

Compilers

CS414-2017S-01

Compiler Basics & Lexical Analysis

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01-0: Syllabus

- Office Hours
- Course Text
- Prerequisites
- Test Dates & Testing Policies
- Projects
 - Teams of up to 2
- Grading Policies
- Questions?

01-1: Notes on the Class

- I tend to talk quickly. Don't be afraid to ask me to slow down!
- We will cover some pretty complex stuff here, which can be difficult to get the first (or even the second) time. *ASK QUESTIONS*
- While specific questions are always preferred, "I don't get it" is always an acceptable question. I am always happy to stop, re-explain a topic in a different way.
 - If you are confused, I can *guarantee* that at least 1 other person in the class would benefit from more explanation

01-2: Notes on the Class

- Projects are non-trivial
 - Using new tools (lex, yacc)
 - Managing a large scale project
 - Lots of pointer chasing, with large, complex data structures

01-3: Notes on the Class

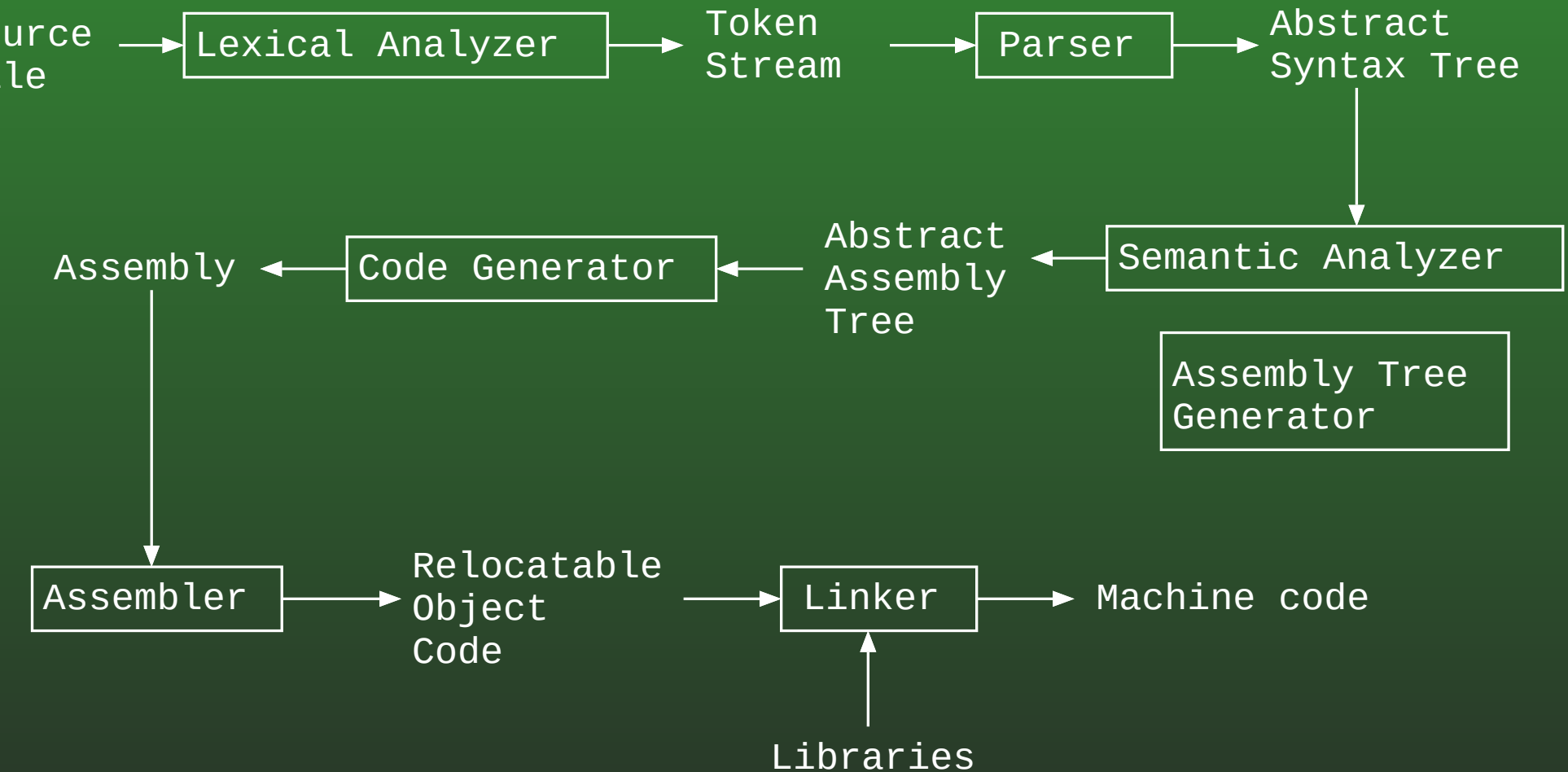
- Projects are non-trivial
 - Using new tools (lex, yacc)
 - Managing a large scale project
 - Lots of pointer chasing, with large, complex data structures
- *START EARLY!*
 - Projects will take longer than you think (especially starting with the semantic analyzer project)
- *ASK QUESTIONS!*

01-4: What is a compiler?



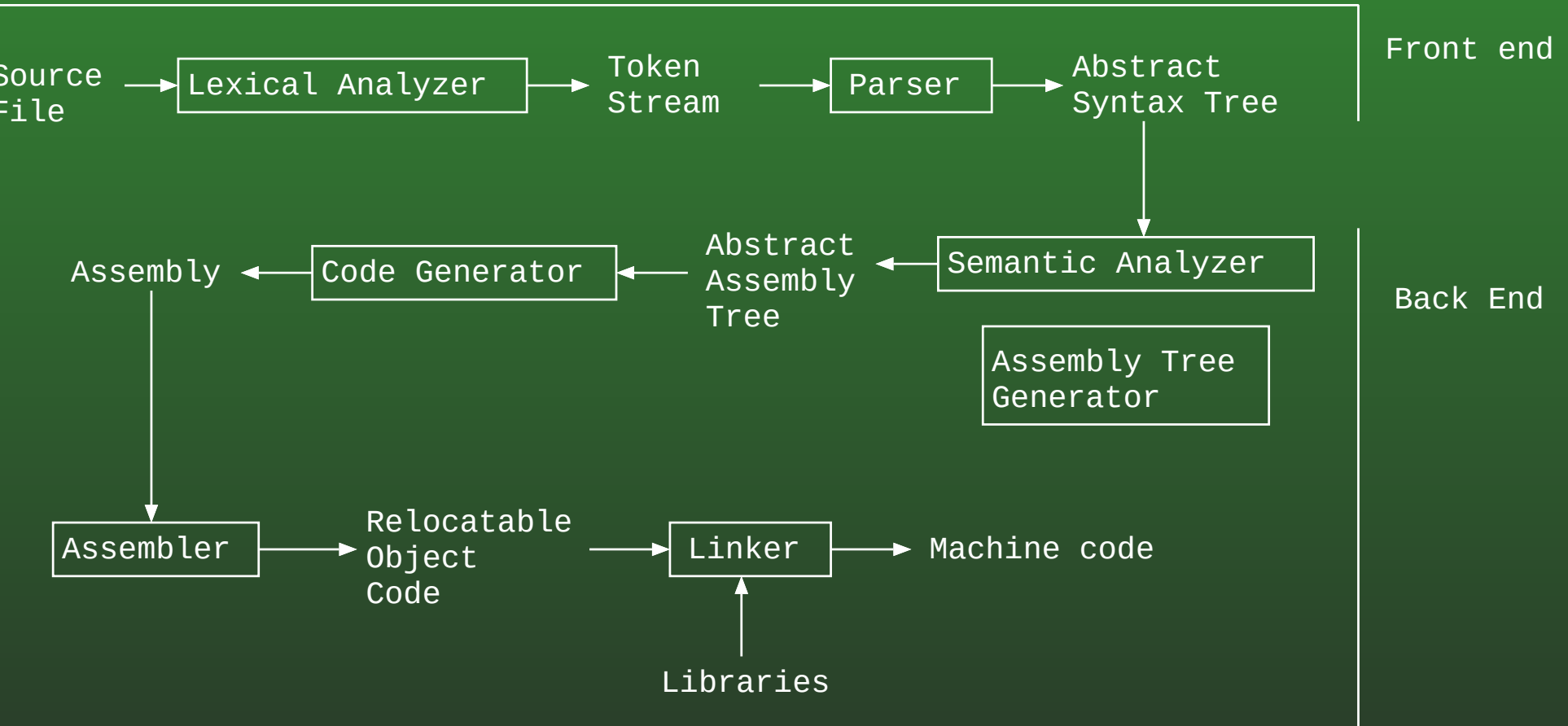
Simplified View

01-5: What is a compiler?

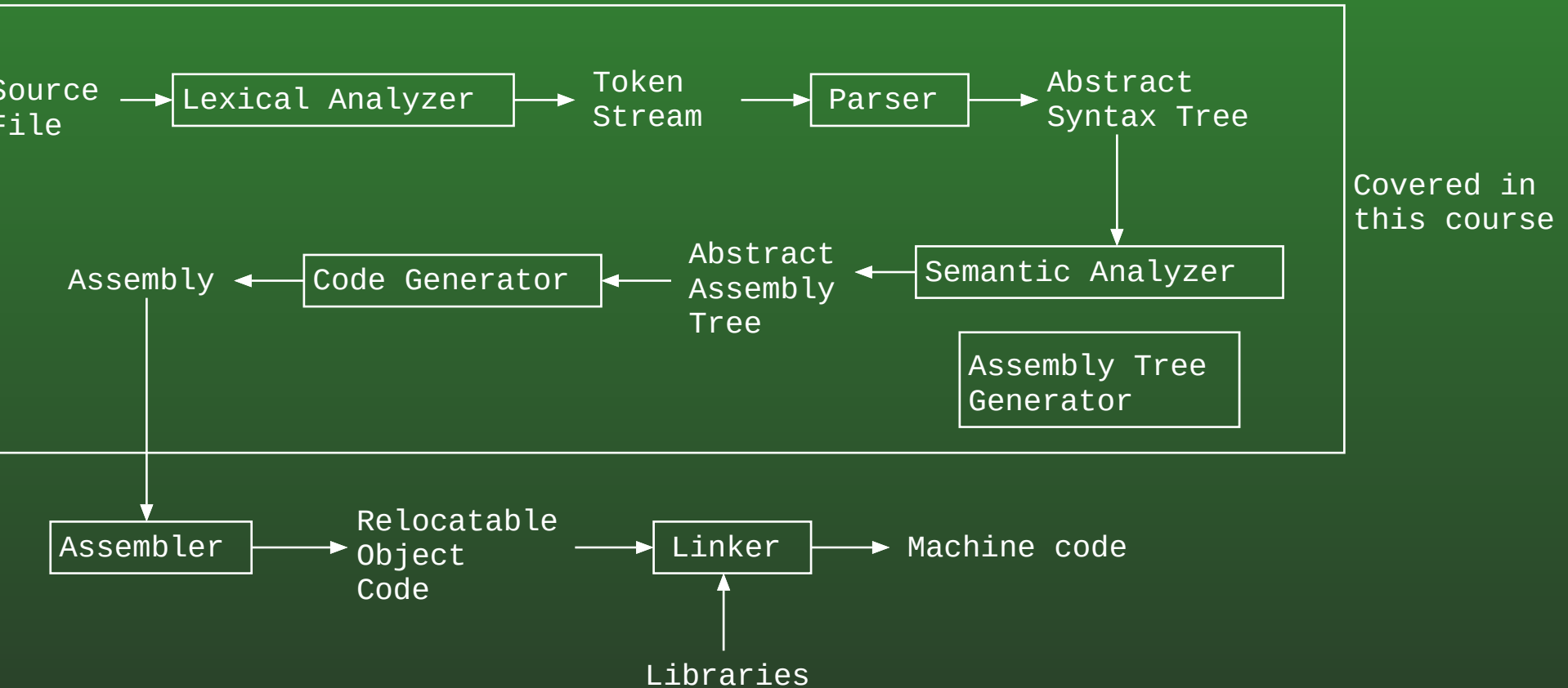


More Accurate View

01-6: What is a compiler?



01-7: What is a compiler?



01-8: Why Use Decomposition?

01-9: Why Use Decomposition?

Software Engineering!

- Smaller units are easier to write, test and debug
- Code Reuse
 - Writing a suite of compilers (C, Fortran, C++, etc) for a new architecture
 - Create a new language – want compilers available for several platforms

01-10: Lexical Analysis

- Converting input file to stream of tokens

```
void main() {  
    print(4);  
}
```

01-11: Lexical Analysis

- Converting input file to stream of tokens

void main() {	IDENTIFIER(void)
print(4);	IDENTIFIER(main)
}	LEFT-PARENTHESIS
	RIGHT-PARENTHESIS
	LEFT-BRACE
	IDENTIFIER(print)
	LEFT-PARENTHESIS
	INTEGER-LITERAL(4)
	RIGHT-PARENTHESIS
	SEMICOLON
	RIGHT-BRACE

01-12: Lexical Analysis

Brute-Force Approach

- Lots of nested if statements

```
if (c = nextchar() == 'P') {  
    if (c = nextchar() == 'R') {  
        if (c = nextchar() == 'O') {  
            if (c = nextchar() == 'G') {  
                /* Code to handle the rest of either  
                 PROGRAM or any identifier that starts  
                 with PROG  
                */  
            } else if (c == 'C') {  
                /* Code to handle the rest of either  
                 PROCEDURE or any identifier that starts  
                 with PROC  
                */  
            }  
        }  
    }  
}
```

...

01-13: Lexical Analysis

Brute-Force Approach

- Break the input file into words, separated by spaces or tabs
 - This can be tricky – not all tokens are separated by whitespace
 - Use string comparison to determine tokens

01-14: Deterministic Finite Automata

- Set of states
- Initial State
- Final State(s)
- Transitions

DFA for else, end, identifiers

Combine DFA

01-15: DFAs and Lexical Analyzers

- Given a DFA, it is easy to create C code to implement it
- DFAs are easier to understand than C code
 - Visual – almost like structure charts
- ... However, creating a DFA for a complete lexical analyzer is still complex

01-16: Automatic Creation of DFAs

We'd like a tool:

- Describe the tokens in the language
- Automatically create DFA for tokens
- Then, automatically create C code that implements the DFA

We need a method for describing tokens

01-17: Formal Languages

- **Alphabet Σ :** Set of all possible symbols (characters) in the input file
 - Think of Σ as the set of symbols on the keyboard
- **String w :** Sequence of symbols from an alphabet
- **String length $|w|$** Number of characters in a string: $|\text{car}| = 3$, $|\text{abba}| = 4$
 - **Empty String ϵ :** String of length 0: $|\epsilon| = 0$
- **Formal Language:** Set of strings over an alphabet

Formal Language $\neq$ Programming language – Formal Language is only a set of strings.

01-18: Formal Languages

Example formal languages:

- Integers $\{0, 23, 44, \dots\}$
- Floating Point Numbers $\{3.4, 5.97, \dots\}$
- Identifiers $\{\text{foo}, \text{bar}, \dots\}$

01-19: Language Concatenation

- **Language Concatenation** Given two formal languages L_1 and L_2 , the concatenation of L_1 and L_2 , $L_1L_2 = \{xy | x \in L_1, y \in L_2\}$

For example:

$\{\text{fire, truck, car}\} \{\text{car, dog}\} =$
 $\{\text{firecar, firedog, truckcar, truckdog, carcar, cardog}\}$

01-20: Kleene Closure

Given a formal language L :

$$L^0 = \{\epsilon\}$$

$$L^1 = L$$

$$L^2 = LL$$

$$L^3 = LLL$$

$$L^4 = LLLL$$

$$L^* = L^0 \cup L^1 \cup L^2 \cup \dots \cup L^n \cup \dots$$

01-21: Regular Expressions

Regular expressions are used to describe formal languages over an alphabet Σ :

Regular Expression	Language
--------------------	----------

ϵ	$L[\epsilon] = \{\epsilon\}$
------------	------------------------------

$a \in \Sigma$	$L[a] = \{a\}$
----------------	----------------

(MR)	$L[MR] = L[M]L[R]$
--------	--------------------

$(M R)$	$L[(M R)] = L[M] \cup L[R]$
---------	-----------------------------

(M^*)	$L[(M^*)] = L[M]^*$
---------	---------------------

01-22: r.e. Precedence

From highest to Lowest:

Kleene Closure *

Concatenation

Alternation |

$$ab^*c|e = (a(b^*)c) \mid e$$

01-23: Regular Expression Examples

$(a b)^*$	all strings over $\{a,b\}$
$(0 1)^*$	binary integers (with leading zeroes)
$a(a b)^*a$	all strings over $\{a,b\}$ that begin and end with a
$(a b)^*aa(a b)^*$	all strings over $\{a,b\}$ that contain aa
$b^*(abb^*)^*(a \epsilon)$	all strings over $\{a,b\}$ that do not contain aa

01-24: r.e. Shorthand

$[abcd] = (a|b|c|d)$

$[b-g] = [befg] = (b|e|f|g)$

$[b-gM-O] = [befgMNO] = (b|e|f|g|M|N|O)$

$M? = (M | \epsilon)$

$M+ = MM^*$

$.$ = Any character except newline

$"vc"$ = The string vc exactly

$4"."5$ = String 4.5 (not $4<\text{any}>5$)

01-25: r.e. Shorthand Examples

Regular Expression	Language
if	{if}
[a-z][0-9a-z]*	Set of legal identifiers
[0-9]*	Set of integers (with leading zeroes)
([0-9]*"."[0-9]*)	Set of real numbers
([0-9]*"."[0-9]+)	

01-26: Lexical Analyzer Generator

Lex is a lexical analyzer generator

- Input: Set of regular expressions (each of which describes a type of token in the language)
- Output: A lexical analyzer, which reads an input file and separates it into tokens

01-27: Structure of lex file

```
%{  
  
/* C declarations */  
  
%}  
  
/* Lex definitions */  
  
%%  
  
/* Regular expressions and actions */
```

01-28: How Lex Works

- Lex creates a function `int yylex()`, which does the following:

```
while(true) {  
    <find a rule that matches the next  
        section of the input>  
  
    <execute action associated with  
        that rule>  
}
```

- Thus, a call to `yylex()` will keep matching rules and performing actions, until an action contains a return statement

01-29: Lex Example

```
%{  
#include "tokens.h"  
%}  
  
%%  
  
for      { return FOR; }  
if       { return IF; }
```

01-30: Lex Example

```
%{  
#include "tokens.h"  
%}  
  
%%  
  
for      { return FOR; }  
if       { return IF; }  
",,"    { return COMMA; }
```


01-31: Lex Example

```
%{  
#include "tokens.h"  
%}  
  
%%  
  
for      { return FOR; }  
if       { return IF; }  
", "     { return COMMA; }  
"  "     { }  
\\n      { }
```

01-32: Lex Example

```
%{  
#include "tokens.h"  
%}  
  
%%  
  
for      { return FOR; }  
if       { return IF; }  
", "     { return COMMA; }  
"  "     { }  
\n        { }  
[0-9]+   { return INTEGER_LITERAL; }
```

01-33: Lex Example

```
%{  
#include <string.h>  
#include "tokens.h"  
%}  
%%  
  
for      { return FOR; }  
if       { return IF; }  
", "     { return COMMA; }  
"  "     { }  
\\n      { }  
[0-9]+   { yylval.integer_value.value =  
           atoi(ytex);  
           return INTEGER_LITERAL; }  
}
```

01-34: Lex Loop

Lex creates a function `yylex`, which does the following:

- Find regular expression that matches the input
- Execute the action associated with that regular expression
- Repeat, until a return statement is reached – return the appropriate value

01-35: Lex Matching

What if more than one regular expression matches?

- Find the longest possible match
 - if382 returns an IDENTIFIER, not an IF followed by an INTEGER_LITERAL
- If two matches are the same length, match the string that appears earliest in the file.
 - if returns an IF, not an IDENTIFIER (as long as the rule for IDENTIFIER appears after the rule for IF)

01-36: Lex States

- We can label each regular expression/action pair with a “state”
- Unlabeled regular expression/action pairs are assumed to be in the default INITIAL state
- Lex will only match regular expressions for the current state
- Can switch between states with a BEGIN action

01-37: Using Lex States

```
%{
#include "tokens.h"
%}

%x COMMENT

%%
else          { return ELSE; }
for           { return FOR; }
" "          { }
\n           { }
" /*"        { BEGIN(COMMENT); }
<COMMENT>" */ { BEGIN(INITIAL); }
<COMMENT>\n   { }
<COMMENT>.    { }
```

01-38: Token Positions

```
%{
#include "tokens.h"
#include <string.h>
int current_line_number = 1;
void newline() {
    current_line_number++;
}
}%
%%
else                { yylval.line_number = current_line_number;
                     return ELSE; }

" "                {      }

\n                 { newline(); }

";"                { yylval.line_number = current_line_number;
                     return SEMICOLON; }

[a-zA-Z][a-zA-Z0-9]* { yylval.string_value.line_number = current_line_number;
                     yylval.string_value.value = malloc(sizeof(char) *
                                                         (strlen(yytext)+1));
                     strcpy(yylval.string_value.value,yytext);
                     return IDENTIFIER; }
```


01-39: Random Details

- tokens.h
- Skeleton .lex file
 - Note Current_Line is defined in errors.h
- errors.h file
- main (driver) program
- makefile