



**CS 220:** Introduction to Parallel Computing

# Review: Mutexes, Condition Variables

Lecture 25

# Mutex Declaration

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- **Where** you declare your mutex is very important
- For example, what happens when each thread creates its own mutex?
  - This is basically like checking if you have the keys to your own house
- In general, mutexes should be a **shared** resource
  - Declared globally

# Mutex: Mental Model

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- You can think of a mutex as a protector of a shared resource that only one thread can access at a time
- It's the gatekeeper for your protected resource
- You'll almost always have:
  1. The mutex
  2. The variable you're protecting

# Mutex: Mental Model

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- Let's say our shared resource is the whiteboard
- Before you can write on the whiteboard, you have to ask the instructor first
- The instructor will only allow one student to write on the board at a time
  - ...if you request to use the whiteboard while someone else is already using it, then the instructor makes you wait

# Checking a Mutex

- Thus far, we've just **locked** or **unlocked** a mutex
- What happens when we try to lock a mutex that is already locked by another thread?
  - We block!
- In some cases, we want to determine whether we can lock the mutex, but move on if we cannot:
  - `pthread_mutex_trylock(&mutex)`
  - Even if the mutex is already locked by another thread, the function call returns immediately

# Back to the Whiteboard Example

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- Now, assume that I've divided the whiteboard up into four quadrants
- I can now let four students have their own part of the whiteboard at a time
- To protect the four quadrants, we could have four unique mutexes
  - This doesn't scale very well... What happens when I buy another whiteboard or divide it up more?

# Array of Mutexes

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- One approach would be to keep a big array of mutexes, one for each part of the whiteboard
- Do we really need all that complexity?
- There is, however, a better way: **condition variables**

# Condition Variables

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- To wait for something to happen, we can use **condition variables**
- Condition variables have two related functions:
  - wait – wait for the condition to become true
  - signal – inform the waiting thread that the condition has changed
- When a thread is waiting, it **blocks**
  - Just like how our MPI programs block when they are waiting for a message to come in



# Initializing Condition Variables

- Initialization is just like a mutex:

```
pthread_cond_t cond_variable =  
    PTHREAD_COND_INITIALIZER;
```

- Note: to use a condition variable, you also need a mutex
  - Why? This protects the condition variable logic

# Using Condition Variables

## Thread A:

```
pthread_mutex_lock(&mutex);
while (num_students_at_board >= 4) {
    /* Note: mutex is released here: */
    pthread_cond_wait(&cond, &mutex);
}
/* Do the work we were waiting to do! */
pthread_mutex_unlock(&mutex);
```

## Thread B:

```
pthread_mutex_lock(&mutex);
/* Do whatever thread A is waiting for us to do ... */
/* Signal the other thread! */
pthread_cond_signal(&cond);
pthread_mutex_unlock(&mutex);
```