

03-0: **Algorithm Analysis**

```
for (i=1; i<=n*n; i++)
  for (j=0; j<i; j++)
    sum++;
```

03-1: **Algorithm Analysis**

```
for (i=1; i<=n*n; i++)   Executed n*n times
  for (j=0; j<i; j++)     Executed <= n*n times
    sum++;                O(1)
```

Running Time:  $O(n^4)$  03-2: **Algorithm Analysis**

```
for (i=1; i<=n*n; i++)
  for (j=0; j<i; j++)
    sum++;
```

Exact # of times sum++ is executed:

$$\begin{aligned} \sum_{i=1}^{n^2} i &= \frac{n^2(n^2 + 1)}{2} \\ &= \frac{n^4 + n^2}{2} \\ &\in \Theta(n^4) \end{aligned}$$

03-3: **Recursive Functions**

```
long power(long x, long n) {
  if (n == 0)
    return 1;
  else
    return x * power(x, n-1);
}
```

03-4: **Recurrence Relations**

$T(n)$  = Time required to solve a problem of size  $n$

Recurrence relations are used to determine the running time of recursive programs – recurrence relations themselves are recursive

$T(0)$  = time to solve problem of size 0  
– Base Case

$T(n)$  = time to solve problem of size  $n$   
– Recursive Case

03-5: **Recurrence Relations**

```

long power(long x, long n) {
    if (n == 0)
        return 1;
    else
        return x * power(x, n-1);
}

```

$$T(0) = c_1 \quad \text{for some constant } c_1$$

$$T(n) = c_2 + T(n-1) \quad \text{for some constant } c_2$$

### 03-6: Solving Recurrence Relations

$$T(0) = c_1$$

$$T(n) = T(n-1) + c_2$$

If we knew  $T(n-1)$ , we could solve  $T(n)$ .

$$T(n) = T(n-1) + c_2$$

### 03-7: Solving Recurrence Relations

$$T(0) = c_1$$

$$T(n) = T(n-1) + c_2$$

If we knew  $T(n-1)$ , we could solve  $T(n)$ .

$$\begin{aligned}
 T(n) &= T(n-1) + c_2 & T(n-1) &= T(n-2) + c_2 \\
 &= T(n-2) + c_2 + c_2 \\
 &= T(n-2) + 2c_2
 \end{aligned}$$

### 03-8: Solving Recurrence Relations

$$T(0) = c_1$$

$$T(n) = T(n-1) + c_2$$

If we knew  $T(n-1)$ , we could solve  $T(n)$ .

$$\begin{aligned}
 T(n) &= T(n-1) + c_2 & T(n-1) &= T(n-2) + c_2 \\
 &= T(n-2) + c_2 + c_2 \\
 &= T(n-2) + 2c_2 & T(n-2) &= T(n-3) + c_2 \\
 &= T(n-3) + c_2 + 2c_2 \\
 &= T(n-3) + 3c_2
 \end{aligned}$$

### 03-9: Solving Recurrence Relations

$$T(0) = c_1$$

$$T(n) = T(n-1) + c_2$$

If we knew  $T(n-1)$ , we could solve  $T(n)$ .

$$\begin{aligned}
 T(n) &= T(n-1) + c_2 & T(n-1) &= T(n-2) + c_2 \\
 &= T(n-2) + c_2 + c_2 \\
 &= T(n-2) + 2c_2 & T(n-2) &= T(n-3) + c_2 \\
 &= T(n-3) + c_2 + 2c_2 \\
 &= T(n-3) + 3c_2 & T(n-3) &= T(n-4) + c_2 \\
 &= T(n-4) + 4c_2
 \end{aligned}$$

### 03-10: Solving Recurrence Relations

$$T(0) = c_1$$

$$T(n) = T(n-1) + c_2$$

If we knew  $T(n-1)$ , we could solve  $T(n)$ .

$$\begin{aligned}
 T(n) &= T(n-1) + c_2 & T(n-1) &= T(n-2) + c_2 \\
 &= T(n-2) + c_2 + c_2 \\
 &= T(n-2) + 2c_2 & T(n-2) &= T(n-3) + c_2 \\
 &= T(n-3) + c_2 + 2c_2 \\
 &= T(n-3) + 3c_2 & T(n-3) &= T(n-4) + c_2 \\
 &= T(n-4) + 4c_2 \\
 &= \dots \\
 &= T(n-k) + kc_2
 \end{aligned}$$

03-11: Solving Recurrence Relations

$$\begin{aligned}
 T(0) &= c_1 \\
 T(n) &= T(n-k) + k * c_2 \quad \text{for all } k
 \end{aligned}$$

If we set  $k = n$ , we have:

$$\begin{aligned}
 T(n) &= T(n-n) + nc_2 \\
 &= T(0) + nc_2 & \text{03-12: Building a Better Power} \\
 &= c_1 + nc_2 \\
 &\in \Theta(n)
 \end{aligned}$$

```

long power(long x, long n) {
    if (n==0) return 1;
    if (n==1) return x;
    if ((n % 2) == 0)
        return power(x*x, n/2);
    else
        return power(x*x, n/2) * x;
}

```

03-13: Building a Better Power

```

long power(long x, long n) {
    if (n==0) return 1;
    if (n==1) return x;
    if ((n % 2) == 0)
        return power(x*x, n/2);
    else
        return power(x*x, n/2) * x;
}

```

$$T(0) = c_1$$

$$T(1) = c_2$$

$$T(n) = T(n/2) + c_3$$

(Assume  $n$  is a power of 2)

03-14: Solving Recurrence Relations

$$T(n) = T(n/2) + c_3 \quad \text{03-15: Solving Recurrence Relations}$$

$$\begin{aligned}
 T(n) &= T(n/2) + c_3 & T(n/2) &= T(n/4) + c_3 \\
 &= T(n/4) + c_3 + c_3 \\
 &= T(n/4) + 2c_3
 \end{aligned}$$

03-16: Solving Recurrence Relations

$$\begin{aligned}
 T(n) &= T(n/2) + c_3 & T(n/2) &= T(n/4) + c_3 \\
 &= T(n/4) + c_3 + c_3 \\
 &= T(n/4) + 2c_3 & T(n/4) &= T(n/8) + c_3 \\
 &= T(n/8) + c_3 + 2c_3 \\
 &= T(n/8) + 3c_3
 \end{aligned}$$

## 03-17: Solving Recurrence Relations

$$\begin{aligned}
T(n) &= T(n/2) + c_3 & T(n/2) &= T(n/4) + c_3 \\
&= T(n/4) + c_3 + c_3 \\
&= T(n/4) + 2c_3 & T(n/4) &= T(n/8) + c_3 \\
&= T(n/8) + c_3 + 2c_3 \\
&= T(n/8) + 3c_3 & T(n/8) &= T(n/16) + c_3 \\
&= T(n/16) + c_3 + 3c_3 \\
&= T(n/16) + 4c_3
\end{aligned}$$

## 03-18: Solving Recurrence Relations

$$\begin{aligned}
T(n) &= T(n/2) + c_3 & T(n/2) &= T(n/4) + c_3 \\
&= T(n/4) + c_3 + c_3 \\
&= T(n/4) + 2c_3 & T(n/4) &= T(n/8) + c_3 \\
&= T(n/8) + c_3 + 2c_3 \\
&= T(n/8) + 3c_3 & T(n/8) &= T(n/16) + c_3 \\
&= T(n/16) + c_3 + 3c_3 \\
&= T(n/16) + 4c_3 & T(n/16) &= T(n/32) + c_3 \\
&= T(n/32) + c_3 + 4c_3 \\
&= T(n/32) + 5c_3
\end{aligned}$$

## 03-19: Solving Recurrence Relations

$$\begin{aligned}
T(n) &= T(n/2) + c_3 & T(n/2) &= T(n/4) + c_3 \\
&= T(n/4) + c_3 + c_3 \\
&= T(n/4) + 2c_3 & T(n/4) &= T(n/8) + c_3 \\
&= T(n/8) + c_3 + 2c_3 \\
&= T(n/8) + 3c_3 & T(n/8) &= T(n/16) + c_3 \\
&= T(n/16) + c_3 + 3c_3 \\
&= T(n/16) + 4c_3 & T(n/16) &= T(n/32) + c_3 \\
&= T(n/32) + c_3 + 4c_3 \\
&= T(n/32) + 5c_3 \\
&= \dots \\
&= T(n/2^k) + kc_3
\end{aligned}$$

## 03-20: Solving Recurrence Relations

$$\begin{aligned}
T(0) &= c_1 \\
T(1) &= c_2 \\
T(n) &= T(n/2) + c_3 \\
\\
T(n) &= T(n/2^k) + kc_3
\end{aligned}$$

We want to get rid of  $T(n/2^k)$ . Since we know  $T(1) \dots$

$$\begin{aligned}
n/2^k &= 1 \\
n &= 2^k \\
\lg n &= k
\end{aligned}$$

## 03-21: Solving Recurrence Relations

$$\begin{aligned}
T(1) &= c_2 \\
T(n) &= T(n/2^k) + kc_3
\end{aligned}$$

Set  $k = \lg n$ :

$$T(n) = T(n/2^{\lg n}) + (\lg n)c_3$$

$$\begin{aligned}
 &= T(n/n) + c_3 \lg n \\
 &= T(1) + c_3 \lg n \\
 &= c_2 + c_3 \lg n \\
 &\in \Theta(\lg n)
 \end{aligned}$$

03-22: Power **Modifications**

```

long power(long x, long n) {
    if (n==0) return 1;
    if (n==1) return x;
    if ((n % 2) == 0)
        return power(x*x, n/2);
    else
        return power(x*x, n/2) * x;
}

```

03-23: Power **Modifications**

```

long power(long x, long n) {
    if (n==0) return 1;
    if (n==1) return x;
    if ((n % 2) == 0)
        return power(power(x,2), n/2);
    else
        return power(power(x,2), n/2) * x;
}

```

This version of power will not work. Why?

03-24: Power **Modifications**

```

long power(long x, long n) {
    if (n==0) return 1;
    if (n==1) return x;
    if ((n % 2) == 0)
        return power(power(x,n/2), 2);
    else
        return power(power(x,n/2), 2) * x;
}

```

This version of power also will not work. Why?

03-25: Power **Modifications**

```

long power(long x, long n) {
    if (n==0) return 1;
    if (n==1) return x;
    if ((n % 2) == 0)
        return power(x,n/2) * power(x,n/2);
    else
        return power(x,n/2) * power(x,n/2) * x;
}

```

This version of power does work.

What is the recurrence relation that describes its running time?

03-26: Power **Modifications**

```
long power(long x, long n) {
    if (n==0) return 1;
    if (n==1) return x;
    if ((n % 2) == 0)
        return power(x, n/2) * power(x, n/2);
    else
        return power(x, n/2) * power(x, n/2) * x;
}
```

$$\begin{aligned} T(0) &= c_1 \\ T(1) &= c_2 \\ T(n) &= T(n/2) + T(n/2) + c_3 \\ &= 2T(n/2) + c_3 \end{aligned}$$

(Again, assume  $n$  is a power of 2)

03-27: **Solving Recurrence Relations**

$$\begin{aligned} T(n) &= 2T(n/2) + c_3 & T(n/2) &= 2T(n/4) + c_3 \\ &= 2[2T(n/4) + c_3] + c_3 \\ &= 4T(n/4) + 3c_3 & T(n/4) &= 2T(n/8) + c_3 \\ &= 4[2T(n/8) + c_3] + 3c_3 \\ &= 8T(n/8) + 7c_3 \\ &= 8[2T(n/16) + c_3] + 7c_3 \\ &= 16T(n/16) + 15c_3 \\ &= 32T(n/32) + 31c_3 \\ &\dots \\ &= 2^k T(n/2^k) + (2^k - 1)c_3 \end{aligned}$$

03-28: **Solving Recurrence Relations**

$$\begin{aligned} T(0) &= c_1 \\ T(1) &= c_2 \\ T(n) &= 2^k T(n/2^k) + (2^k - 1)c_3 \end{aligned}$$

Pick a value for  $k$  such that  $n/2^k = 1$ :

$$\begin{aligned} n/2^k &= 1 \\ n &= 2^k \\ \lg n &= k \\ T(n) &= 2^{\lg n} T(n/2^{\lg n}) + (2^{\lg n} - 1)c_3 \\ &= nT(n/n) + (n - 1)c_3 \\ &= nT(1) + (n - 1)c_3 \\ &= nc_2 + (n - 1)c_3 \\ &\in \Theta(n) \end{aligned}$$

03-29: **Recursion Trees**

- We can also do this substitution visually, leads to Recursion Trees
- Consider:

$$\begin{aligned} T(n) &= 2T(n/2) + Cn \\ T(1) &= C_2 \\ T(0) &= C_2 \end{aligned}$$

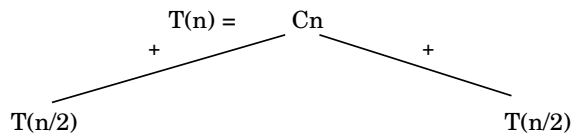
03-30: **Recursion Trees**

- Start with the recursive definition

$$T(n) = Cn + 2T(n/2)$$

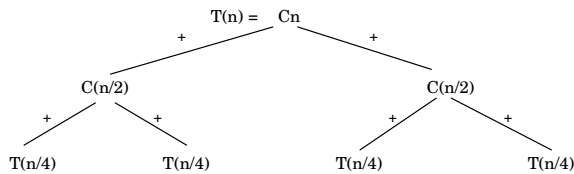
03-31: **Recursion Trees**

- Move the equation around a bit to get:



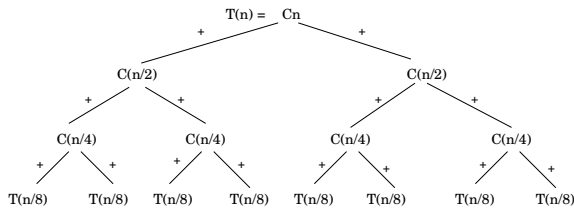
- Replace each occurrence of  $T(n/2)$  with  $T(n/4) + T(n/4) + C(n/2)$

03-32: **Recursion Trees**



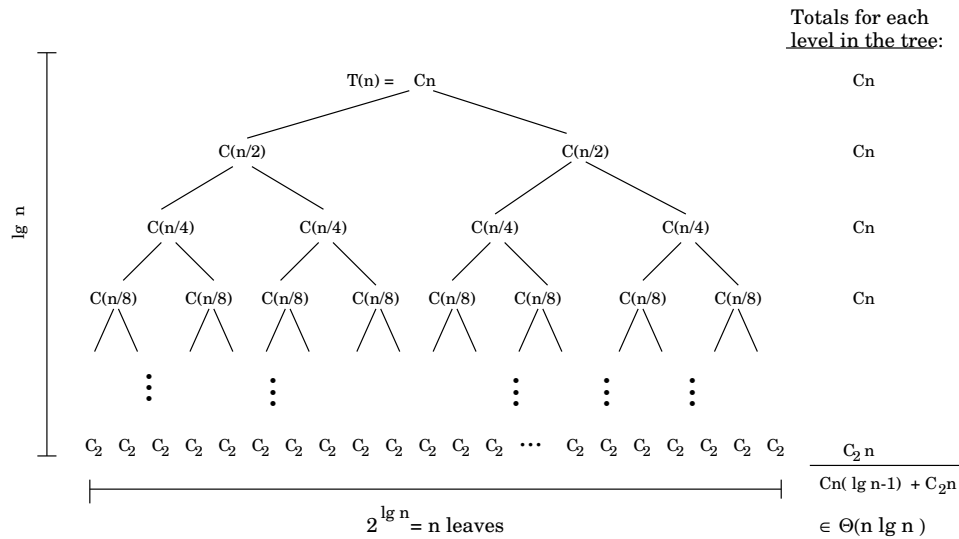
- Replace again, using  $T(n) = 2T(n/2) + Cn$

03-33: **Recursion Trees**



- If we continue replacing ...

03-34: **Recursion Trees**



03-35: Recursion Trees

$$T(1) = C_1$$

$$T(n) = T(n - 1) + C_2$$

03-36: Recursion Trees

$$T(0) = C_1$$

$$T(1) = C_1$$

$$T(n) = T(n/2) + C_2$$

03-37: Substitution Method

- We can prove that a bound is correct using induction, this is the substitution method

$$T(1) = C_1$$

$$T(n) = T(n - 1) + C_2$$

Show:  $T(n) \in O(?)$

03-38: Substitution Method

- We can prove that a bound is correct using induction, this is the substitution method

$$T(1) = C_1$$

$$T(n) = T(n - 1) + C_2$$



Show:  $T(n) \in O(n)$ , that is:

$$T(n) \leq C * n \text{ for all } n > n_0,$$

for some pair of constants  $C, n_0$

03-39: **Substitution Method**

$$T(1) = C_1$$

$$T(n) = T(n-1) + C_2$$

Show:  $T(n) \in O(n)$ , that is,  $T(n) \leq C * n$

- Base case:  $T(1) = C_1 \leq C * 1$  for some constant  $C$

This is true as long as  $C \geq C_1$ .

03-40: **Substitution Method**

$$T(1) = C_1$$

$$T(n) = T(n-1) + C_2$$

Show:  $T(n) \in O(n)$ , that is,  $T(n) \leq C * n$

- Recursive case:

$$T(n) = T(n-1) + C_2 \quad \text{Recurrence definition}$$

03-41: **Substitution Method**

$$T(1) = C_1$$

$$T(n) = T(n-1) + C_2$$

Show:  $T(n) \in O(n)$ , that is,  $T(n) \leq C * n$

- Recursive case:

$$T(n) = T(n-1) + C_2 \quad \text{Recurrence definition}$$

$$\leq C(n-1) + C_2 \quad \text{Inductive hypothesis}$$

03-42: **Substitution Method**

$$T(1) = C_1$$

$$T(n) = T(n-1) + C_2$$

Show:  $T(n) \in O(n)$ , that is,  $T(n) \leq C * n$

- Recursive case:

$$T(n) = T(n-1) + C_2 \quad \text{Recurrence definition}$$

$$\leq C(n-1) + C_2 \quad \text{Inductive hypothesis}$$

$$\leq Cn + (C_2 - C) \quad \text{Algebra}$$

$$\leq Cn \quad \text{If } C > C_2$$

This is true as long as  $C \geq C_1$ .

03-43: **Substitution Method**

- We can prove that a bound is correct using induction, this is the substitution method

$$T(1) = C_1$$

$$T(n) = T(n-1) + C_2$$

Show:  $T(n) \in \Omega(n)$

$T(n) \geq C * n$  for all  $n > n_0$ ,  
for some pair of constants  $C, n_0$

03-44: **Substitution Method**

$$\begin{aligned} T(1) &= C_1 \\ T(n) &= T(n-1) + C_2 \end{aligned}$$

Show:  $T(n) \in \Omega(n)$ , that is,  $T(n) \geq C * n$

- Base case:  $T(1) = C_1 \geq C * 1$  for some constant  $C$

This is true as long as  $C \leq C_1$ .

03-45: **Substitution Method**

$$\begin{aligned} T(1) &= C_1 \\ T(n) &= T(n-1) + C_2 \end{aligned}$$

Show:  $T(n) \in \Omega(n)$ , that is,  $T(n) \geq C * n$

- Recursive case:

$$T(n) = T(n-1) + C_2 \quad \text{Recurrence definition}$$

03-46: **Substitution Method**

$$\begin{aligned} T(1) &= C_1 \\ T(n) &= T(n-1) + C_2 \end{aligned}$$

Show:  $T(n) \in \Omega(n)$ , that is,  $T(n) \geq C * n$

- Recursive case:

$$\begin{aligned} T(n) &= T(n-1) + C_2 && \text{Recurrence definition} \\ &\geq C(n-1) + C_2 && \text{Inductive hypothesis} \end{aligned}$$

03-47: **Substitution Method**

$$\begin{aligned} T(1) &= C_1 \\ T(n) &= T(n-1) + C_2 \end{aligned}$$

Show:  $T(n) \in \Omega(n)$ , that is,  $T(n) \geq C * n$

- Recursive case:

$$\begin{aligned} T(n) &= T(n-1) + C_2 && \text{Recurrence definition} \\ &\geq C(n-1) + C_2 && \text{Inductive hypothesis} \\ &\geq Cn + (C_2 - C) && \text{Algebra} \\ &\geq Cn && \text{If } C \leq C_2 \end{aligned}$$

This is true as long as  $C \leq C_1$ .

03-48: **Substitution Method**

$$\begin{aligned} T(0) &= C_2 \\ T(1) &= C_2 \\ T(n) &= 2T(n/2) + C_1n \end{aligned}$$

Show:  $T(n) \in O(n \lg n)$ , that is,  $T(n) \leq C * n \lg n$

03-49: **Substitution Method**

$$\begin{aligned} T(0) &= C_2 \\ T(1) &= C_2 \\ T(n) &= 2T(n/2) + C_1n \end{aligned}$$

Show:  $T(n) \in O(n \lg n)$ , that is,  $T(n) \leq C * n \lg n$

- Base cases:
  - $T(0) = C_1 \leq C * 0 \lg 0$  for some constant  $C$
  - $T(1) = C_1 \leq C * 1 \lg 1$  for some constant  $C$

- 

Hmmm....

03-50: **Substitution Method**

$$T(0) = C_2$$

$$T(1) = C_2$$

$$T(n) = 2T(n/2) + C_1n$$

Show:  $T(n) \in O(n \lg n)$ , that is,  $T(n) \leq C * n \lg n$

- Only care about  $n > n_0$ . We can pick 2, 3 as base cases (why?)
  - $T(2) = C_1 \leq C * 2 \lg 2$  for some constant  $C$
  - $T(3) = C_1 \leq C * 3 \lg 3$  for some constant  $C$

- 

03-51: **Substitution Method**

$$T(0) = C_2$$

$$T(1) = C_2$$

$$T(n) = 2T(n/2) + C_1n$$

$$T(n) = 2T(n/2) + C_1n \quad \text{Recurrence Definition}$$

03-52: **Substitution Method**

$$T(0) = C_2$$

$$T(1) = C_2$$

$$T(n) = 2T(n/2) + C_1n$$

$$T(n) = 2T(n/2) + C_1n \quad \text{Recurrence Definition}$$

$$\leq 2C(n/2) \lg(n/2) + C_1n \quad \text{Inductive hypothesis}$$

03-53: **Substitution Method**

$$T(0) = C_2$$

$$T(1) = C_2$$

$$T(n) = 2T(n/2) + C_1n$$

$$T(n) = 2T(n/2) + C_1n \quad \text{Recurrence Definition}$$

$$\leq 2C(n/2) \lg(n/2) + C_1n \quad \text{Inductive hypothesis}$$

$$\leq Cn \lg n/2 + C_1n \quad \text{Algebra}$$

$$\leq Cn \lg n - Cn \lg 2 + C_1n \quad \text{Algebra}$$

$$\leq Cn \lg n - Cn + C_1n \quad \text{Algebra}$$

03-54: **Substitution Method**

$$T(0) = C_2$$

$$T(1) = C_2$$

$$T(n) = 2T(n/2) + C_1n$$

$$T(n) = 2T(n/2) + C_1n \quad \text{Recurrence Definition}$$

$$\leq 2C(n/2) \lg(n/2) + C_1n \quad \text{Inductive hypothesis}$$

$$\leq Cn \lg n/2 + C_1n \quad \text{Algebra}$$

$$\leq Cn \lg n - Cn \lg 2 + C_1n \quad \text{Algebra}$$

$$\leq Cn \lg n - Cn + C_1n \quad \text{Algebra}$$

$$\leq Cn \lg n \quad \text{If } C > C_1$$

03-55: **Substitution Method**

- Sometimes, the math doesn't work out in the substitution method:

$$T(1) = 1$$

$$T(n) = 2T\left(\frac{n}{2}\right) + 1$$

(Work on board)

03-56: **Substitution Method** Try  $T(n) \leq cn$ :

$$T(n) = 2T\left(\frac{n}{2}\right) + 1$$

$$\leq 2c\left(\frac{n}{2}\right) + 1$$

$$\leq cn + 1$$

We did not get back  $T(n) \leq cn$  – that extra +1 term means the proof is not valid. We need to get back *exactly* what we started with (see invalid proof of  $\sum_{i=1}^n i \in O(n)$  for why this is true)

03-57: **Substitution Method** Try  $T(n) \leq cn - b$ :

$$T(n) = 2T\left(\frac{n}{2}\right) + 1$$

$$\leq 2\left(c\left(\frac{n}{2}\right) - b\right) + 1$$

$$\leq cn - 2b + 1$$

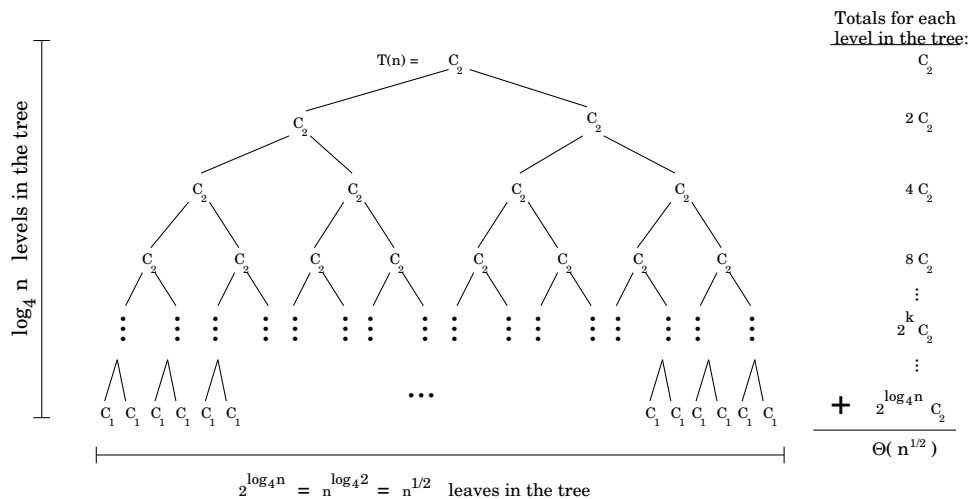
$$\leq cn - b$$

As long as  $b \geq 1$

03-58: **Master Method**

Recursion Tree for:  $T(n) = 2T(n/4) + C_2$

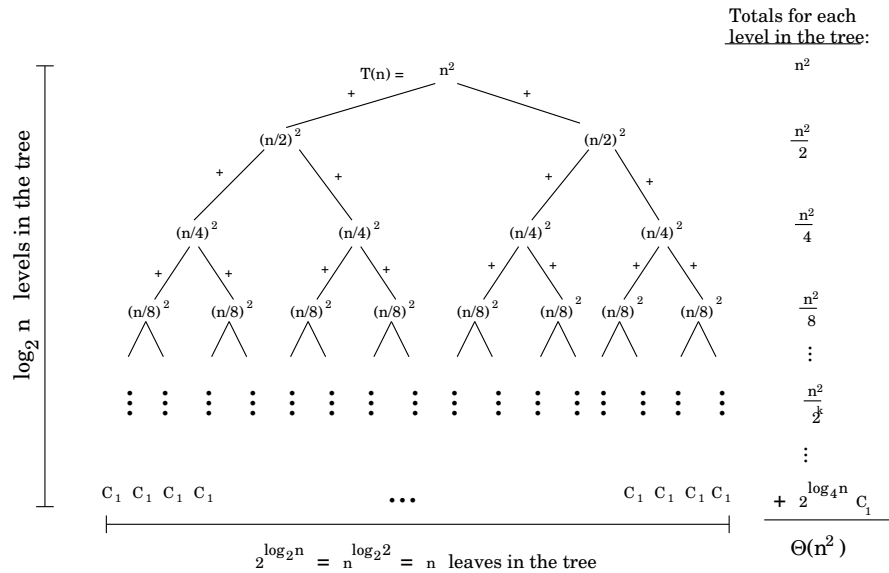
03-59: **Master Method**



03-60: **Master Method**

Recursion Tree for:  $T(n) = 2T(n/2) + n^2$

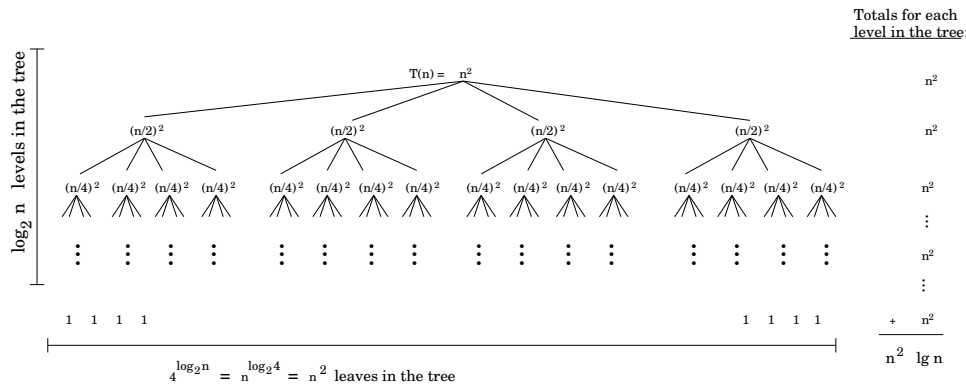
03-61: **Master Method**



03-62: Master Method

Recursion Tree for:  $T(n) = 4T(n/2) + n^2$

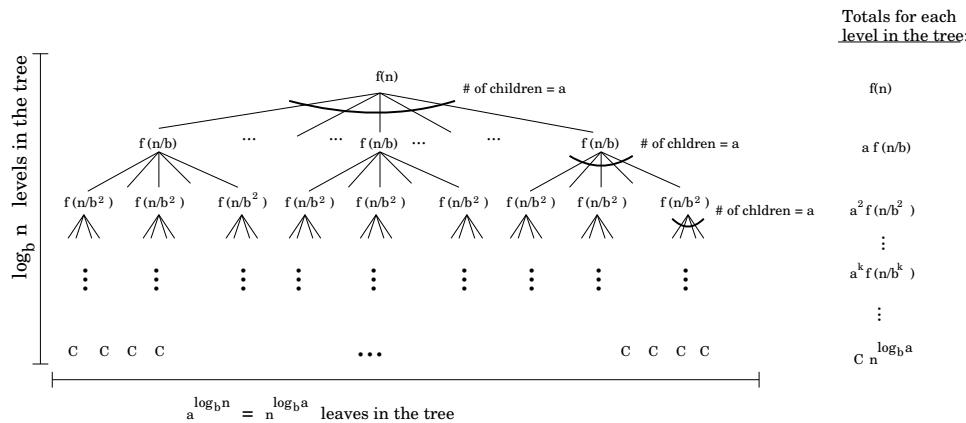
03-63: Master Method



03-64: Master Method

Recursion Tree for:  $T(n) = aT(n/b) + f(n)$

03-65: Master Method



03-66: Master Method

$$T(n) = aT(n/b) + f(n)$$

1. if  $f(n) \in O(n^{\log_b a - \epsilon})$  for some  $\epsilon > 0$ , then  $T(n) \in \Theta(n^{\log_b a})$
2. if  $f(n) \in \Theta(n^{\log_b a})$  then  $T(n) \in \Theta(n^{\log_b a} * \lg n)$
3. if  $f(n) \in \Omega(n^{\log_b a + \epsilon})$  for some  $\epsilon > 0$ , and if  $a f(n/b) \leq c f(n)$  for some  $c < 1$  and large  $n$ , then  $T(n) \in \Theta(f(n))$

03-67: **Master Method**

$$T(n) = 9T(n/3) + n$$

03-68: **Master Method**

$$T(n) = 9T(n/3) + n$$

- $a = 9, b = 3, f(n) = n$
- $n^{\log_b a} = n^{\log_3 9} = n^2$
- $n \in O(n^{2-\epsilon})$

$$T(n) = \Theta(n^2)$$

03-69: **Master Method**

$$T(n) = T(2n/3) + 1$$

03-70: **Master Method**

$$T(n) = T(2n/3) + 1$$

- $a = 1, b = 3/2, f(n) = 1$
- $n^{\log_b a} = n^{\log_{3/2} 1} = n^0 = 1$
- $1 \in O(1)$

$$T(n) = \Theta(1 * \lg n) = \Theta(\lg n)$$

03-71: **Master Method**

$$T(n) = 3T(n/4) + n \lg n$$

03-72: **Master Method**

$$T(n) = 3T(n/4) + n \lg n$$

- $a = 3, b = 4, f(n) = n \lg n$
- $n^{\log_b a} = n^{\log_4 3} = n^{0.792}$
- $n \lg n \in \Omega(n^{0.792+\epsilon})$
- $3(n/4) \lg(n/4) \leq c * n \lg n$

$$T(n) \in \Theta(n \lg n)$$

03-73: **Master Method**

$$T(n) = 2T(n/2) + n \lg n$$

03-74: **Master Method**

$$T(n) = 2T(n/2) + n \lg n$$

- $a = 2, b = 2, f(n) = n \lg n$
- $n^{\log_b a} = n^{\log_2 2} = n^1$

Master method does not apply!

$n^{1+\epsilon}$  grows faster than  $n \lg n$  for any  $\epsilon > 0$

Logs grow *incredibly* slowly!  $\lg n \in o(n^\epsilon)$  for any  $\epsilon > 0$