]$